

## 7.2 OWL

- the OWL versions use certain DL semantics:
- Base:  $\mathcal{ALC}_{\mathcal{R}^+}$ : (i.e., with transitive roles). This logic is called  $\mathcal{S}$  (reminiscent to its similarity to the modal logic  $S$ ).
- roles can be ordered hierarchically (`rdfs:subPropertyOf`;  $\mathcal{H}$ ).
- OWL Lite:  $\mathcal{SHIF}(D)$ , Reasoning in EXPTIME.
- OWL DL:  $\mathcal{SHOIN}(D)$ , decidable.  
Pellet (2007) implements  $\mathcal{SHOIQ}(D)$ . Decidability is in NEXPTIME (combined complexity wrt. TBox+ABox), but the actual complexity of a given task is constrained by the maximal used cardinality and use of nominals and inverses and behaves like the simpler classes.  
(Ian Horrocks and Ulrike Sattler: A Tableau Decision Procedure for SHOIQ(D); In IJCAI, 2005, pp. 448-453; available via <http://dblp.uni-trier.de>)
- progress in OWL 2.0 towards  $\mathcal{SROIQ}(D)$  and more datatypes ...

## OWL NOTIONS

- Classes and Properties are nodes in an RDF model, their characteristics are specified by OWL properties.

### OWL Class Definitions and Axioms (Overview)

- owl:Class
- The properties of an owl:Class (including owl:Restriction) node describe the properties of that class.

An owl:Class is required to satisfy the conjunction of all constraints (implicit: intersection) stated by its subelements.

These characterizations are roughly the same as discussed for DL class definitions:

- Constructors: owl:unionOf, owl:intersectionOf, owl:complementOf ( $\mathcal{ALC}$ )
- Enumeration Constructor: owl:oneOf (enumeration of elements;  $\mathcal{O}$ )
- Axioms rdfs:subClassOf, owl:equivalentClass,
- Axiom owl:disjointWith (also expressible in  $\mathcal{ALC}$ :  $C$  disjoint with  $D$  is equivalent to  $C \sqsubseteq \neg D$ )

## OWL NOTIONS (CONT'D)

### OWL Restriction Classes (Overview)

- owl:Restriction is a subclass of owl:Class, allowing for specification of a **constraint on one property**.
- one property is restricted by an owl:onProperty specifier and a constraint on this property:
  - ( $\mathcal{N}$ ,  $\mathcal{Q}$ ,  $\mathcal{F}$ ) owl:cardinality, owl:minCardinality or owl:maxCardinality,
  - owl:allValuesFrom ( $\forall R.C$ ), owl:someValuesFrom ( $\exists R.C$ ),
  - owl:hasValue ( $\mathcal{O}$ ),
  - including datatype restrictions for the range ( $\mathcal{D}$ )
- by defining an owl:Restriction as subclass of another owl:Restriction, multiple such constraints can be defined.

## OWL NOTIONS (CONT'D)

### OWL Property Axioms (Overview)

- atomic constructor: `rdf:Property`
- from RDFS: `rdfs:domain/rdfs:range` assertions, `rdfs:subPropertyOf`
- Axiom `owl:equivalentProperty`
- Axioms: subclasses of `rdf:Property`:  
`owl:TransitiveProperty`, `owl:SymmetricProperty`, `owl:FunctionalProperty`,  
`owl:InverseFunctionalProperty` (see Slide 238)

### OWL Individual Axioms (Overview)

- Individuals are modeled by unary classes
- `owl:sameAs`, `owl:differentFrom`, `owl:AllDifferent(o1, . . . , on)`.

## FIRST-ORDER LOGIC EQUIVALENTS

OWL : $x \in C$	DL Syntax	FOL
$C$	$C$	$C(x)$
intersectionOf( $C_1, C_2$ )	$C_1 \sqcap \dots \sqcap C_n$	$C_1(x) \wedge \dots \wedge C_n(x)$
unionOf( $C_1, C_2$ )	$C_1 \sqcup \dots \sqcup C_n$	$C_1(x) \vee \dots \vee C_n(x)$
complementOf( $C_1$ )	$\neg C_1$	$\neg C_1(x)$
oneOf( $x_1, \dots, x_n$ )	$\{x_1\} \sqcup \dots \sqcup \{x_n\}$	$x = x_1 \vee \dots \vee x = x_n$
OWL : $x \in C$ , Restriction on $P$	DL Syntax	FOL
someValuesFrom( $C'$ )	$\exists P.C'$	$\exists y : P(x, y) \wedge C'(y)$
allValuesFrom( $C'$ )	$\forall P.C'$	$\forall y : P(x, y) \rightarrow C'(y)$
hasValue( $y$ )	$\exists P.\{y\}$	$P(x, y)$
maxCardinality( $n$ )	$\leq n.P$	$\exists^{\leq n} y : P(x, y)$
minCardinality( $n$ )	$\geq n.P$	$\exists^{\geq n} y : P(x, y)$
cardinality( $n$ )	$n.P$	$\exists^=n y : P(x, y)$

## FIRST-ORDER LOGIC EQUIVALENTS (CONT'D)

OWL Class Axioms for $C$	DL Syntax	FOL
<code>rdfs:subClassOf(<math>C_1</math>)</code>	$C \sqsubseteq C_1$	$\forall x : C(x) \rightarrow C_1(x)$
<code>equivalentClass(<math>C_1</math>)</code>	$C \equiv C_1$	$\forall x : C(x) \leftrightarrow C_1(x)$
<code>disjointWith(<math>C_1</math>)</code>	$C \sqsubseteq \neg C_1$	$\forall x : C(x) \rightarrow \neg C_1(x)$
<hr/>		
OWL Individual Axioms	DL Syntax	FOL
$x_1$ <code>sameAs</code> $x_2$	$\{x_1\} \equiv \{x_2\}$	$x_1 = x_2$
$x_1$ <code>differentFrom</code> $x_2$	$\{x_1\} \sqsubseteq \neg\{x_2\}$	$x_1 \neq x_2$
<code>AllDifferent(<math>x_1, \dots, x_n</math>)</code>	$\bigwedge_{i \neq j} \{x_i\} \sqsubseteq \neg\{x_j\}$	$\bigwedge_{i \neq j} x_i \neq x_j$

