Chapter 5 RDF Schema

Schema Information and Reasoning in an Open World

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ONTOLOGIES

Schema languages, metadata languages, modeling languages, ontologies ...

Classical Data Models: seen as Specification and Constraints

- every schema description defines a (more or less complete) ontology:
- ER Model (1976, entity types, attributes, relationships with cardinalities),
- UML (1997, classes with subclasses, associations with cardinalities, OCL assertions to schema components etc.).

Knowledge Representation

Metadata provides additional information about resources of a type, or about a property.

- F-Logic signatures (1989),
- ... RDFS and OWL (Web Ontology Language)

SCHEMA INFORMATION IN AN OPEN WORLD

- schema describes
 - allowed properties for an object,
 - datatype constraints for literal properties [Here: XSD literal types],
 - allowed types/classes for reference properties,
 - cardinality constraints.

Closed World: Schema as Constraints

• a database must satisfy the constraints. It must be a *model* of the formulas – *the given* data alone must be a model.

Open World: potentially incomplete knowledge

- schema information as additional information.
- since the world must be a model of the schema, some information can be *derived* from the schema.
- complain only if information is *contradictory* to the schema.

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METADATA INFORMATION: Types, Properties, and Ontologies

- Types and properties (i.e., everything that is used in a namespace) are not only "names", but are resources "somewhere in the Web", identified by a URI (used in RDF or in XML via namespaces).
- ⇒ a domain ontology describes the notions used in a namespace.

Schema and Ontology Information

- what types/classes are there,
- subclass information,
- what properties objects of a given type must/can have,
- to what types some property is applicable and what range it has,
- cardinalities of properties,
- default values,
- that some properties are transitive, symmetric, subproperties of another or excluding each other etc.

REASONING WITH RDF, RDF SCHEMA AND OWL

- theoretical details will be discussed later. The underlying thing is either
 - graph completion by rules (RDFS, OWL Lite),
 (can be translated to Datalog)
 - Description Logic (DL) Reasoning (OWL DL)
 (requires a DL reasoner, based on Tableaux techniques)
- there are reasoners available for the Jena Framework:
 - an internal one:

```
jena -q -inf -qf sparql-file
for invoking SPARQL with its internal reasoner
```

an external one:

```
(integrated into the semweb.jar used in the lecture as plug-in) jena -q -pellet -qf sparql-file for invoking SPARQL with the Pellet DL reasoner class
```

- external ones as Web Services ...

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USE OF THE JENA TOOL

- option "-t": transform (between N3 and RDF/XML)
 jena -t -pellet -if rdf-file .
 (-t is not complete for checking inconsistencies)
- option "-q": query jena -q -pellet [-if rdf-input-file] -qf query-file .
- option "-e": export the class tree (available only when the pellet reasoner is activated).
 Input is an RDF or OWL file:

```
jena -e -pellet -if rdf-file. (for checking consistency, use -e)
```

• [note: since Jan. 2008, the former [-il RDF/XML] for indicating RDF/XML vs N3 input can be omitted in most cases]

PELLET COMMANDLINE FOR SPARQL-DL QUERIES

- download pellet, set alias for pellet/pellet.sh
- see pellet help for further information
- pellet query -q query-file input-file
 - does not use FROM line(s) in SPARQL, input file must be given explicitly,
 - only one input file possible.

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ASIDE: DIG INTERFACE - DESCRIPTION LOGIC IMPLEMENTATION GROUP

- Web page: http://dl.kr.org/dig/
- agreed "tell-and-ask-interface" of DL Reasoners as Web Service:
- tell them the facts and ask them queries, or for the whole inferred model
- e.g. supported by "Pellet"
- URL for download see Lecture Web page

```
may@dbis01:~/SemWeb-Tools/pellet-1.3$ ./pellet-dig.sh &
PelletDIGServer Version 1.3 (April 17 2006)
Port: 8081
```

invoke the SPARQL Jena interface by
 jena -q -qf sparql-file -inf -r reasoner-url

```
(e.g.: http://localhost:8081)
```

 note: the tell-functionality seems to transfer only part of the knowledge → incomplete reasoning → currently not recommended.

