7.2 **OWL**

- the OWL versions use certain DL semantics:
- Base: $\mathcal{ALC}_{\mathcal{R}^+}$: (i.e., with transitive roles). This logic is called \mathcal{S} (reminiscent to its similarity to the modal logic S).
- roles can be ordered hierarchically (rdfs:subPropertyOf; \mathcal{H}).
- OWL Lite: SHIF(D), Reasoning in EXPTIME.
- OWL DL: SHOIN(D), decidable.
 Pellet (2007) implements SHOIQ(D). Decidability is in NEXPTIME (combined complexity wrt. TBox+ABox), but the actual complexity of a given task is constrained by the maximal used cardinality and use of nominals and inverses and behaves like the simpler classes.
 (Ian Horrocks and Ulrike Sattler: A Tableau Decision Procedure for SHOIQ(D); In IJCAI, 2005, pp. 448-453; available via http://dblp.uni-trier.de)
- OWL 2.0 towards $\mathcal{SROIQ}(D)$ and more datatypes ...

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OWL NOTIONS; OWL-DL vs. RDF/RDFS; Model vs. Graph

- OWL is defined based on (Description Logics) model theory,
- OWL ontologies can be represented by RDF graphs,
- Only certain RDF graphs are allowed OWL-DL ontologies: those, where class names, property names, individuals etc. occur in a well-organized way.
- Reasoning works on the (Description Logic) model, the RDF graph is only a means to represent it.

(recall: RDF/RDFS "reasoning" works on the graph level)

OWL VOCABULARIES

- An OWL-DL vocabulary \mathcal{V} is a 7-tuple (= a sorted vocabulary)
 - $V = (V_{cls}, V_{objprop}, V_{dtprop}, V_{annprop}, V_{indiv}, V_{DT}, V_{lit})$:
- \mathcal{V}_{cls} is the set of URIs denoting class names,
 - <http://.../mondial/10/meta#Country>
- $V_{objprop}$ is the set of URIs denoting object property names,
 - <http://.../mondial/10/meta#capital>
- \mathcal{V}_{dtprop} is the set of URIs denoting datatype property names,
 - <http://.../mondial/10/meta#population>
- ($\mathcal{V}_{annprop}$ is the set of URIs denoting annotation property names,)
- V_{indiv} is the set of URIs denoting individuals, http://.../mondial/10/countries/D
- V_{DT} is the set of URIs denoting datatype names,
 - <http://www.w3.org/2001/XMLSchema#int>
- \mathcal{V}_{lit} is the set of literals;
- the builtin notions (=URIs) from RDF, RDFS, OWL namespaces do not belong to the vocabulary of the ontology (they are only used for describing the ontology in RDF).

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OWL INTERPRETATIONS

Since DL is a subset of FOL, the interpretation of an OWL-DL vocabulary can be given as a FOL interpretation

$$\mathcal{I} = (I_{indiv} \cup I_{cls} \cup I_{objprop} \cup I_{dtprop} \cup I_{annprop} \cup I_{DT}, \ \mathcal{U}_{obj} \cup \mathcal{U}_{DT})$$

where I interprets the vocabulary as

- I_{indiv} constant symbols (individuals),
- I_{cls} , I_{DT} unary predicates (classes and datatypes),
- $I_{objprop}$, I_{dtprop} , $I_{annprop}$ binary predicates (properties),

and the universe \mathcal{U} is partitioned into

- an object domain U_{obj}
- and a data domain \mathcal{U}_{DT} (of all values of datatypes).

OWL INTERPRETATIONS

The interpretation I is as follows:

```
each individual a \in \mathcal{V}_{indiv} to an object I(a) \in \mathcal{U}_{obj},
I_{indiv}:
               (e.g., I(\langle http://.../mondial/10/countries/D\rangle) = germany)
               each class C \in \mathcal{V}_{cls} to a set I(C) \subseteq \mathcal{U}_{obj},
I_{cls}:
               (e.g., germany \in I(\langle http://.../mondial/10/meta\#Country>))
               each datatype D \in \mathcal{V}_{DT} to a set I(D) \subseteq \mathcal{U}_{DT},
I_{DT}:
              (e.g., I(<http://www.w3.org/2001/XMLSchema#int>) = \{..., -2, -1, 0, 1, 2, ...\})
              each object property p \in \mathcal{V}_{objprop} to a binary relation I(p) \subseteq \mathcal{U}_{obj} \times \mathcal{U}_{obj},
I_{objprop}:
              (e.g., (germany, berlin) \in I(\langle http://.../mondial/10/meta#capital \rangle)
              each datatype property p \in \mathcal{V}_{dtprop} to a binary relation I(p) \subseteq \mathcal{U}_{obj} \times \mathcal{U}_D,
I_{dtprop}:
              (e.g., (germany, 83536115) \in I(\langle http://.../mondial/10/meta#population>))
              each annotation property p \in \mathcal{V}_{annprop} to a binary relation I(p) \subseteq \mathcal{U} \times \mathcal{U}.
I_{dtprop}:
```

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OWL Class Definitions and Axioms (Overview)

- owl:Class
- The properties of an owl:Class (including owl:Restriction) node describe the properties of that class.

An owl:Class is required to satisfy the conjunction of all constraints (implicit: intersection) stated about it.

These characterizations are roughly the same as discussed for DL class definitions:

- Constructors: owl:unionOf, owl:intersectionOf, owl:complementOf (\mathcal{ALC})
- Enumeration Constructor: owl:oneOf (enumeration of elements; O)
- Axioms rdfs:subClassOf, owl:equivalentClass,
- Axiom owl:disjointWith (also expressible in \mathcal{ALC} : C disjoint with D is equivalent to $C \sqsubseteq \neg D$)

OWL Notions (Cont'D)

OWL Restriction Classes (Overview)

- owl:Restriction is a subclass of owl:Class, allowing for specification of a constraint on one property.
- one property is restricted by an owl:onProperty specifier and a constraint on this property:
 - $(\mathcal{N}, \mathcal{Q}, \mathcal{F})$ owl:cardinality, owl:minCardinality or owl:maxCardinality,
 - owl:allValuesFrom ($\forall R.C$), owl:someValuesFrom ($\exists R.C$),
 - owl:hasValue (\mathcal{O}),
 - including datatype restrictions for the range (D)
- by defining intersections of owl:Restrictions, classes having multiple such constraints can be specified.

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OWL Notions (Cont'D)

OWL Property Axioms (Overview)

- Distinction between owl:ObjectProperty and owl:DatatypeProperty
- from RDFS: rdfs:domain/rdfs:range assertions, rdfs:subPropertyOf
- Axiom owl:equivalentProperty
- Axioms: subclasses of rdf:Property: owl:TransitiveProperty, owl:SymmetricProperty, owl:FunctionalProperty, owl:InverseFunctionalProperty (see Slide 327)

OWL Individual Axioms (Overview)

- Individuals are modeled by unary classes
- owl:sameAs, owl:differentFrom, owl:AllDifferent(o₁,...,o_n).