## FIRST-ORDER LOGIC EQUIVALENTS (CONT'D)

OWL Properties	DL Syntax	FOL
$P$	$P$	$P(x, y)$
OWL Property Axioms for $P$	DL Syntax	FOL
<code>rdfs:range(<math>C</math>)</code>	$\top \sqsubseteq \forall P.C$	$\forall x, y : P(x, y) \rightarrow C(y)$
<code>rdfs:domain(<math>C</math>)</code>	$C \sqsubseteq \exists P.\top$	$\forall x, y : P(x, y) \rightarrow C(x)$
<code>subPropertyOf(<math>P_2</math>)</code>	$P \sqsubseteq P_2$	$\forall x, y : P(x, y) \rightarrow P_2(x, y)$
<code>equivalentProperty(<math>P_2</math>)</code>	$P \equiv P_2$	$\forall x, y : P(x, y) \leftrightarrow P_2(x, y)$
<code>inverseOf(<math>P_2</math>)</code>	$P \equiv P_2^-$	$\forall x, y : P(x, y) \leftrightarrow P_2(y, x)$
<code>TransitiveProperty</code>	$P^+ \equiv P$	$\forall x, y, z : ((P(x, y) \wedge P(y, z)) \rightarrow P(x, z))$ $\forall x, z : ((\exists y : P(x, y) \wedge P(y, z)) \rightarrow P(x, z))$
<code>FunctionalProperty</code>	$\top \sqsubseteq \leq 1P.\top$	$\forall x, y_1, y_2 : P(x, y_1) \wedge P(x, y_2) \rightarrow y_1 = y_2$
<code>InverseFunctionalProperty</code>	$\top \sqsubseteq \leq 1P^-. \top$	$\forall x, y_1, y_2 : P(y_1, x) \wedge P(y_2, x) \rightarrow y_1 = y_2$

## REPRESENTATION

- most OWL constructs have a straightforward representation in RDF/XML and N3.
- OWL in RDF/XML format: usage of class, property, and individual names:
  - as `@rdf:about` when used as identifier of a subject (`owl:Class`, `rdf:Property` and their subclasses),
  - as `@rdf:resource` as the object of a property.
- some constructs need auxiliary structures (collections):  
`owl:unionOf`, `owl:intersectionOf`, and `owl:oneOf` are based on Collections
  - representation in RDF/XML by `rdf:parseType="Collection"`.
  - representation in N3 by  $(x_1 \ x_2 \ \dots \ x_n)$
  - as RDF lists: `rdf:List`, `rdf:first`, `rdf:rest`



## EXAMPLE: PARADOX

```
<?xml version="1.0"?>
<rdf:RDF xmlns:owl="http://www.w3.org/2002/07/owl#"
  xmlns:rdf="http://www.w3.org/1999/02/22-rdf-syntax-ns#"
  xml:base="foo://bla/">
  <owl:Class rdf:about="Paradox">
    <owl:complementOf rdf:resource="Paradox"/>
  </owl:Class>
</rdf:RDF>
```

[Filename: RDF/paradox.rdf]

- without reasoner:

```
jena -t -if paradox.rdf
```

Outputs the same RDF facts in N3 without checking consistency.

- with reasoner:

```
jena -t -pellet -if paradox.rdf
```

reads the RDF file, creates a model (and checks consistency) and in this case reports that it is not consistent.

## EXAMPLE: UNION AND SUBCLASS; T-BOX REASONING

```
<?xml version="1.0"?>
<rdf:RDF xmlns:owl="http://www.w3.org/2002/07/owl#"
  xmlns:rdf="http://www.w3.org/1999/02/22-rdf-syntax-ns#"
  xmlns:f="foo://bla/"
  xml:base="foo://bla/">
  <owl:Class rdf:about="Person">
    <owl:unionOf rdf:parseType="Collection">
      <owl:Class rdf:about="Male"/>
      <owl:Class rdf:about="Female"/>
    </owl:unionOf>
  </owl:Class>
  <owl:Class rdf:about="EqToPerson">
    <owl:unionOf rdf:parseType="Collection">
      <owl:Class rdf:about="Female"/>
      <owl:Class rdf:about="Male"/>
    </owl:unionOf>
  </owl:Class>
  <f:Person rdf:about="unknownPerson"/>
</rdf:RDF>
```

[Filename: RDF/union-subclass.rdf]

## Example (Cont'd)

- print class tree (with jena -e -pellet):

```
owl:Thing
  bla:Person = bla:EqToPerson - (bla:unknownPerson)
    bla:Female
    bla:Male
```

- Male and Female are derived to be subclasses of Person.
- Person and EqToPerson are equivalent classes.
- unknownPerson is a member of Person and EqToPerson.

```
prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>
prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>
prefix owl: <http://www.w3.org/2002/07/owl#>
prefix : <foo://bla/>
select ?SC ?C ?T ?CC ?CD
from <file:union-subclass.rdf>
where {{?SC rdfs:subClassOf ?C} UNION
       { :unknownPerson rdf:type ?T } UNION
       {?CC owl:equivalentClass ?CD}}
```

[Filename: RDF/union-subclass.sparql]

## EXERCISE

Consider

```
<owl:Class rdf:about="C1">  
  <owl:intersectionOf rdf:parseType="Collection">  
    <owl:Class rdf:about="A"/>  
    <owl:Class rdf:about="B"/>  
  </owl:intersectionOf>  
</owl:Class>
```

and

```
<owl:Class rdf:about="C2">  
  <rdfs:subClassOf rdf:resource="A"/>  
  <rdfs:subClassOf rdf:resource="B"/>  
</owl:Class>
```

- give mathematical characterizations of both cases.
- discuss whether both fragments are equivalent or not.