5.1 RDF Schema Notions

- RDF is the instance level
- XML: DTDs and XML Schema for describing the structure/schema of the instance
- RDF Schema: stronger than DTD/XML "semantic-level"
 - describe the structure of the RDF instance (i.e. the "schema" of the RDF graph, not of the RDF/XML file):
 - describes the schema <u>semantically</u> in terms of an (lightweight) ontology (OWL provides then much more features):
 - * class/subclass
 - * property/subproperty, domains and ranges

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PREDEFINED RDFS CLASSES

The obvious ones

rdfs:Resource is "everything". All things described by RDF are called resources, and are instances of the class rdfs:Resource. This is the class of everything. All other classes are subclasses of this class. rdfs:Resource is an instance of rdfs:Class.

rdfs:Class: all things (resources and literals) are of rdf:type of some rdfs:Class. rdf:Properties have an rdfs:Class as domain and another rdfs:Class or rdfs:Datatype as range.

mon:Country rdf:type rdfs:Class.

An rdfs:Class is simply a resource X that is of (X rdf:type rdfs:Class). Usually, class names start with a capital letter.

Later, **owl:Class** will provide more interesting concepts of *intensionally defined* classes – like "the class father is the class of things that are male and have children".

rdf:Property is a subset of rdfs:Resource that contains all properties.

mon:capital rdf:type rdf:Property.

Usually, property names start with a non-capital letter.

[note: it's rdf:Property, not rdfs:Property!]

PREDEFINED RDFS CLASSES

rdfs:Datatype is the class of datatypes.

rdfs:Literal is the subclass of rdfs:Resource that contains all literals (i.e., values of rdfs:Datatypes).

Literals do (usually) not have a URI, but a literal representation (as already discussed for integers and strings).

E.g. the following holds

@prefix xsd: <http://www.w3.org/2001/XMLSchema#>.
xsd:int rdf:type rdfs:Datatype .

• Note that *reification* takes place here: rdfs:Datatype is both an instance of and a subclass of rdfs:Class! Each instance of rdfs:Datatype is a subclass of rdfs:Literal.

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SEMANTICS OF SUBCLASSES AND SUBPROPERTIES

rdfs:subClassOf specifies that one rdfs:Class is an rdfs:subClassOf another: for any model \mathcal{M} of the RDFS model theory,

$$\mathcal{M} \models \forall C_1, C_2 : (\mathsf{holds}(C_1, \mathsf{rdfs:subClassOf}, C_2) \rightarrow (\forall x : (\mathsf{holds}(x, \mathsf{rdf:type}, C_1) \rightarrow \mathsf{holds}(x, \mathsf{rdf:type}, C_2))))$$

rdfs:subPropertyOf specifies that one rdf:Property is an rdfs:subPropertyOf another:

$$\mathcal{M} \models \forall P_1, P_2 : (\ \mathsf{holds}(P_1, \mathsf{rdfs} : \mathsf{subPropertyOf}, P_2) \rightarrow \\ (\forall x, y : (\mathsf{holds}(x, P_1, y) \rightarrow \mathsf{holds}(x, P_2, y))))$$

SEMANTICS OF DOMAIN AND RANGE

rdfs:domain specifies that the domain of an rdf:Property is a certain rdfs:Class:

```
\mathcal{M} \models \forall C, P : (\mathsf{holds}(P, \mathsf{rdfs:domain}, C) \rightarrow (\forall x : (\exists y : \mathsf{holds}(x, P, y)) \rightarrow \mathsf{holds}(x, \mathsf{rdf:type}, C)))
```

rdfs:range specifies that the range of an rdf:Property is a certain rdfs:Class (note that rdfs:Datatype is a subclass (and an instance) of rdfs:Class):

$$\mathcal{M} \models \forall C, P : (\ \mathsf{holds}(P, \mathsf{rdfs:range}, C) \rightarrow \\ (\forall y : (\exists x : \mathsf{holds}(x, P, y)) \rightarrow \mathsf{holds}(y, \mathsf{rdf:type}, C)))$$

Exercise

Give an implementation by Datalog Rules for RDFS constructs.

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INFERENCE RULES

- The above are built-in inference rules of the RDFS Model Theory
- until now, the SPARQL query language was applied to pure RDF facts (extensional knowledge)
- for the *inference rules* (= *intensional knowledge*), a *reasoner* is required.
- Queries are then not evaluated against the fact base, but against the model of the factbase and the rules.