FIRST-ORDER LOGIC EQUIVALENTS

$OWL: x \in C$	DL Syntax	FOL
C	C	C(x)
$intersectionOf(C_1,C_2)$	$C_1 \sqcap \ldots \sqcap C_n$	$C_1(x) \wedge \ldots \wedge C_n(x)$
$unionOf(C_1,C_2)$	$C_1 \sqcup \ldots \sqcup C_n$	$C_1(x) \vee \ldots \vee C_n(x)$
$complementOf(C_1)$	$\neg C_1$	$\neg C_1(x)$
$oneOf(x_1,\ldots,x_n)$	$\{x_1\} \sqcup \ldots \sqcup \{x_n\}$	$x = x_1 \vee \ldots \vee x = x_n$

$OWL: x \in C, Restriction \; on \; P$	DL Syntax	FOL
someValuesFrom(C')	$\exists P.C'$	$\exists y: P(x,y) \land C'(y)$
allValuesFrom(C')	$\forall P.C'$	$\forall y: P(x,y) \to C'(y)$
hasValue(y)	$\exists P.\{y\}$	P(x,y)
maxCardinality(n)	$\leq n.P$	$\exists^{\leq n} y : P(x,y)$
${\sf minCardinality}(n)$	$\geq n.P$	$\exists^{\geq n} y: P(x,y)$
cardinality(n)	n.P	$\exists^{=n} y : P(x,y)$

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FIRST-ORDER LOGIC EQUIVALENTS (CONT'D)

OWL Class Axioms for ${\cal C}$	DL Syntax	FOL
$rdfs.subClassOf(C_1)$	$C \sqsubseteq C_1$	$\forall x: C(x) \to C_1(x)$
$equivalentClass(C_1)$	$C \equiv C_1$	$\forall x: C(x) \leftrightarrow C_1(x)$
${\sf disjointWith}(C_1)$	$C \sqsubseteq \neg C_1$	$\forall x: C(x) \to \neg C_1(x)$
OWL Individual Axioms	DL Syntax	FOL
OWL Individual Axioms x_1 sameAs x_2	DL Syntax $\{x_1\} \equiv \{x_2\}$	FOL $x_1 = x_2$

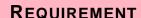
FIRST-ORDER LOGIC EQUIVALENTS (CONT'D)

OWL Properties	DL Syntax	FOL
\overline{P}	P	P(x,y)
OWL Property Axioms for P	DL Syntax	FOL
rdfs:range(C)	$\top \sqsubseteq \forall P.C$	$\forall x, y : P(x, y) \to C(y)$
$rdfs \mathpunct{:} domain(C)$	$C \sqsupseteq \exists P. \top$	$\forall x, y : P(x, y) \to C(x)$
$subPropertyOf(P_2)$	$P \sqsubseteq P_2$	$\forall x, y : P(x, y) \to P_2(x, y)$
${\sf equivalentProperty}(P_2)$	$P \equiv P_2$	$\forall x, y : P(x, y) \leftrightarrow P_2(x, y)$
$inverseOf(P_2)$	$P \equiv P_2^-$	$\forall x, y : P(x, y) \leftrightarrow P_2(y, x)$
TransitiveProperty	$P^+ \equiv P$	$\forall x,y,z: ((P(x,y) \land P(y,z)) \rightarrow P(x,z))$
		$\forall x, z: ((\exists y: P(x,y) \land P(y,z)) \rightarrow P(x,z))$
FunctionalProperty	$\top \sqsubseteq \leq 1P. \top$	$\forall x, y_1, y_2 : P(x, y_1) \land P(x, y_2) \to y_1 = y_2$
InverseFunctionalProperty	$\top \sqsubseteq \leq 1P^\top$	$\forall x, y_1, y_2 : P(y_1, x) \land P(y_2, x) \to y_1 = y_2$

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SYNTACTICAL REPRESENTATION

- OWL specifications can be represented by graphs: OWL constructs have a straightforward representation as triples in RDF/XML and N3.
- there are several logic-based representations (e.g. *Manchester OWL Syntax*); TERP (which can be used with pellet) is a combination of Turtle and Manchester syntax.
- OWL in RDF/XML format: usage of class, property, and individual names:
 - as @rdf:about when used as identifier of a subject (owl:Class, rdf:Property and their subclasses),
 - as @rdf:resource as the object of a property.
- some constructs need auxiliary structures (collections):
 owl:unionOf, owl:intersectionOf, and owl:oneOf are based on Collections
 - representation in RDF/XML by rdf:parseType="Collection".
 - representation in N3 by $(x_1 x_2 ... x_n)$
 - as RDF lists: rdf:List, rdf:first, rdf:rest



every entity in an OWL ontology must be explicitly typed (i.e., as a class, an object property, a datatype property, . . . , or an instance of some class).
 (for reasons of space this is not always done in the examples; in general, it may lead to incomplete results)

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QUERYING OWL DATA

- queries are atomic and conjunctive DL queries against the underlying OWL-DL model.
- this model can still be seen as a graph:
 - many of the edges are those known from the basic RDF graph
 - some edges (and collections) are only there for encoding OWL stuff (describing owl:unionOf, owl:propertyChain etc.) – these should not be queried
- SPARQL-DL is a subset of SPARQL: not every SPARQL query pattern is allowed for use on an OWL ontology (but the reasonable ones are, so in practice this is not a problem.)
- the query language SPARQL-DL allows exactly such well-sorted patterns using the notions of OWL.

SOME TBOX-ONLY REASONING EXAMPLES ON SETS

EXAMPLE: PARADOX

· without reasoner:

jena -t -if paradox.rdf

Outputs the same RDF facts in N3 without checking consistency.

with reasoner:

jena -e -pellet -if paradox.rdf

reads the RDF file, creates a model (and checks consistency) and in this case reports that it is not consistent.

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Union as $A \sqcup B \equiv \neg((\neg A) \sqcap (\neg B))$

```
Oprefix : <foo://bla/>.
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
:A rdf:type owl:Class.
                             :B rdf:type owl:Class.
:Union1 owl:unionOf (:A :B).
:CompA owl:complementOf :A. :CompB owl:complementOf :B.
:IntersectComps owl:intersectionOf (:CompA :CompB).
:Union2 owl:complementOf :IntersectComps.
:x rdf:type :A.
                              :x rdf:type :B.
:y rdf:type :CompA. # a negative assertion y not in A would be better -> OWL 2
                                                      [Filename: RDF/union.n3]
:y rdf:type :CompB.
prefix owl: <http://www.w3.org/2002/07/owl#>
prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>
prefix : <foo://bla/>
select ?X ?C ?D
                                                      [Filename: RDF/union.sparql]
from <file:union.n3>
where {{?X rdf:type ?C} UNION {:Union1 owl:equivalentClass ?D}}
```