## DISCUSSION

- Two classes are *equivalent* (wrt. the knowledge base) if they have the same interpretation in every *model* of the KB.
- $C_1$  is characterized to be the intersection of classes  $A$  and  $B$ .
- for  $C_2$ , it is asserted that  $C_1$  is a subset of  $A$  and that it is a subset of  $B$ .
- Thus there can be some  $c$  that is in  $A$ ,  $B$ ,  $C_1$ , but not in  $C_2$ .
- Thus,  $C_1$  and  $C_2$  are not equivalent.

## DISCUSSION: FORMAL NOTATION

The DL equivalent to the knowledge base (TBox) is

$$\mathcal{T} = \{C_1 \sqsubseteq (A \sqcap B), \quad C_2 \sqsubseteq A, \quad C_2 \sqsubseteq B\}$$

The First-Order Logic equivalent is

$$\mathcal{KB} = \{\forall x : A(x) \wedge B(x) \leftrightarrow C_1(x), \quad \forall x : C_2(x) \rightarrow A(x) \wedge B(x)\}$$

Thus,  $\mathcal{KB} \models \forall x : C_2(x) \rightarrow A(x) \wedge B(x)$ .

Or, in DL:  $\mathcal{T} \models C_2 \sqsubseteq C_1$ .

On the other hand,  $\mathcal{M} = (\mathcal{D}, \mathcal{I})$  with  $\mathcal{D} = \{c\}$  and

$$\mathcal{I}(A) = \{c\}, \quad \mathcal{I}(B) = \{c\}, \quad \mathcal{I}(C_1) = \{c\}, \quad \mathcal{I}(C_2) = \emptyset$$

is a model of  $\mathcal{KB}$  (wrt. first-order logic) and  $\mathcal{T}$  (wrt. DL) that shows that  $C_1$  and  $C_2$  are not equivalent.

## OWL:RESTRICTION – EXAMPLE

```
<?xml version="1.0"?>
<rdf:RDF xmlns:owl="http://www.w3.org/2002/07/owl#"
  xmlns:rdf="http://www.w3.org/1999/02/22-rdf-syntax-ns#"
  xmlns:f="foo://bla/"
  xml:base="foo://bla/">
  <owl:Class rdf:about="Parent">
    <owl:intersectionOf rdf:parseType="Collection">
      <owl:Class rdf:about="Person"/>
      <owl:Restriction>
        <owl:onProperty rdf:resource="hasChild"/>
        <owl:minCardinality>1</owl:minCardinality>
      </owl:Restriction>
    </owl:intersectionOf>
  </owl:Class>
  <f:Person rdf:about="john">
    <f:hasChild><f:Person rdf:about="alice"/></f:hasChild>
  </f:Person>
</rdf:RDF>
```

```
prefix : <foo://bla/>
select ?C
from <file:restriction.rdf>
where { :john a ?C }
```

[Filename: RDF/restriction.sparql]

[Filename: RDF/restriction.rdf]

## RESTRICTIONS ONLY AS BLANK NODES

Consider the following (bad) specification:

```
:badIdea a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1.
```

This is not allowed in OWL-DL.

Correct specification:

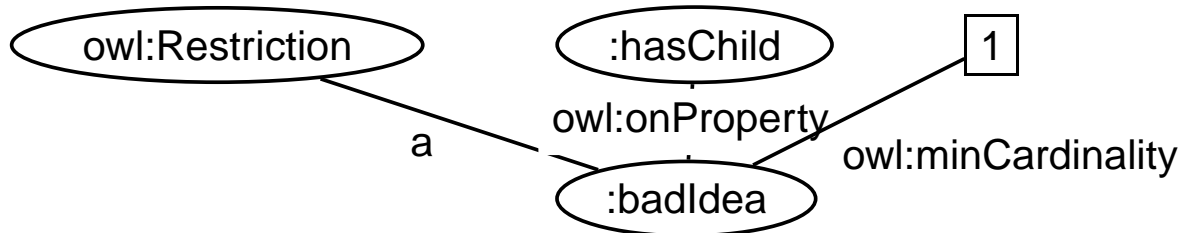
```
:badIdea owl:equivalentClass  
  [a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1].
```

Why? ... there are many reasons, for one of them see next slide.



## Restrictions Only as Blank Nodes (Cont'd)

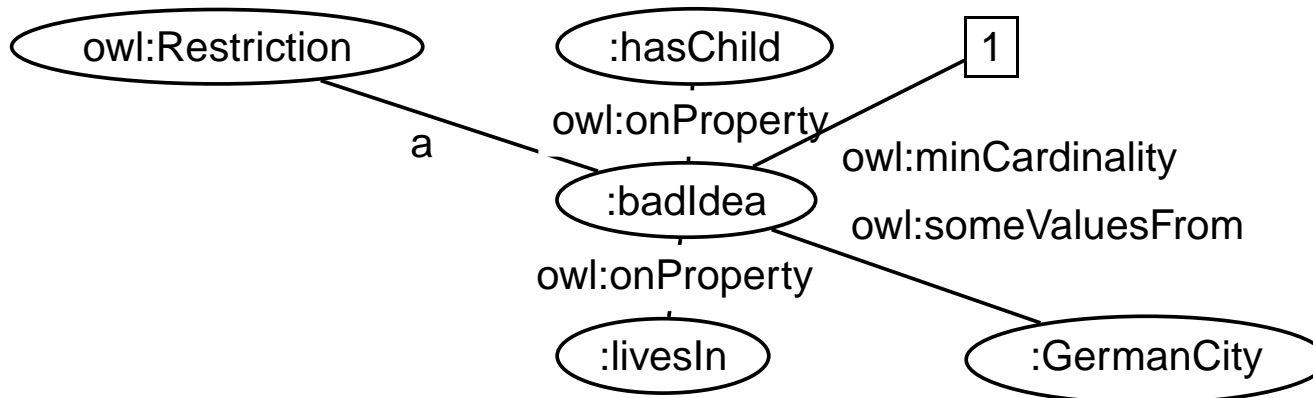
A class with two such specifications:



```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/>.
:badIdea a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1.
:badIdea a owl:Restriction; owl:onProperty :livesIn; owl:someValuesFrom :GermanCity.
```

[Filename: RDF/badIdea.n3]

- call `jena -t -pellet -if badIdea.n3:`



The two restriction specifications are messed up.

