

## 8.1 Bridge Section:

- The lecture “Database Theory” or “Deductive Databases” (MSc or advanced BSc) builds upon the “Introduction to Databases” lecture and requires knowledge about First-Order Logic (e.g., courses “Formal Systems” or “Artificial Intelligence”)
- This section summarizes that knowledge and motivates the main idea of the lecture.
- a database can be seen as a purely relational FOL structure
  - predicate symbols of different arities,
  - only 0-ary functions = constants
    - \* in relational DB: these are the literals (numbers, strings, dates ...)
    - \* in object-relational DB: also object identifiers
    - \* in RDF: also URIs, which basically serve as object identifiers

SOME CONTENTS STILL MISSING

## TYPES OF KNOWLEDGE

- (positive) atomic facts:
  - DB: tuples in an  $n$ -column table of the database
  - FOL:  $\mathcal{S} = (\mathcal{I}, \mathcal{D})$ : for an  $n$ -ary predicate,  $\mathcal{I}(p) \subset \mathcal{D}^n$
  - atoms in a formula

⇒ conjunctions/sets of atomic facts
- negative atomic facts/knowledge:
  - rather “implicit”: the  $n$ -tuples “not there” in a DB or not in  $\mathcal{I}(p)$ .

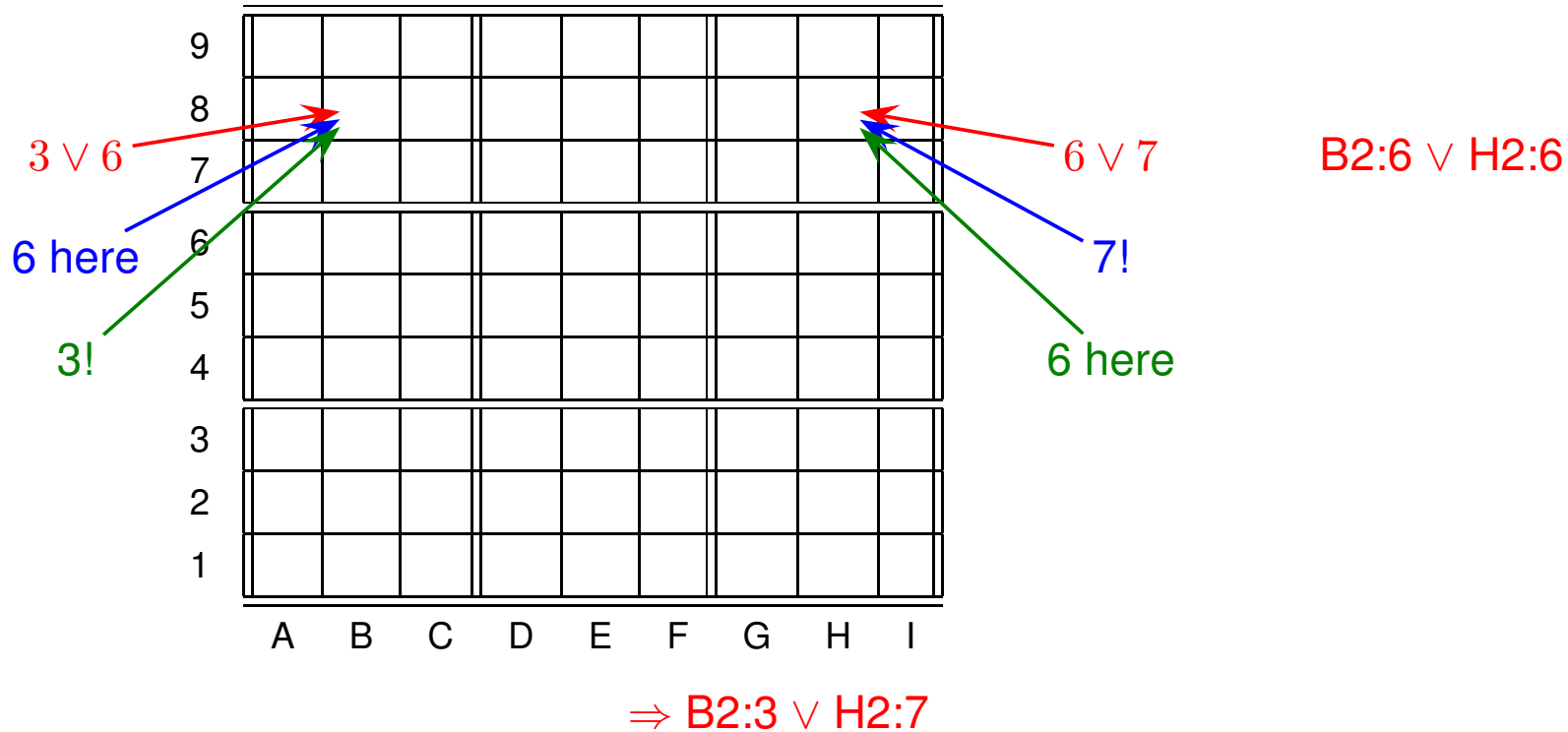
⇒ queries under CWA and  $\mathcal{S} \models \phi$ .
- atomic positive conclusions: INSERT into DB, Views
- atomic negative conclusions: DELETE, or inconsistencies

## Disjunctive Knowledge

- “ $p(x)$  or  $q(y)$  does hold”
- cannot be represented by a database or a single FOL interpretation, only by formulas

⇒ conclusions in “knowledge base”

## Disjunctive Knowledge in Human Reasoning: Sudoku



## Existential Knowledge

- “every country has some city that is its capital (and which is located in this country)”

$\text{country}(x) \rightarrow \exists y: ( \text{city}(y) \wedge \text{located\_in}(y, x) )$

- Mondial SQL: not null and foreign key to primary key reference:  
country.(code, capital, capprov) references city.(country, capital, province)

- “everything which is a parent has *some* child (which is a person)”

ER model: Parent is a subclass of Person, minCardinality of “hasChild” is 1

OWL/Semantic Web: Parent  $\equiv \exists \text{child}.\text{Person}$

$\Leftarrow$ : SQL: view

$\Rightarrow$ : SQL: not possible

FOL, e.g. tableau calculus: skolem function  $\text{hasChild}(\text{alice}, f_{\text{child}}(\text{alice}))$

- “every person has a parent (which is a person)”