EXAMPLE: UNION AND SUBCLASS

```
<?xml version="1.0"?>
<rdf:RDF xmlns:owl="http://www.w3.org/2002/07/owl#"
    xmlns:rdf="http://www.w3.org/1999/02/22-rdf-syntax-ns#"
    xmlns:f="foo://bla/"
    xml:base="foo://bla/">
 <owl:Class rdf:about="Person">
  <owl:unionOf rdf:parseType="Collection">
  <owl:Class rdf:about="Male"/>
  <owl:Class rdf:about="Female"/>
  </owl:unionOf>
 </owl:Class>
 <owl:Class rdf:about="EqToPerson">
  <owl:unionOf rdf:parseType="Collection">
  <owl:Class rdf:about="Female"/>
  <owl:Class rdf:about="Male"/>
  </owl:unionOf>
 </owl:Class>
 <f:Person rdf:about="unknownPerson"/>
                                                [Filename: RDF/union-subclass.rdf]
</rdf:RDF>
```

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Example (Cont'd)

print class tree (with jena -e -pellet):

```
owl:Thing
bla:Person = bla:EqToPerson - (bla:unknownPerson)
bla:Female
bla:Male
```

- Male and Female are derived to be subclasses of Person.
- Person and EqToPerson are equivalent classes.
- unknownPerson is a member of Person and EqToPerson.

EXERCISE

Consider

```
<owl:Class rdf:about="C1">
  <owl:intersectionOf rdf:parseType="Collection">
    <owl:Class rdf:about="A"/>
    <owl:Class rdf:about="B"/>
    </owl:intersectionOf>
  </owl:Class>
and
  <owl:Class rdf:about="C2">
    <rdfs:subClassOf rdf:resource="A"/>
    <rdfs:subClassOf rdf:resource="B"/>
  </owl:Class>
```

- · give mathematical characterizations of both cases.
- discuss whether both fragments are equivalent or not.

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DISCUSSION

- Two classes are *equivalent* (wrt. the knowledge base) if they have the same interpretation in every *model* of the KB.
- C_1 is characterized to be the intersection of classes A and B.
- for C_2 , it is asserted that C_2 is a subset of A and that it is a subset of B.
- Thus there can be some c that is in A, B, C_1 , but not in C_2 .
- Thus, C_1 and C_2 are not equivalent.

DISCUSSION: FORMAL NOTATION

The DL equivalent to the knowledge base (TBox) is

$$\mathcal{T} = \{ C_1 \equiv (A \sqcap B) , \quad C_2 \sqsubseteq A , \quad C_2 \sqsubseteq B \}$$

The First-Order Logic equivalent is

$$\mathcal{KB} = \{ \forall x : A(x) \land B(x) \leftrightarrow C_1(x) , \forall x : C_2(x) \rightarrow A(x) \land B(x) \}$$

Thus, $\mathcal{KB} \models \forall x : C_2(x) \rightarrow A(x) \land B(x)$.

Or, in DL: $\mathcal{T} \models C_2 \sqsubseteq C_1$.

On the other hand, $\mathcal{M} = (\mathcal{D}, \mathcal{I})$ with $\mathcal{D} = \{c\}$ and

$$\mathcal{I}(A) = \{c\}, \ \mathcal{I}(B) = \{c\}, \ \mathcal{I}(C_1) = \{c\}, \ \mathcal{I}(C_2) = \emptyset$$

is a model of KB (wrt. first-order logic) and T (wrt. DL) that shows that C_1 and C_2 are not equivalent.

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SUBCLASSES OF PROPERTIES

Triple syntax: some property rdf:type a specific type of property

According to their ranges

- owl:ObjectProperty subclass of rdf:Property; object-valued (i.e. rdfs:range must be an Object class)
- owl:DatatypeProperty subclass of rdf:Property; datatype-valued (i.e. its rdfs:range must be an rdfs:Datatype)
- ⇒ OWL ontologies require each property to be typed in such a way! (for reasons of space sometimes omitted in examples)

According to their Cardinality

- specifying n:1 or 1:n cardinality: owl:FunctionalProperty, owl:InverseFunctionalProperty
- ⇒ useful for deriving that objects must be different from each other.

According to their Properties

owl:TransitiveProperty, owl:SymmetricProperty see later ...

FUNCTIONAL CARDINALITY SPECIFICATION

property rdf:type owl:FunctionalProperty

- not a constraint, but
- if such a property results in two things ... these things are inferred to be the same.

```
@prefix : <foo://bla/names#>.
@prefix person: <foo://bla/persons/>.
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#> .
@prefix owl: <http://www.w3.org/2002/07/owl#>.
   :world :has_pope person:josephratzinger .
   :world :has_pope [ :name "Benedikt XVI" ] .
   :has_pope rdf:type owl:FunctionalProperty.
```

[Filename: RDF/popes.n3]

```
prefix : <foo://bla/names#>
prefix person: <foo://bla/persons/>
select ?N from <file:popes.n3>
where { person:josephratzinger :name ?N }
```

[Filename: RDF/pope.sparql]

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OWL: RESTRICTION - EXAMPLE

• owl:Restriction for $\exists p.C$ and $\forall p.C$. (cf. earlier examples)

```
<?xml version="1.0"?>
<rdf:RDF xmlns:owl="http://www.w3.org/2002/07/owl#"
   xmlns:rdf="http://www.w3.org/1999/02/22-rdf-syntax-ns#"
   xmlns:f="foo://bla/"
   xml:base="foo://bla/">
<owl:Class rdf:about="Parent">
  <owl:intersectionOf rdf:parseType="Collection">
  <owl:Class rdf:about="Person"/>
                                                      prefix : <foo://bla/>
  <owl:Restriction>
                                                      select ?C
   <owl:onProperty rdf:resource="hasChild"/>
                                                      from <file:restriction.rdf>
   <owl:minCardinality>1</owl:minCardinality>
                                                      where {: john a ?C}
  </owl:Restriction>
                                                     [Filename: RDF/restriction.sparql]
 </owl:intersectionOf>
 </owl:Class>
 <f:Person rdf:about="john">
  <f:hasChild><f:Person rdf:about="alice"/></f:hasChild>
 </f:Person>
                                                       [Filename: RDF/restriction.rdf]
</rdf:RDF>
```

RESTRICTIONS ONLY AS BLANK NODES

Consider the following (bad) specification:

```
:badIdea a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1.
```

This is not allowed in OWL-DL.

Correct specification:

```
:badIdea owl:equivalentClass
[a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1].
```

Why? ... there are many reasons, for one of them see next slide.