## Restrictions Only as Blank Nodes (Cont'd)

- Thus specify each Restriction specification with a separate blank node:

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/>.
:TwoRestrictions owl:intersectionOf
  ( [ a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1]
    [ a owl:Restriction; owl:onProperty :livesIn; owl:someValuesFrom :GermanCity.] ) .
```

[Filename: RDF/twoRestrictions.n3]

The DL equivalent:  $\text{TwoRestrictions} \equiv (\exists \text{hasChild}.\top) \sqcap (\exists \text{livesIn}.\text{GermanCity})$

## Another reason:

```
:AnotherBadDesignExample a owl:Restriction;
  owl:onProperty :hasChild; owl:minCardinality 1;
  rdfs:subClassOf :Person.
```

... mixes the *definition* of the Restriction with an axiom; the meaning is unclear (and the outcome is up to the strategy of the Reasoner). Obviously, the designer intended to specify an intersection,  $\text{ABDE} \equiv \exists \geq 1 \text{hasChild}.\top \sqcap \text{Person}$ , but the DL translation actually specifies a definition and an assertive axiom:  $\text{ABDE} \equiv \exists \geq 1 \text{hasChild}.\top \wedge \text{ABDE} \sqsubseteq \text{Person}$

## MULTIPLE RESTRICTIONS ON A PROPERTY

- “All persons that have at least two children, and one of them is male”

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix : <foo://bla/>.
### Test: multiple restrictions: the cardinality condition is then ignored
:HasTwoChildrenOneMale owl:intersectionOf (:Person
  [ a owl:Restriction; owl:onProperty :hasChild;
    owl:someValuesFrom :Male; owl:minCardinality 2]).
:name a owl:FunctionalProperty.
:Male rdfs:subClassOf :Person; owl:disjointWith :Female.
:Female rdfs:subClassOf :Person.
:kate a :Female; :name "Kate"; :hasChild :john.
:john a :Male; :name "John";
  :hasChild [a :Female; :name "Alice"], [a :Male; :name "Bob"].
:sue a :Female; :name "Sue";
  :hasChild [a :Female; :name "Anne"], [a :Female; :name "Barbara"].
```

```
prefix : <foo://bla/>
select ?X
from <file:restriction-double.n3>
where {?X a :HasTwoChildrenOneMale}
```

[Filename: RDF/restriction-double.sparql]

[Filename: RDF/restriction-double.n3]

- The cardinality condition is ignored in this case (Result: John and Sue).
- Solution: intersection of restrictions

## MULTIPLE RESTRICTIONS ON A PROPERTY

- “All persons that have at least two children, and one of them is male”

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix : <foo://bla/>.
:HasTwoChildrenOneMale owl:intersectionOf (:Person
  [ a owl:Restriction; owl:onProperty :hasChild; owl:someValuesFrom :Male]
  [ a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 2]).
:name a owl:FunctionalProperty.
:Male rdfs:subClassOf :Person; owl:disjointWith :Female.
:Female rdfs:subClassOf :Person.
:kate a :Female; :name "Kate"; :hasChild :john.
:john a :Male; :name "John";
  :hasChild [a :Female; :name "Alice"], [a :Male; :name "Bob"].
:sue a :Female; :name "Sue";
  :hasChild [a :Female; :name "Anne"], [a :Female; :name "Barbara"].
```

```
prefix : <foo://bla/>
select ?X
from <file:intersect-restrictions.n3>
where {?X a :HasTwoChildrenOneMale}
[Filename: RDF/intersect-restrictions.sparql]
```

[Filename: RDF/intersect-restrictions.n3]

- Note: this is different from Qualified Range Restrictions such as “All persons that have at least two male children” – see Slide 358.

## USE OF A DERIVED CLASS

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix : <foo://bla/names#>.
:kate :name "Kate"; :child :john.
:john :name "John"; :child :alice.
:alice :name "Alice".
:Parent a owl:Class; owl:equivalentClass
  [ a owl:Restriction; owl:onProperty :child; owl:minCardinality 1 ].
:Grandparent owl:equivalentClass
  [a owl:Restriction; owl:onProperty :child; owl:someValuesFrom :Parent].
```

[Filename: RDF/grandparent.n3]

```
prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>
prefix : <foo://bla/names#>
select ?A ?B
from <file:grandparent.n3>
where {{?A a :Parent} UNION
       {?B a :Grandparent} UNION
       {:Grandparent rdfs:subClassOf :Parent}}
```

[Filename: RDF/grandparent.sparql]

## UNION AS $A \sqcup B \equiv \neg((\neg A) \sqcap (\neg B))$

```
@prefix : <foo://bla/>.
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
:A rdf:type owl:Class.      :B rdf:type owl:Class.
:Union1 owl:unionOf (:A :B).
:CompA owl:complementOf :A.  :CompB owl:complementOf :B.
:IntersectComps owl:intersectionOf (:CompA :CompB).
:Union2 owl:complementOf :IntersectComps.
:x rdf:type :A.              :x rdf:type :B.
:y rdf:type :CompA. # a negative assertion y not in A would be better -> OWL 2
:y rdf:type :CompB. [Filename: RDF/union.n3]
```

```
prefix owl: <http://www.w3.org/2002/07/owl#>
prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>
prefix : <foo://bla/>
select ?X ?C ?D
from <file:union.n3> [Filename: RDF/union.sparql]
where {{?X rdf:type ?C} UNION {:Union1 owl:equivalentClass ?D}}
```