SUBCLASS, DOMAIN, RANGE: EXAMPLE

```
@prefix : <foo://bla/names#> .
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#> .
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
   :has_cat rdfs:domain :Person .
   :has_cat rdfs:range :Cat .
   :Person rdfs:subClassOf :LivingBeing .
   :Cat rdfs:subClassOf :LivingBeing .
   <foo://bla/persons/john> :has_cat <foo://bla/cats/garfield>.
   <foo://bla/persons/mary> rdf:type :Person.
```

[Filename: RDF/subclass.n3]

```
prefix : <foo://bla/names#>
prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>
select ?X ?T
from <file:subclass.n3>
where {?X rdf:type ?T}
```

[Filename: RDF/subclass.sparql]

activate the (internal) reasoner when invoking Jena.

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SUBCLASS, DOMAIN, RANGE: EXAMPLE (CONT'D)

Recall the previous example. Given the following facts:

```
:has_cat rdfs:domain :Person .
:has_cat rdfs:range :Cat .
:Person rdfs:subClassOf :LivingBeing .
:Cat rdfs:subClassOf :LivingBeing .
<foo://bla/persons/john> :has_cat <foo://bla/cats/garfield>.
<foo://bla/persons/mary> rdf:type :Person.
```

The domain/range information does not act as a constraint, but as information. From that knowledge, the following facts can be *inferred*:

- :has_cat implies that the subject (John) is a Person, and the object (Garfield) is a cat,
- both are thus LivingBeings.

SUBPROPERTIES

• outlook: combine it with owl:TransitiveProperty.

```
@prefix : <foo://bla/names#> .
@prefix person: <foo://bla/persons/> .
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#> .
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
   person:john :child person:alice, person:bob.
   person:kate :child person:john.
   :child rdfs:subPropertyOf :descendant.
   :descendant rdf:type owl:TransitiveProperty.
```

[Filename: RDF/descendants.n3]

```
prefix : <foo://bla/names#>
select ?X ?Y
from <file:descendants.n3>
where {?X :descendant ?Y}
```

[Filename: RDF/descendants.sparql]

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5.2 Datatypes

- Strings: xsd:string (by default, every string literal is handled as a string)
- XML Schema Simple Types xsd:int etc. can be used.
- standard notations for numeric values do not need annotation.
- required etc. for time/date values.
- Further datatypes can be defined in OWL.
- Can be used in the TBox and in the ABox (with rdfs:range).

Representation in the TBox

declare xsd prefix/entity as http://www.w3.org/2001/XMLSchema#

```
    N3: p :birthday "1999-12-31" xsd:date .
    b mon:longitude 13 xsd:int .
    b mon:longitude 13 .
```

• RDF/XML: <mon:longitude rdf:datatype="%xsd;int">13</mon:longitude>

DATATYPES: DATE

use notation from XML/XML Schema for xsd:date/time/datetime

```
@prefix : <foo://bla#> .
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix xsd: <http://www.w3.org/2001/XMLSchema#>.
:birthdate rdfs:range xsd:date.
:john a :Person; :name "John"; :age 32;
:birthdate "1970-12-31"^^xsd:date .
:alice a :Person; :name "Alice"; :birthdate "2000-01-01"^^xsd:date .
```

[Filename: RDF/datatype-date.n3]

• if "xsd:date is omitted, the ontology is detected to be inconsistent!

[RDF/datatype-date.sparql]

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STRING DATATYPES: ESCAPING

- as usual with "...\" ...", or
- using """ as delimiter, escaping inside is not necessary:

DATATYPES

• it also accepts non-existing datatypes:

```
@prefix : <foo://bla#> .
@prefix xsd: <http://www.w3.org/2001/XMLSchema#>.
:john a :Person; :name "John";
:age "35"^^xsd:integer, "36"^^xsd:bla, 37, "38".
```

[Filename: RDF/datatype-casting.n3]

• use jena -t for transform.

```
prefix : <foo://bla#>
select ?Y
from <file:datatype-casting.n3>
where {:john :age ?Y}

[RDF/datatype-casting.sparql]
```

Υ	comment
"38"	string in standard notation
37	integer in standard notation
"36"~ <http: 2001="" www.w3.org="" xmlschema#bla=""></http:>	
35	integer in standard notation

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COMPARISON

SQL

- queries only against the database (no intensional knowledge),
- equivalent to tree expressions in relational algebra, based on set theory,
- formal semantics can be given purely syntactically with the algebra,
- ⇒ in the DB lecture, we did not need logic.
 - equivalent to the relational calculus, semantics of queries can be given by the calculus. Equivalent to *nonrecursive Datalog* (cf. Slide 102) with "negation as failure" (top-down) stratification (bottom-up).