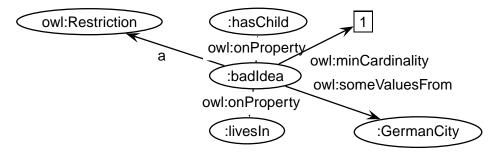
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Restrictions Only as Blank Nodes (Cont'd)

A class with two such specifications:

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/>.
:badIdea a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1 .
:badIdea a owl:Restriction; owl:onProperty :livesIn; owl:someValuesFrom :GermanCity.
[Filename: RDF/badIdea.n3]
```

• call jena -t -pellet -if badldea.n3:



The two restriction specifications are messed up.

Restrictions Only as Blank Nodes (Cont'd)

Thus specify each Restriction specification with a separate blank node:

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
Oprefix : <foo://bla/>.
 :TwoRestrictions owl:intersectionOf
  ( [ a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1]
    [ a owl:Restriction; owl:onProperty :livesIn; owl:someValuesFrom :GermanCity] ) .
[Filename: RDF/twoRestrictions.n3]
```

The DL equivalent: TwoRestrictions \equiv (\exists hasChild. \top) \sqcap (\exists livesIn.GermanCity)

Another reason:

```
:AnotherBadDesignExample a owl:Restriction;
  owl:onProperty :hasChild; owl:minCardinality 1;
  rdfs:subClassOf :Person.
```

... mixes the *definition* of the Restriction with an assertive axiom: ABDE $\equiv \exists \geq 1$ hasChild. $\top \land$ ABDE

□ Person

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MULTIPLE RESTRICTIONS ON A PROPERTY

"All persons that have at least two children, and one of them is male"

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix : <foo://bla/>.
### Test: multiple restrictions: the cardinality condition is then ignored
:HasTwoChildrenOneMale owl:intersectionOf (:Person
                                                         prefix : <foo://bla/>
  [ a owl:Restriction; owl:onProperty :hasChild;
                                                         from <file:restriction-double.n3>
    owl:someValuesFrom :Male; owl:minCardinality 2]).
                                                         where {?X a :HasTwoChildrenOneMale}
:name a owl:FunctionalProperty.
                                                         [Filename: RDF/restriction-double.sparql]
:Male rdfs:subClassOf :Person; owl:disjointWith :Female.
:Female rdfs:subClassOf :Person.
:kate a :Female; :name "Kate"; :hasChild :john.
:john a :Male; :name "John";
  :hasChild [a :Female; :name "Alice"], [a :Male; :name "Bob"].
:sue a :Female; :name "Sue";
                                                        [Filename: RDF/restriction-double.n3]
  :hasChild [a :Female; :name "Anne"], [a :Female; :name "Barbara"].
```

- The cardinality condition is ignored in this case (Result: John and Sue).
- Solution: intersection of restrictions

MULTIPLE RESTRICTIONS ON A PROPERTY

• "All persons that have at least two children, and one of them is male"

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
                                                         prefix : <foo://bla/>
                                                         select ?X
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
                                                          from <file:intersect-restrictions.n3>
Oprefix : <foo://bla/>.
                                                          where {?X a :HasTwoChildrenOneMale}
:HasTwoChildrenOneMale owl:intersectionOf (:Person
                                                         [Filename: RDF/intersect-restrictions.sparql]
  [ a owl:Restriction; owl:onProperty :hasChild; owl:someValuesFrom :Male]
  [ a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 2]).
:name a owl:FunctionalProperty.
:Male rdfs:subClassOf :Person; owl:disjointWith :Female.
:Female rdfs:subClassOf :Person.
:kate a :Female; :name "Kate"; :hasChild :john.
:john a :Male; :name "John";
  :hasChild [a :Female; :name "Alice"], [a :Male; :name "Bob"].
:sue a :Female; :name "Sue";
  :hasChild [a :Female; :name "Anne"], [a :Female; :name "Barbara"].
```

[Filename: RDF/intersect-restrictions.n3]

• Note: this is different from Qualified Range Restrictions such as "All persons that have at least two male children" – see Slide 380.

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USE OF A DERIVED CLASS

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix : <foo://bla/names#>.
:kate :name "Kate"; :child :john.
:john :name "John"; :child :alice.
:alice :name "Alice".
:Parent a owl:Class; owl:equivalentClass
  [a owl:Restriction; owl:onProperty :child; owl:minCardinality 1].
:Grandparent owl:equivalentClass
  [a owl:Restriction; owl:onProperty :child; owl:someValuesFrom :Parent].
```

[Filename: RDF/grandparent.n3]

[Filename: RDF/grandparent.sparql]

Non-Existence of a Property Fillers

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/names#>.
                                                       prefix : <foo://bla/names#>
:kate a :Person; :name "Kate"; :hasChild :john.
                                                        select ?X ?Y
:john a :Person; :name "John"; :hasChild :alice, :bob.
                                                       from <file:childless.n3>
:alice a :Person; :name "Alice".
                                                        where {{?X a :ChildlessA}
:bob a :Person; :name "Bob".
                                                               union {?Y a :ParentA}}
:name a owl:FunctionalProperty.
                                                       [Filename: RDF/childless.sparql]
:ChildlessA owl:intersectionOf (:Person
  [ a owl:Restriction; owl:onProperty :hasChild; owl:maxCardinality 0]).
:ChildlessB owl:intersectionOf (:Person
  [ a owl:Restriction; owl:onProperty :hasChild; owl:allValuesFrom owl:Nothing]).
:ParentA owl:intersectionOf (:Person [owl:complementOf :ChildlessA]).
:ParentB owl:intersectionOf (:Person
  [ a owl:Restriction; owl:onProperty :hasChild; owl:minCardinality 1]).
```

[Filename: RDF/childless.n3]

- export class tree: ChildlessA and ChildlessB are equivalent,
- note: due to the Open World Assumption, both classes are empty.
- Persons where no children are known are neither in ChildlessA or in Parent!

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INVERSE PROPERTIES

- owl:ObjectProperty owl:inverseOf owl:ObjectProperty
- owl:DatatypeProperties cannot have an inverse (this would define properties of objects, cf. next slide)

[Filename: RDF/inverse.n3] [Filename: RDF/inverse.sparql]

No Inverses of owl:DatatypeProperties!

an owl:DatatypeProperty must not have an inverse:

and foo://bla/names#age (DatatypeProperty)

• ":john :age 35" would imply "35 :ageOf :john" which would mean that a literal has a property, which is not allowed.

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#> .
@prefix : <foo://bla/names#> .
# :john :name "John"; :age 35; :child [:name "Alice"], [:name "Bob"; :age 8].
:age a owl:DatatypeProperty.
:child a rdf:Property.
:child of owl:inverseOf :child.
:ageOf owl:inverseOf :age.