## NON-EXISTENCE OF A PROPERTY

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
```

```
@prefix : <foo://bla/names#>.
```

```
:kate a :Person; :name "Kate"; :hasChild :john.
```

```
:john a :Person; :name "John"; :hasChild :alice, :bob.
```

```
:alice a :Person; :name "Alice".
```

```
:bob a :Person; :name "Bob".
```

```
:name a owl:FunctionalProperty.
```

```
:ChildlessA owl:intersectionOf (:Person
```

```
  [ a owl:Restriction; owl:onProperty :hasChild; owl:maxCardinality 0]).
```

```
:ChildlessB owl:intersectionOf (:Person
```

```
  [ a owl:Restriction; owl:onProperty :hasChild; owl:allValuesFrom owl:Nothing]).
```

```
:ParentA owl:intersectionOf (:Person [owl:complementOf :ChildlessA]).
```

```
:ParentB owl:intersectionOf (:Person
```

```
  [ a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1]).
```

```
prefix : <foo://bla/names#>
select ?X ?Y
from <file:childless.n3>
where {{?X a :ChildlessA}
       union {?Y a :ParentA}}
```

[Filename: RDF/childless.sparql]

[Filename: RDF/childless.n3]

- export class tree: ChildlessA and ChildlessB are equivalent,
- note: due to the Open World Assumption, both classes are empty.
- Persons where no children are known are neither in ChildlessA or in Parent!

## TBox vs. ABox

DL makes a clean separation between TBox and ABox vocabulary:

- TBox: RDFS/OWL vocabulary for information about classes and properties (further partitioned into definitions and axioms),
- ABox: Domain vocabulary and `rdf:type`.

RDFS/OWL allows to mix everything in a set of triples.



## NOMINALS

- use individuals (that usually occur only in the ABox) in the TBox:
- as individuals `:Italy` (that are often implemented in the reasoner as unary classes) with *property* `owl:hasValue o` (the class of all things such that  $\{?x \text{ property } o\}$  holds).
- in enumerated classes *class* `owl:oneOf (o1, ..., on)` (*class* is defined to be the set  $\{o_1, \dots, o_n\}$ ).

### Difference to Reification

- Reification treats a class (e.g. `:Penguin`) or a property as an individual (`:Penguin a :Species`)
  - without reification, only specific RDFS and OWL properties are allowed for classes and properties only
  - reification assigns properties from an application domain to classes and properties.
- useful when talking about metadata notions,
- risk: allows for paradoxes

## USING NOMINALS: ITALIAN CITIES

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix mon: <http://www.semwebtech.org/mondial/10/meta#>.
@prefix it: <foo://italian/>.
it:Italy owl:sameAs <http://www.semwebtech.org/mondial/10/countries/I/>.
it:ItalianProvince owl:intersectionOf
  (mon:Province
   [a owl:Restriction; owl:onProperty mon:isProvinceOf;
    owl:hasValue it:Italy]).          # Nominal: an individual in a TBox axiom
it:ItalianCity owl:intersectionOf
  (mon:City
   [a owl:Restriction;
    owl:onProperty mon:belongsTo;
    owl:someValuesFrom it:ItalianProvince]).          [Filename: RDF/italiancities.n3]
```

```
prefix it: <foo://italian/>
select ?X
from <file:mondial-meta.n3>
from <file:mondial-europe.n3>
from <file:italiancities.n3>
where {?X a it:ItalianCity}          [Filename: RDF/italiancities.sparql]
```

## AN ONTOLOGY IN OWL

Consider the Italian-English-Ontology from Slide 109.

```
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix f: <foo://bla/>.
f:Italian rdfs:subClassOf f:Person;
  owl:disjointWith f:English;
  owl:unionOf (f:Lazy f:Latin Lover).
f:Lazy owl:disjointWith f:Latin Lover.
f:English rdfs:subClassOf f:Person.
f:Gentleman rdfs:subClassOf f:English.
f:Hooligan rdfs:subClassOf f:English.
f:Latin Lover rdfs:subClassOf f:Gentleman.
```

[Filename: RDF/italian-english.n3]

Class tree with jena -e:

```
owl:Thing
  bla:Person
    bla:English
      bla:Hooligan
      bla:Gentleman
        bla:Italian = bla:Lazy
        owl:Nothing = bla:Latin Lover
```

- Latin Lover is empty,  
thus Italian  $\equiv$  Lazy.

## Italians and Englishmen (Cont'd)

- the conclusions apply to the instance level:

```
@prefix : <foo://bla/>.
:mario a :Italian.
```

[Filename: RDF/mario.n3]

```
prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>
prefix : <foo://bla/>
select ?C
from <file:italian-english.n3>
from <file:mario.n3>
where { :mario rdf:type ?C }
```

[Filename: RDF/italian-english.sparql]

## AN ONTOLOGY IN OWL

Consider the Italian-Ontology from Slide 110.

```
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix it: <foo://italian/>.
it:Bolzano owl:sameAs
<http://www.semwebtech.org/mondial/10/countries/I/provinces/TrentinoAltoAdige/cities/Bolzano/>
it:Italian owl:intersectionOf
  (it:Person
    [a owl:Restriction; owl:onProperty it:livesIn;
     owl:someValuesFrom it:ItalianCity]);
  owl:unionOf (it:Lazy it:Mafioso it:Latin Lover).
it:Professor rdfs:subClassOf it:Person.
it:Lazy owl:disjointWith it:ItalianProf;
  owl:disjointWith it:Mafioso;
  owl:disjointWith it:Latin Lover.
it:Mafioso owl:disjointWith it:ItalianProf;
  owl:disjointWith it:Latin Lover.
it:ItalianProf owl:intersectionOf (it:Italian it:Professor).
it:enrico a it:Professor; it:livesIn it:Bolzano.      [Filename: RDF/italian-prof.n3]
```

```
prefix : <foo://italian/>
select ?C
from <file:italian-prof.n3>
from <file:mondial-meta.n3>
from <file:mondial-europe.n3>
from <file:italiancities.n3>
where { :enrico a ?C }
```