RDFS + SPARQL

- only restricted negation
- RDFS: built-in rules (positive, recursive Datalog)
- SPARQL: positive, nonrecursive Datalog
- intuitive bottom-up semantics

RDFS AXIOMATIC TRIPLES

See RDF Semantics and Model Theory, http://www.w3.org/TR/rdf-mt.

Axioms: expected to hold in any RDFS model:

rdf:type rdfs:domain rdfs:Resource .

rdfs:domain rdfs:domain rdf:Property .

rdfs:range rdfs:domain rdf:Property .

rdfs:subPropertyOf rdfs:domain rdf:Property .

rdfs:subClassOf rdfs:domain rdfs:Class .

rdf:type rdfs:range rdfs:Class .

rdfs:range rdfs:range rdfs:Class .

rdfs:range rdfs:range rdfs:Class .

rdfs:subPropertyOf rdfs:range rdf:Property .

rdfs:subPropertyOf rdfs:range rdf:Property .

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USING RDF IN THE WORLD WIDE WEB

- The (Semantic) Web is not seen as a collection of documents, but as a collection of correlated information (described via documents)
- using RDF, everybody can make statements about any resource (cf. link-bases in XLink)
 - incremental, world wide data and meta-data

rdfs:Datatype rdfs:subClassOf rdfs:Class .

- distributed RDFS.
- distributed RDF.

... and some more.

- often using only virtual resources (URIs).
- not assumed that complete information about any resource is available.
- Open world, no notion of (implicit) negation.

REASONING BASED ON RDFS

- RDF/RDFS model theory as above,
- rather simple Datalog rules, graph completion,
- queries: against the (completed) graph by matching (SPARQL).
- incomplete knowledge when reasoning: "open world assumption"

Further Aspects

- potentially inconsistent information;
- statements can be equipped with probabilities or labeled as opinions;
 fuzzy reasoning, belief revision ...
- · ... lots of artificial intelligence applications ...

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EXAMPLE/EXERCISE

Consider again the employee-manages-departments example (Slide 23).

- Give the RDF Graph.
- give the N3 triples and feed them into the Jena tool.

ADDITIONAL RDF/RDFS VOCABULARY

The rdf/rdfs namespaces provide some more vocabulary:

Like most data models, RDF provides a representation for *Collections*:

- Collections: rdf:Alt, rdf:Bag, rdf:Seq, rdf:List are collections.
 Lists have properties rdf:first (a resource) and rdf:rest (a list). Others have properties _1, _2, ... that refer to their members.
- (rdfs:Container, rdfs:member, rdfs:ContainerMembershipProperty)

... these are partially used implicitly (e.g., collections in owl:intersectionOf, owl:OneOf), but often not supported by OWL reasoners if used explicitly (see Slides 402 ff.).

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EXAMPLE: THE MONDIAL ONTOLOGY

See mondial.n3, mondial-europe.n3 and mondial-meta.n3 on the Web page.

Note that it is highly redundant: defining just rdfs:domain and rdfs:range of properties implies most of the classes (and also most of the rdfs:type relationships in mondial.n3).

```
prefix mon: <http://www.semwebtech.org/mondial/10/meta#>
prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>
select ?X
from <file:mondial.n3>
from <file:mondial-meta.n3>
where {?X rdf:type mon:Country}
```

[Filename: RDF/mondial-meta-query.sparql]

• activate Jena with reasoner (if mondial.n3 is too big, use mondial-europe.n3 instead)

Mondial is not an interesting example for RDFS (and OWL):

- it's mainly data, no intensional knowledge, no complex ontology
- for that reason it is a good example for SQL and XML.
- RDFS and OWL is interesting when information is combined and additional knowledge can be derived.

Developing Ontologies

- have an idea of the required concepts and relationships (ER, UML, ...),
- generate a (draft) n3 or RDF/XML instance,
- write a separate file for the metadata,
- load it into Jena with activating a reasoner.
- If the reasoner complains about an inconsistent ontology, check the metadata file alone. If this is consistent, and it complains only when also data is loaded:
 - it may be due to populating a class whose definition is inconsistent and that thus must be empty.
 - often it is due to wrong datatypes. Recall that datatype specification is not interpreted
 as a constraint (that is violated for a given value), but as additional knowledge.