[Filename: RDF/inverseDTProp.n3]

jena -e -pellet -if inverseDTProp.n3

WARN [main] (OWLLoader.java:352) - Unsupported axiom:
Ignoring inverseOf axiom between foo://bla/names#ageOf (ObjectProperty)
```

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Specification of Inverse Functional Properties

- Mathematics: a mapping m is inverse-functional if the inverse of m is functional: x p y is inverse-functional, if for every y, there is at most one x such that xpy holds.
- Example:
 - hasCarCode is functional: every country has one car code,
 - hasCarCode is also inverse functional: every car code uniquely identifies a country.
- OWL:

```
:m-inverse owl:inverseOf :m .
:m-inverse a owl:FunctionalProperty .
not allowed for e.g. mon:carCode a owl:DatatypeProperty:
```

```
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo:bla#>.
:carCode a owl:DatatypeProperty; rdfs:domain :Country;
  owl:inverseOf :isCarCodeOf.
# :Germany :carCode "D". [Filename: RDF/noinverse.n3]
```

the statement is rejected.

OWL: INVERSE FUNCTIONAL PROPERTY

- such cases are described with owl:InverseFunctionalProperty
- \bullet a property P is an owl:InverseFunctionalProperty if

```
\forall x, y_1, y_2 : P(y_1, x) \land P(y_2, x) \to y_1 = y_2 \text{ holds}
```

```
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo:bla#>.
:carCode rdfs:domain :Country; a owl:DatatypeProperty;
    a owl:FunctionalProperty; a owl:InverseFunctionalProperty.
:name a owl:DatatypeProperty; a owl:FunctionalProperty.
:Germany :carCode "D"; :name "Germany".
:DominicanRepublic :carCode "D"; :name "Dominican Republic".
```

[Filename: RDF/invfunctional.n3]

the fragment is detected to be inconsistent.

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OWL: HASKEY (OWL 2)

• description of key attributes (k_1, \ldots, k_n) is a relevant issue in data modeling. Note that InverseFunctionalProperty covers the simple case that n = 1.

```
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo:bla#>.
:name a owl:DatatypeProperty; a owl:FunctionalProperty.
:Country owl:hasKey (:carCode).
:Germany a :Country; :carCode "D"; :name "Germany".
:DominicanRepublic a :Country; :carCode "D"; :name "Dominican Republic".
:Duesseldorf a :City; :carCode "D"; :name "Duesseldorf".
```

[Filename: RDF/haskey.n3]

the fragment is inconsistent.

OWL:HASKEY (OWL 2)

 keys can also be used to detect that two resources (e.g. described by different Web sources) are actually the same:

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo:bla#>.
:Person owl:hasKey (:givenName :familyName).
   _:b1 a :Person; :firstName "John"; :familyName "Doe"; :age 32 .
   _:b2 a :Person; :firstName "John"; :familyName "Doe"; :address "Main Street 1" .
```

[Filename: RDF/haskey2.n3]

```
prefix : <foo:bla#>
select ?X ?P ?Y
from <file:haskey2.n3>
where {?X a :Person ; ?P ?Y}
```

[Filename: RDF/haskey2.sparql]

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TBox vs. ABox

DL makes a clean separation between TBox and ABox vocabulary:

- TBox: RDFS/OWL vocabulary for information about classes and properties (further partitioned into definitions and axioms),
- ABox: Domain vocabulary and rdf:type.

RDFS/OWL allows to mix everything in a set of triples.

Nominals

- use individuals (that usually occur only in the ABox) in the TBox:
- as individuals: Italy (that are often implemented in the reasoner as unary classes) with property owl:hasValue o (the class of all things such that {?x property o} holds).
- in enumerated classes class owl:oneOf (o₁,...,o_n)
 (class is defined to be the set {o₁,...,o_n}).

Difference to Reification

- Reification treats a class (e.g. :Penguin) or a property as an individual (:Penguin a :Species)
 - without reification, only specific RDFS and OWL properties are allowed for classes and properties only
 - reification assigns properties from an application domain to classes and properties.
- useful when talking about metadata notions,
- risk: allows for paradoxes

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USING NOMINALS: ITALIAN CITIES

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix mon: <http://www.semwebtech.org/mondial/10/meta#>.
@prefix it: <foo://italian/>.
it:Italy owl:sameAs <a href="http://www.semwebtech.org/mondial/10/countries/I/">http://www.semwebtech.org/mondial/10/countries/I/>.
it:ItalianProvince owl:intersectionOf
  (mon:Province
   [a owl:Restriction; owl:onProperty mon:isProvinceOf;
                                  # Nominal: an individual in a TBox axiom
    owl:hasValue it:Italy]).
it:ItalianCity owl:intersectionOf
  (mon:City
   [a owl:Restriction;
    owl:onProperty mon:belongsTo;
                                                           [Filename: RDF/italiancities.n3]
    owl:someValuesFrom it:ItalianProvince]).
prefix it: <foo://italian/>
select ?X
from <file:mondial-meta.n3>
from <file:mondial-europe.n3>
from <file:italiancities.n3>
where {?X a it:ItalianCity}
                                                           [Filename: RDF/italiancities.sparql]
```

AN ONTOLOGY IN OWL

Consider the Italian-English-Ontology from Slide 111.

```
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#> Class tree with jena -e:
@prefix owl: <http://www.w3.org/2002/07/owl#>.
                                                         owl:Thing
Oprefix f: <foo://bla/>.
                                                            bla:Person
f:Italian rdfs:subClassOf f:Person;
                                                               bla:English
  owl:disjointWith f:English;
                                                                  bla:Hooligan
  owl:unionOf (f:Lazy f:LatinLover).
                                                                   bla:Gentleman
f:Lazy owl:disjointWith f:LatinLover.
                                                               bla:Italian = bla:Lazy
f:English rdfs:subClassOf f:Person.
                                                         owl:Nothing = bla:LatinLover
f:Gentleman rdfs:subClassOf f:English.
f:Hooligan rdfs:subClassOf f:English.

    LatinLover is empty,

f:LatinLover rdfs:subClassOf f:Gentleman.
                                                            thus Italian \equiv Lazy.
```

[Filename: RDF/italian-english.n3]

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Italians and Englishmen (Cont'd)

the conclusions apply to the instance level:

```
@prefix : <foo://bla/>.
:mario a :Italian.
```

[Filename: RDF/mario.n3]

```
prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>
prefix : <foo://bla/>
select ?C
from <file:italian-english.n3>
from <file:mario.n3>
where {:mario rdf:type ?C}
[Filename: RDF/italian-english.sparql]
```

AN ONTOLOGY IN OWL

Consider the Italian-Ontology from Slide 112.

```
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix it: <foo://italian/>.
it:Bolzano owl:sameAs
<a href="http://www.semwebtech.org/mondial/10/countries/I/provinces/TrentinoAltoAdige/cities/Bdlzano/">http://www.semwebtech.org/mondial/10/countries/I/provinces/TrentinoAltoAdige/cities/Bdlzano/</a>
it:Italian owl:intersectionOf
  (it:Person
                                                              prefix : <foo://italian/>
   [a owl:Restriction; owl:onProperty it:livesIn;
                                                              select ?C
    owl:someValuesFrom it:ItalianCity]);
                                                              from <file:italian-prof.n3>
  owl:unionOf (it:Lazy it:Mafioso it:LatinLover).
                                                              from <file:mondial-meta.n3>
it:Professor rdfs:subClassOf it:Person.
                                                              from <file:mondial-europe.n3>
it:Lazy owl:disjointWith it:ItalianProf;
                                                              from <file:italiancities.n3>
   owl:disjointWith it:Mafioso;
                                                              where {:enrico a ?C}
                                                             [Filename: RDF/italian-prof.sparql]
   owl:disjointWith it:LatinLover.
it:Mafioso owl:disjointWith it:ItalianProf;
   owl:disjointWith it:LatinLover.
it:ItalianProf owl:intersectionOf (it:Italian it:Professor).
                                                            [Filename: RDF/italian-prof.n3]
it:enrico a it:Professor; it:livesIn it:Bolzano.
```

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ENUMERATED CLASSES: ONEOF