[Filename: RDF/italian-prof.sparql]

## ENUMERATED CLASSES: ONEOF

```
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
```

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
```

```
@prefix mon: <http://www.semwebtech.org/mondial/10/meta#>.
```

```
<bla:MontanunionMembers> owl:intersectionOf
```

```
  (mon:Country
```

```
    [owl:oneOf
```

```
      (<http://www.semwebtech.org/mondial/10/countries/NL/>
```

```
      <http://www.semwebtech.org/mondial/10/countries/B/>
```

```
      <http://www.semwebtech.org/mondial/10/countries/L/>
```

```
      <http://www.semwebtech.org/mondial/10/countries/F/>
```

```
      <http://www.semwebtech.org/mondial/10/countries/I/>
```

```
      <http://www.semwebtech.org/mondial/10/countries/D/>))]).
```

```
<bla:Result> owl:intersectionOf (mon:Organization
```

```
  [a owl:Restriction; owl:onProperty mon:hasMember;
```

```
  owl:someValuesFrom <bla:MontanunionMembers>]).
```

```
select ?X
from <file:montanunion.n3>
from <file:mondial-europe.n3>
from <file:mondial-meta.n3>
where {?X a <bla:Result>}
```

[RDF/montanunion.sparql]

- Query: all organizations that **share** a member with the Montanunion.

## ONEOF (EXAMPLE CONT'D)

- previous example: “all organizations that share a member with the Montanunion.”  
(DL:  $x \in \exists \text{hasMember.MontanunionMembers}$ )
- “all organizations where *all* members are also members of the Montanunion.”  
(DL:  $x \in \forall \text{hasMember.MontanunionMembers}$ )

The result is empty (although there is e.g. BeNeLux) due to open world: it is not known whether there may exist additional members of e.g. BeNeLux.

Only if the membership is “closed”, results can be proven:

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix mon: <http://www.semwebtech.org/mondial/10/meta#>.
<http://www.semwebtech.org/mondial/10/organizations/Benelux/>
  a [a owl:Restriction;
      owl:onProperty mon:hasMember; owl:cardinality 3].
<bla:SupersetOfMU> owl:intersectionOf (mon:Organization
  [a owl:Restriction; owl:onProperty mon:hasMember;
    owl:allValuesFrom <bla:MontanunionMembers>]).
mon:name a owl:FunctionalProperty. # not yet given in th
```

```
select ?X
from <file:montanunion.n3>
from <file:montanunion2.n3>
from <file:mondial-europe.n3>
from <file:mondial-meta.n3>
where {?X a <bla:SupersetOfMU>}
```

[Filename: RDF/montanunion2.n3]

[RDF/montanunion2.sparql]

## ONEOF (EXAMPLE CONT'D)

- “all organizations that cover *all* members of the Montanunion.”  
(DL:  $x \in \forall \text{hasMember.MontanunionMembers}$ )  
owl:oneOf is closed, so there is no problem with the universal quantifier.

```
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix mon: <http://www.semwebtech.org/mondial/10/meta#>.
<bla:EUMembers> owl:equivalentClass [a owl:Restriction;
    owl:onProperty mon:isMember; owl:hasValue
    <http://www.semwebtech.org/mondial/10/organizations/EU/>].
```

[Filename: RDF/montanunion3.n3]

```
prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>
select ?X # ?Y ?Z
from <file:montanunion.n3>
from <file:montanunion3.n3>
from <file:mondial-europe.n3>
from <file:mondial-meta.n3>
where {#{?Y a <bla:EUMembers>} UNION {?Z a <bla:MontanunionMembers>} UNION
    {<bla:MontanunionMembers> rdfs:subClassOf ?X}} [RDF/montanunion3.sparql]
```



## ONE OF (EXAMPLE CONT'D)

Previous example:

- only for one organization
- defined a class that contains all members of the organization
- not possible to define a *family of classes* – one class for each organization.
- this would require a *parameterized constructor*:

“ $C_{org}$  is the set of all members of  $org$ ”

Second-Order Logic: each organization can be seen as a unary predicate (=set):

$\forall Org : Org(c) \leftrightarrow \text{hasMember}(Org, c)$

or in F-Logic syntax: `C isa Org :- Org:organization[hasMember->C]`

yields e.g.

$I(eu) = \{germany, france, \dots\}$ ,

$I(nato) = \{usa, canada, germany, \dots\}$

Recall that “organization” itself is a predicate:

$I(organization) = \{eu, nato, \dots, \}$

So we have again reification: organizations are both first-order-individuals and classes.

## CONVENIENCE CONSTRUCT: owl:ALLDIFFERENT

- owl:oneOf defines a class as a closed set;
- in owl:oneOf ( $x_1, \dots, x_n$ ), two items may be the same (open world),

### owl:AllDifferent

- Triples of the form `:a owl:differentFrom :b` state that two individuals are different.

For a database with  $n$  elements, one needs

$$(n - 1) + (n - 2) + \dots + 2 + 1 = \sum_{i=1..n} i = n \cdot (n + 1) / 2 = O(n^2) \text{ such statements.}$$

- The –purely syntactical– convenience construct

`[ a owl:AllDifferent; owl:members ( $r_1 r_2 \dots r_n$ ) ]`

provides a shorthand notation.

- it is *immediately* translated into the set of all statements

$$\{r_i \text{ owl:differentFrom } r_j \mid i \neq j \in 1..n\}$$

- `[ a owl:AllDifferent; owl:members (...) ]`

is to be understood as a (blank node) that acts as a *specification* that the listed things are different that does not actually exist in the model.

## [SYNTAX] OWL:ALLDIFFERENT IN RDF/XML

```
<?xml version="1.0"?>
<rdf:RDF xmlns:owl="http://www.w3.org/2002/07/owl#"
  xmlns:rdf="http://www.w3.org/1999/02/22-rdf-syntax-ns#"
  xmlns:f="foo://bla/" xml:base="foo://bla/">
  <owl:Class rdf:about="Foo">
    <owl:equivalentClass> <owl:Class>
      <owl:oneOf rdf:parseType="Collection">
        <owl:Thing rdf:about="a"/> <owl:Thing rdf:about="b"/>
        <owl:Thing rdf:about="c"/> <owl:Thing rdf:about="d"/>
      </owl:oneOf>
    </owl:Class> </owl:equivalentClass>
  </owl:Class>
  <owl:AllDifferent> <!-- use like a class, but is only a shorthand -->
    <owl:members rdf:parseType="Collection">
      <owl:Thing rdf:about="a"/> <owl:Thing rdf:about="b"/>
      <owl:Thing rdf:about="c"/> <owl:Thing rdf:about="d"/>
    </owl:members>
  </owl:AllDifferent>
  <owl:Thing rdf:about="a"> <owl:sameAs rdf:resource="b"/> </owl:Thing>
</rdf:RDF>
```

[Filename: RDF/alldiff.rdf]

```
prefix : <foo://bla/>
prefix owl:
  <http://www.w3.org/2002/07/owl#>
select ?X ?Y
from <file:alldiff.rdf>
where {?X a :Foo}
```

[Filename: RDF/alldiffxml.sparql]

- AllDifferent is only intended as a kind of command to the application to add all pairwise “different-from” statements, it does not actually introduce itself as triples:
- trying to add things like `{?X a owl:AllDifferent}` to the query results in an error.

## [SYNTAX] OWL:ALLDIFFERENT IN N3

Example:

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/>.
:Foo owl:equivalentClass [ owl:oneOf (:a :b :c :d) ].
# noth the following syntax are equivalent and correct:
[ a owl:AllDifferent; owl:members (:a :b)].
[] a owl:AllDifferent; owl:members (:c :d).
:a owl:sameAs :b.
# :b owl:sameAs :d.
```

[Filename: RDF/alldiff.n3]

```
prefix : <foo://bla/>
select ?X ?Y
from <file:alldiff.n3>
where {?X a :Foo}
```

[Filename: RDF/alldiff.sparql]

## ONEOF: A TEST

- owl:oneOf defines a “closed set” (use with anonymous class; see below):
- note that in owl:oneOf ( $x_1, \dots, x_n$ ), two items may be the same (open world),
- optional owl:AllDifferent to guarantee that ( $x_1, \dots, x_n$ ) are pairwise distinct.

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/>.
:Person owl:equivalentClass [ owl:oneOf (:john :alice :bob) ].
# :john owl:sameAs :alice. # to show that it is consistent that they are the same
[] a owl:AllDifferent; owl:members (:john :alice :bob). # to guarantee distinctness
# :name a owl:FunctionalProperty. # this also guarantees distinctness ;)
:john :name "John".
:alice :name "Alice".
:bob :name "Bob".
:d a :Person.
:d owl:differentFrom :john; owl:differentFrom :alice.
# :d owl:differentFrom :bob. ### adding this makes the ontology inconsistent
```

[Filename: RDF/three.n3]

- Who is :d?

## oneOf: a Test (cont'd)

Who is :d?

- check the class tree:  
bla:Person - (bla:bob, bla:alice, bla:d, bla:john)

- and ask it:

```
prefix : <foo://bla/>
select ?N
from <file:three.n3>
where { :d :name ?N }
```

[Filename: RDF/three.sparql]

The answer is ?N/"Bob".

## ANSWER SETS TO QUERIES AS AD-HOC CONCEPTS

- all organizations whose headquarter city is a capital:

```
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#> .
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <http://www.semwebtech.org/mondial/10/meta#> .
:CountryCapital owl:intersectionOf
  (:City [a owl:Restriction; owl:onProperty :isCapitalOf;
         owl:someValuesFrom :Country])).
<bla:Result> owl:intersectionOf
  (:Organization [a owl:Restriction; owl:onProperty :hasHeadq;
                 owl:someValuesFrom :CountryCapital])). [Filename: RDF/organizations-query.n3]
```

```
prefix : <http://www.semwebtech.org/mondial/10/meta#>
select ?A ?N
from <file:organizations-query.n3>
from <file:mondial-europe.n3>
from <file:mondial-meta.n3>
where {?X a <bla:Result> . ?X :abbrev ?A . ?X :hasHeadq ?C . ?C :name ?N}
```

[Filename:RDF/organizations-query.sparql]

## HOW TO DEAL WITH OWL:ALLVALUESFROM IN AN OPEN WORLD?

- “forall items” is only applicable if additional items can be excluded ( $\Rightarrow$  locally closed predicate/property),
- often, RDF data is generated from a database,
- certain predicates can be closed by defining restriction classes with maxCardinality.