```
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix mon: <http://www.semwebtech.org/mondial/10/meta#>.
<bla:MontanunionMembers> owl:intersectionOf
                                                         select ?X
 (mon:Country
                                                         from <file:montanunion.n3>
  [owl:oneOf
                                                         from <file:mondial-europe.n3>
   (<http://www.semwebtech.org/mondial/10/countries/NL/
                                                         from <file:mondial-meta.n3>
   <http://www.semwebtech.org/mondial/10/countries/B/>
                                                         where {?X a <bla:Result>}
   <http://www.semwebtech.org/mondial/10/countries/L/>
                                                              [RDF/montanunion.sparql]
   <http://www.semwebtech.org/mondial/10/countries/F/>
   <http://www.semwebtech.org/mondial/10/countries/I/>
   <http://www.semwebtech.org/mondial/10/countries/D/>)]).
<bla:Result> owl:intersectionOf (mon:Organization
  [a owl:Restriction; owl:onProperty mon:hasMember;
                                                      [Filename: RDF/montanunion.n3]
  owl:someValuesFrom <bla:MontanunionMembers>]).
```

Query: all organizations that share a member with the Montanunion.

oneOf (Example Cont'd)

- previous example: "all organizations that share a member with the Montanunion." (DL: $x \in \exists$ hasMember.MontanunionMembers)
- "all organizations where *all* members are also members of the Montanunion." (DL: $x \in \forall$ hasMember.MontanunionMembers)
- The result is empty (although there is e.g. BeNeLux) due to open world: it is not known whether there may exist additional members of e.g. BeNeLux.
- Only if the membership of Benelux is "closed", results can be proven:

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix mon: <http://www.semwebtech.org/mondial/10/meta#>.
<http://www.semwebtech.org/mondial/10/organizations/Benelux/>
 a [a owl:Restriction;
                                                          select ?X
     owl:onProperty mon:hasMember; owl:cardinality 3].
                                                          from <file:montanunion.n3>
<bla:SubsetOfMU> owl:intersectionOf (mon:Organization)
                                                          from <file:montanunion2.n3>
  [a owl:Restriction; owl:onProperty mon:hasMember;
                                                          from <file:mondial-europe.n3>
     owl:allValuesFrom <bla:MontanunionMembers>]).
                                                          from <file:mondial-meta.n3>
mon:name a owl:FunctionalProperty. # not yet given in th where {?X a <bla:SubsetOfMU>}
                                                               [RDF/montanunion2.sparql]
[Filename: RDF/montanunion2.n3]
```

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oneOf (Example Cont'd)

"all organizations that cover all members of the Montanunion."

```
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix mon: <http://www.semwebtech.org/mondial/10/meta#>.
<bla:EUMembers> owl:equivalentClass [a owl:Restriction;
   owl:onProperty mon:isMember; owl:hasValue
   <http://www.semwebtech.org/mondial/10/organizations/EU/>].
```

[Filename: RDF/montanunion3.n3]

ONEOF (EXAMPLE CONT'D)

Previous example:

- only for one organization
- defined a class that contains all members of the organization
- not possible to define a family of classes one class for each organization.
- this would require a parameterized constructor:

```
"c_{org} is the set of all members of org"
```

Second-Order Logic: each organization can be seen as a unary predicate (=set):

```
\forall Org: Org(c) \leftrightarrow \mathsf{hasMember}(Org,c) or in F-Logic syntax: C isa Org :- Org:organization[hasMember->C] yields e.g.
```

```
I(eu) = \{germany, france, ...\},\ I(nato) = \{usa, canada, germany, ...\}
```

Recall that "organization" itself is a predicate:

```
I(organization) = \{eu, nato, \dots, \}
```

So we have again reification: organizations are both first-order-individuals and classes.

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CONVENIENCE CONSTRUCT: OWL: ALL DIFFERENT

- owl:oneOf defines a class as a closed set;
- in owl:oneOf (x_1, \ldots, x_n) , two items may be the same (open world),

owl:AllDifferent

Triples of the form :a owl:differentFrom :b state that two individuals are different.
 For a database with n elements, one needs

$$(n-1) + (n-2) + \ldots + 2 + 1 = \sum_{i=1..n} i = n \cdot (n+1)/2 = O(n^2)$$
 such statements.

The –purely syntactical– convenience construct

```
[ a owl:AllDifferent; owl:members (r_1 r_2 \dots r_n) ]
```

provides a shorthand notation.

- it is *immediately* translated into the set of all statements $\{r_i \text{ owl:differentFrom } r_j \mid i \neq j \in 1..n\}$
- [a owl:AllDifferent; owl:members (...)]
 is to be understood as a (blank node) that acts as a specification that the listed things are different that does not actually exist in the model.

[SYNTAX] OWL: ALL DIFFERENT IN RDF/XML

```
<?xml version="1.0"?>
<rdf:RDF xmlns:owl="http://www.w3.org/2002/07/owl#"
    xmlns:rdf="http://www.w3.org/1999/02/22-rdf-syntax-ns#"
    xmlns:f="foo://bla/" xml:base="foo://bla/">
                                                                   prefix : <foo://bla/>
                                                                   prefix owl:
 <owl:Class rdf:about="Foo">
                                                                    <http://www.w3.org/2002/07/owl#>
  <owl:equivalentClass> <owl:Class>
                                                                   select ?X ?P ?P2 ?V
    <owl:oneOf rdf:parseType="Collection">
                                                                  from <file:alldiff.rdf>
     <owl:Thing rdf:about="a"/> <owl:Thing rdf:about="b"/>
                                                                   where {?X a owl:AllDifferent ;
     <owl:Thing rdf:about="c"/> <owl:Thing rdf:about="d"/>
                                                                         ?P [?P2 ?V]}
    </owl:oneOf>
   </owl:Class> </owl:equivalentClass>
                                                                  [Filename: RDF/alldiffxml.sparql]
 </owl:Class>
 <owl:AllDifferent> <!-- use like a class, but is only a shorthand -->
    <owl:members rdf:parseType="Collection">
     <owl:Thing rdf:about="a"/> <owl:Thing rdf:about="b"/>
     <owl:Thing rdf:about="c"/> <owl:Thing rdf:about="d"/>
    </owl:members>
 </owl:AllDifferent>
 <owl:Thing rdf:about="a"> <owl:sameAs rdf:resource="b"/> </owl:Thing>
</rdf:RDF>
```

[Filename: RDF/alldiff.rdf]

- AllDifferent is only intended as a kind of command to the application to add all pairwise "different-from" statements, it does not actually introduce itself as triples:
- querying {?X a owl:AllDifferent} is actually not intended.

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[SYNTAX] OWL: ALL DIFFERENT IN N3

Example:

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/>.
:Foo owl:equivalentClass [ owl:oneOf (:a :b :c :d) ].
# both the following syntaxes are equivalent and correct:
[ a owl:AllDifferent; owl:members (:a :b)].
[] a owl:AllDifferent; owl:members (:c :d).
:a owl:sameAs :b.
# :b owl:sameAs :d.
[Filename: RDF/alldiff.n3]
```

```
prefix : <foo://bla/>
select ?X ?Y
from <file:alldiff.n3>
where {?X a owl:AllDifferent ; ?P [?P2 ?V]} [Filename: RDF/alldiff.sparql]
```

ONEOF: A TEST

- owl:oneOf defines a "closed set" (use with anonymous class; see below):
- note that in owl:oneOf (x_1, \ldots, x_n) , two items may be the same (open world),
- optional owl:AllDifferent to guarantee that (x_1, \ldots, x_n) are pairwise distinct.

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/>.
:Person owl:equivalentClass [ owl:oneOf (:john :alice :bob) ].
# :john owl:sameAs :alice. # to show that it is consistent that they are the same
[] a owl:AllDifferent; owl:members (:john :alice :bob). # to guarantee distinctness
# :name a owl:FunctionalProperty. # this also guarantees distinctness;)
:john :name "John".
:alice :name "Alice".
:bob :name "Bob".
:d a :Person.
:d owl:differentFrom :john; owl:differentFrom :alice.
# :d owl:differentFrom :bob. ### adding this makes the ontology inconsistent
```

[Filename: RDF/three.n3]

• Who is :d?

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oneOf: a Test (cont'd)

Who is :d?

- check the class tree:
 bla:Person (bla:bob, bla:alice, bla:d, bla:john)
- and ask it:

```
prefix : <foo://bla/>
select ?N
from <file:three.n3>
where {:d :name ?N}
```

[Filename: RDF/three.sparql]
The answer is ?N/"Bob".

ANSWER SETS TO QUERIES AS AD-HOC CONCEPTS

all organizations whose headquarter city is a capital:

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <http://www.semwebtech.org/mondial/10/meta#> .
:CountryCapital owl:intersectionOf
  (:City [a owl:Restriction; owl:onProperty :isCapitalOf;
           owl:someValuesFrom :Country]).
<bla:Result> owl:intersectionOf
  (:Organization [a owl:Restriction; owl:onProperty:hasHeadq;
    owl:someValuesFrom :CountryCapital]).
                                                     [Filename: RDF/organizations-query.n3]
prefix : <http://www.semwebtech.org/mondial/10/meta#>
select ?A ?N
from <file:organizations-query.n3>
from <file:mondial-europe.n3>
from <file:mondial-meta.n3>
where \{?X \text{ a } < bla: Result> . ?X : abbrev ?A . ?X : hasHeadq ?C . ?C : name ?N\}
[Filename:RDF/organizations-query.sparql]
```

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COMPLEX TERMS IN SPARQL QUERIES

- example: all cities that are a capital
- use pellet! jena does not support this:

COMPLEX TERMS IN SPARQL QUERIES (CONT'D)

- all organizations whose headquarter city is a capital:
- use pellet! jena does not support this:

```
pellet query -query-file organizations-query3.sparql \
  mondial-europe.n3 mondial-meta.n3
```

[Filename:RDF/organizations-query3.sparql]

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NAMED AND UNNAMED RESOURCES

(from the DL reasoner's perspective)

Named Resources

- resources with explicit global URIs
 - ">http://www.semwebtech.org/mo
- resources with local IDs/named blank nodes
- unnamed blank nodes

Unnamed (implicit) Resources

things that exist only implicitly:
 John's child in

```
:Parent a owl:Class; owl:equivalentClass
   [ a owl:Restriction; owl:onProperty :child; owl:minCardinality 1].
:john a Parent.
```

• such resources can even have properties (see next slides).

Implicit Resources

- "every person has a father who is a person" and "john is a person".
- the standard model is infinite: john, john's father, john's father's father, ...
- pure RDF graphs are always finite,
- only with OWL axioms, one can specify such infinite models,
- \Rightarrow they have only finitely many *locally to path length* n different nodes,
 - the reasoner can detect the necessary *n* ("blocking", cf. Slides 434 ff) and create "typical" different structures.

Aside: "standard model" vs "nonstandard model"

- the term "standard model" is not only "what we understand (in this case)", but is a notion of mathematical theory which –roughly– means "the simplest model of a specification"
- nonstandard models of the above are those where there is a cycle in the ancestors relation.
 - (as the length of the cycle is arbitrary, this would not make it easier for the reasoner there is only the possibility to have an owl:sameAs somewhere)

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Implicit Resources

```
@prefix : <foo://bla/names#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
:Person owl:equivalentClass [a owl:Restriction;
   owl:onProperty :father; owl:someValuesFrom :Person].
:bob :name "Bob"; a :Person; :father :john.
:john :name "John"; a :Person.
```

[Filename: RDF/fathers-and-forefathers.n3]

```
prefix : <foo://bla/names#>
prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>
select ?X ?F ?C
from <file:fathers-and-forefathers.n3>
where {{ ?X :father ?F } UNION { ?X a :Person }}
```

[Filename: RDF/fathers-and-forefathers.sparql]

- Reasoner: works on the model, including blocking, i.e. modulo equivalence up to paths of length n.
- SPARQL (and SWRL) rules: works on the graph without the unnamed/implicit resorces.

SPARQL ON THE GRAPH

 SPARQL does not return any answer related with nodes (=resources) that are only implicitly known (=non-named resources)

```
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
:ParentOf12YOChild owl:equivalentClass [a owl:Restriction;
   owl:onProperty :child; owl:someValuesFrom :12YOPerson].
:12YOPerson owl:equivalentClass [a owl:Restriction;
   owl:onProperty :age; owl:hasValue 12].
[ :name "John"; :age 35; a :ParentOf12YOChild;
        :child [:name "Alice"; :age 10], [:name "Bob"; :age 8]].
:age rdf:type owl:FunctionalProperty.
# :12YOPerson owl:equivalentClass owl:Nothing.