## OWL:ALLVALUESFROM

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/names#>.
[ a :Male; a :ThreeChildrenParent; :name "John";
  :child [a :Female; :name "Alice"], [a :Male; :name "Bob"],
         [a :Female; :name "Carol"]].
[ a :Female; a :TwoChildrenParent; :name "Sue";
  :child [a :Female; :name "Anne";], [a :Female; :name "Barbara"]].
:name a owl:FunctionalProperty.
:OneChildParent owl:equivalentClass [a owl:Restriction;
  owl:onProperty :child; owl:cardinality 1].
:TwoChildrenParent owl:equivalentClass [a owl:Restriction;
  owl:onProperty :child; owl:cardinality 2].
:ThreeChildrenParent owl:equivalentClass [a owl:Restriction;
  owl:onProperty :child; owl:cardinality 3].
:OnlyFemaleChildrenParent owl:equivalentClass [a owl:Restriction;
  owl:onProperty :child; owl:allValuesFrom :Female].
```

```
prefix : <foo://bla/names#>
select ?N
from <file:allvaluesfrom.n3>
where {?X :name ?N .
       ?X a :OnlyFemaleChildrenParent}
```

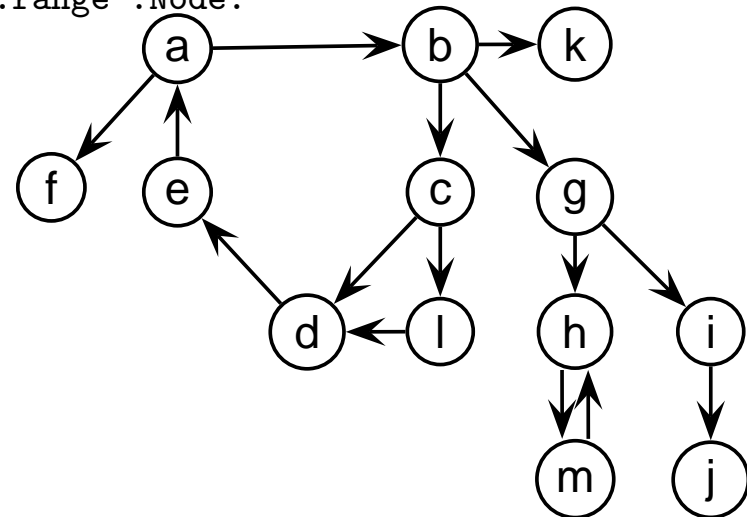
[Filename: RDF/allvaluesfrom.sparql]

[Filename: RDF/allvaluesfrom.n3]

## EXAMPLE: WIN-MOVE-GAME IN OWL

```
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/>.

:Node a owl:Class; owl:equivalentClass
  [ a owl:Class; owl:oneOf (:a :b :c :d :e :f :g :h :i :j :k :l :m)].
:edge a owl:ObjectProperty; rdfs:domain :Node; rdfs:range :Node.
:out a owl:DatatypeProperty.
:a a :Node; :out 2; :edge :b, :f.
:b a :Node; :out 3; :edge :c, :g, :k.
:c a :Node; :out 2; :edge :d, :l.
:d a :Node; :out 1; :edge :e.
:e a :Node; :out 1; :edge :a.
:f a :Node; :out 0.
:g a :Node; :out 2; :edge :i, :h.
:h a :Node; :out 1; :edge :m.
:i a :Node; :out 1; :edge :j.
:j a :Node; :out 0.
:k a :Node; :out 0.
:l a :Node; :out 1; :edge :d.
:m a :Node; :out 1; :edge :h.
```



[Filename: RDF/winmove-graph.n3]

## Win-Move-Game in OWL – the Game Axioms

“If a player cannot move, he loses.”

Which nodes are WinNodes, which one are LoseNodes (i.e., the player who has to move wins/loses)?

- if a player can move to some LoseNode (for the other), he will win.
- if a player can move only to WinNodes (for the other), he will lose.
- recall that there can be nodes that are neither WinNodes nor LoseNodes.

```
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/>.

:WinNode a owl:Class; owl:intersectionOf ( :Node
  [a owl:Restriction; owl:onProperty :edge; owl:someValuesFrom :LoseNode]).
:LoseNode a owl:Class; owl:intersectionOf ( :Node
  [a owl:Restriction; owl:onProperty :edge; owl:allValuesFrom :WinNode]).
```

[Filename: RDF/winmove-axioms.n3]

## Win-Move-Game in OWL – Closure

```
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/>.

:DeadEndNode a owl:Class; rdfs:subClassOf :Node;
  owl:equivalentClass [ a owl:Restriction; owl:onProperty :out; owl:hasValue 0 ],
    [ a owl:Restriction; owl:onProperty :edge; owl:cardinality 0 ].

:OneExitNode a owl:Class; rdfs:subClassOf :Node;
  owl:equivalentClass [ a owl:Restriction; owl:onProperty :out; owl:hasValue 1 ],
    [ a owl:Restriction; owl:onProperty :edge; owl:cardinality 1 ].

:TwoExitsNode a owl:Class; rdfs:subClassOf :Node;
  owl:equivalentClass [ a owl:Restriction; owl:onProperty :out; owl:hasValue 2 ],
    [ a owl:Restriction; owl:onProperty :edge; owl:cardinality 2 ].

:ThreeExitsNode a owl:Class; rdfs:subClassOf :Node;
  owl:equivalentClass [ a owl:Restriction; owl:onProperty :out; owl:hasValue 3 ],
    [ a owl:Restriction; owl:onProperty :edge; owl:cardinality 3 ].
```

[Filename: RDF/winmove-closure.n3]

## Win-Move-Game in OWL: DeadEndNodes

Prove that DeadEndNodes are LoseNodes:

- obvious: Player cannot move from there
- exercise: give a formal (Tableau) proof
- The OWL Reasoner does it:

```
prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>
prefix : <foo://bla/>
select ?X
from <file:winmove-axioms.n3>
from <file:winmove-closure.n3>
where { :DeadEndNode rdfs:subClassOf :LoseNode }
```

[Filename: RDF/deadendnodes.sparql]

The answer contains an (empty) tuple which means “yes”.

## Win-Move-Game in OWL

```
prefix : <foo://bla/>
select ?W ?L ?DE
from <file:winmove-graph.n3>
from <file:winmove-axioms.n3>
from <file:winmove-closure.n3>
where {{?W a :WinNode} UNION
       {?L a :LoseNode} UNION
       {?DE a :DeadEndNode}}
```

[Filename: RDF/winmove.sparql]

## Exercise

- Is it possible to characterize DrawNodes in OWL?