:TwoChildrenParent owl:equivalentClass [a owl:Restriction;
   owl:onProperty :child; owl:cardinality 2].
:ThreeChildrenParent owl:equivalentClass [a owl:Restriction;
   owl:onProperty :child; owl:minCardinality 3]. [Filename: RDF/john-three-children-impl.n3]
```

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SPARQL and Non-Named Resources (Cont'd)

- implicit resources exist only on the reasoning level,
- not considered by SPARQL queries:

[Filename: RDF/john-three-children-impl.sparql]

- John is a ThreeChildrenParent,
- no person known who is 12 years old
- adding :12YOPerson owl:equivalentClass owl:Nothing makes it inconsistent.
- same applies to owl:hasKey (cf. Slide 341) and SWRL rules (cf. Slides 437 ff).

OWL: HASKEY AND NON-NAMED RESOURCES

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix swrl: <http://www.w3.org/2003/11/swrl#>.
@prefix : <foo:bla#>.
:XYThing owl:hasKey (:x :y).
:child rdfs:range :XYThing.
:xy10 a :XYThing; :x 10; :y 10; :text "free".
:TenThing rdfs:subClassOf [ a owl:Restriction;
   owl:onProperty :child; owl:onClass :XYTen; owl:qualifiedCardinality 1].
:XYTen owl:intersectionOf
  ([ a owl:Restriction; owl:onProperty :x; owl:hasValue 10]
   [ a owl:Restriction; owl:onProperty:y; owl:hasValue 10]
   [ a owl:Restriction; owl:onProperty :text; owl:hasValue "thechild"]).
:OneChildThing rdfs:subClassOf [ a owl:Restriction;
    owl:onProperty :child; owl:qualifiedCardinality 1].
:tenTen a :TenThing; a :OneChildThing. # forces the implicit node
# make this implicit node named by forcing "another" child
# of :tenTen without any properties: via hasKey, xyxy = xy10
# :tenTen :child :xyxy.
                                               [Filename: RDF/easykeys-impl.n3]
```

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OWL:HASKEY AND NON-NAMED RESOURCES (CONT'D)

[Filename: RDF/easykeys-impl.sparql]

 as long as the relevant node is only implicit (although quite some information about it is known), it is not considered in the answers.

CLOSING PARTS OF THE OPEN WORLD

- "forall items" is only applicable if additional items can be excluded (⇒ locally closed predicate/property),
- often, RDF data is generated from a database,
- certain predicates can be closed by defining restriction classes with maxCardinality.

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OWL:ALLVALUESFROM

```
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/names#>.
[ a :Male; a :ThreeChildrenParent; :name "John";
    :child [a :Female; :name "Alice"], [a :Male; :name "Bob"],
           [a :Female; :name "Carol"]].
[ a :Female; a :TwoChildrenParent; :name "Sue";
   :child [a :Female; :name "Anne";], [a :Female; :name "Barbara"]].
:name a owl:FunctionalProperty.
:OneChildParent owl:equivalentClass [a owl:Re prefix : <foo://bla/names#>
                                                select ?N
  owl:onProperty :child; owl:cardinality 1].
                                                from <file:allvaluesfrom.n3>
:TwoChildrenParent owl:equivalentClass [a owl
                                                where {?X :name ?N .
  owl:onProperty :child; owl:cardinality 2].
                                                  ?X a :OnlyFemaleChildrenParent}
:ThreeChildrenParent owl:equivalentClass [a o Filename: RDF/allvaluesfrom.sparql]
  owl:onProperty :child; owl:cardinality 3].
:OnlyFemaleChildrenParent owl:equivalentClass [a owl:Restriction;
  owl:onProperty :child; owl:allValuesFrom :Female].
[Filename: RDF/allvaluesfrom.n3]
```

EXAMPLE: WIN-MOVE-GAME IN OWL

```
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/>.
:Node a owl:Class; owl:equivalentClass
  [ a owl:Class; owl:oneOf (:a :b :c :d :e :f :g :h :i :j :k :l :m)].
:edge a owl:ObjectProperty; rdfs:domain :Node; rdfs:range :Node.
:out a owl:DatatypeProperty.
:a a :Node; :out 2; :edge :b, :f.
:b a :Node; :out 3; :edge :c, :g, :k.
:c a :Node; :out 2; :edge :d, :1.
:d a :Node; :out 1; :edge :e.
:e a :Node; :out 1; :edge :a.
:f a :Node; :out 0 .
:g a :Node; :out 2; :edge :i, :h.
:h a :Node; :out 1; :edge :m.
:i a :Node; :out 1; :edge :j.
:j a :Node; :out 0 .
:k a :Node; :out 0 .
:1 a :Node; :out 1; :edge :d.
                                                         [Filename: RDF/winmove-graph.n3]
:m a :Node; :out 1; :edge :h.
```

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Win-Move-Game in OWL – the Game Axioms

"If a player cannot move, he loses."

Which nodes are WinNodes, which one are LoseNodes (i.e., the player who has to move wins/loses)?

- if a player can move to some LoseNode (for the other), he will win.
- if a player can move only to WinNodes (for the other), he will lose.
- recall that there can be nodes that are neither WinNodes nor LoseNodes.

```
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/>.

:WinNode a owl:Class; owl:intersectionOf ( :Node
    [a owl:Restriction; owl:onProperty :edge; owl:someValuesFrom :LoseNode]).
:LoseNode a owl:Class; owl:intersectionOf ( :Node
    [a owl:Restriction; owl:onProperty :edge; owl:allValuesFrom :WinNode]).
```

[Filename: RDF/winmove-axioms.n3]

Win-Move-Game in OWL - Closure

```
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>.
@prefix owl: <http://www.w3.org/2002/07/owl#>.
@prefix : <foo://bla/>.
:DeadEndNode a owl:Class; rdfs:subClassOf :Node;
  owl:equivalentClass [ a owl:Restriction; owl:onProperty :out; owl:hasValue 0],
                      [ a owl:Restriction; owl:onProperty :edge; owl:cardinality 0].
:OneExitNode a owl:Class; rdfs:subClassOf :Node;
  owl:equivalentClass [ a owl:Restriction; owl:onProperty :out; owl:hasValue 1],
                      [ a owl:Restriction; owl:onProperty :edge; owl:cardinality 1].
:TwoExitsNode a owl:Class; rdfs:subClassOf :Node;
 owl:equivalentClass [ a owl:Restriction; owl:onProperty :out; owl:hasValue 2],
                      [ a owl:Restriction; owl:onProperty :edge; owl:cardinality 2].
:ThreeExitsNode a owl:Class; rdfs:subClassOf :Node;
  owl:equivalentClass [ a owl:Restriction; owl:onProperty :out; owl:hasValue 3],
                      [ a owl:Restriction; owl:onProperty :edge; owl:cardinality 3].
```

[Filename: RDF/winmove-closure.n3]

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Win-Move-Game in OWL: DeadEndNodes

Prove that DeadEndNodes are LoseNodes:

- obvious: Player cannot move from there
- exercise: give a formal (Tableau) proof
- The OWL Reasoner does it:

```
prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#>
prefix : <foo://bla/>
select ?X
from <file:winmove-axioms.n3>
from <file:winmove-closure.n3>
where {:DeadEndNode rdfs:subClassOf :LoseNode}
```

[Filename: RDF/deadendnodes.sparql]

The answer contains an (empty) tuple which means "yes".

Win-Move-Game in OWL

[Filename: RDF/winmove.sparql]

lose: f, k, j, e, l win: c, a, i, b, d

Exercise

• Is it possible to characterize DrawNodes in OWL?

note: code and proof see LaTeX source