

$win(a) :- move(a,b), \neg win(b).$   
 $win(a) :- move(a,f), \neg win(f).$   
 $win(b) :- move(b,c), \neg win(c).$   
 $win(b) :- \quad \quad \quad \neg win(g).$   
 $win(b) :- \quad \quad \quad \neg win(e).$   
 $win(c) :- \quad \quad \quad \neg win(d).$   
 $win(c) :- \quad \quad \quad \neg win(l).$   
 $win(d) :- \quad \quad \quad \neg win(e).$   
 $win(e) :- \quad \quad \quad \neg win(a).$   
 $win(f) :-$   
 $win(g) :-$   
 $win(g) :-$   
 $win(h) :-$   
 $win(i) :-$   
 $win(j) :-$   
 $win(k) :-$   
 $win(l) :-$   
 $win(m) :-$   
 $win(n) :-$

no move, no satisfiable ground instance of rule.

$\neg win(h)$   
 $\neg win(i)$   
 $\neg win(m)$   
 $\neg win(j)$   
 $\neg win(d)$   
 $\neg win(h)$

$win(a) :- move(a,b), \neg win(b).$   
 $win(a) :- move(a,f), \neg win(f).$   
 $win(b) :- move(b,c), \neg win(c).$   
 $win(b) :- \quad \quad \quad \neg win(g).$   
 $win(b) :- \quad \quad \quad \neg win(e).$   
 $win(c) :- \quad \quad \quad \neg win(d).$   
 $win(c) :- \quad \quad \quad \neg win(l).$   
 $win(d) :- \quad \quad \quad \neg win(e).$   
 $win(e) :- \quad \quad \quad \neg win(a).$   
 $win(f) :-$   
 $win(g) :-$   
 $win(g) :-$   
 $win(h) :-$   
 $win(i) :-$   
 $win(j) :-$   
 $win(k) :-$   
 $win(l) :-$   
 $win(m) :-$   
 $win(n) :-$

$\neg win:$   
 $P_0 = \{facts + \text{all } win(x) \text{ false}\}$   
 $T_0^1 = \{ \text{all "move" facts from the game} \}$   
 $T_0^2 = \{ \text{all "move"s} \}$   
 $T_0^3 = \dots$   
 $T_0^4 = \dots$   
 $T_0^5 = \dots$

$\neg win(h)$   
 $\neg win(i)$   
 $\neg win(m)$   
 $\neg win(j)$   
 $\neg win(d)$   
 $\neg win(h)$

no move, no satisfiable ground instance of rule

no grnd. instance  
no move  
finite, n have no chance to be in "win"  
=> avoid this

⇒ Define the reduct wrt.  $\mathcal{I}_0$  ← fix all negative literals

$wih(a) :- move(a,b), \neg win(b).$   
 $wih(a) :- move(a,f), \neg win(f).$   
 $wih(b) :- move(b,c), \neg win(c).$   
 $wih(b) :- " \neg win(g).$   
 $wih(b) :- " \neg win(e).$   
 $wih(c) :- " \neg win(d).$   
 $wih(c) :- " \neg win(f).$   
 $wih(d) :- " \neg win(e).$   
 $wih(e) :- " \neg win(a).$   
 $wih(f) :- " \neg win(a).$   
 $wih(g) :- " \neg win(f).$   
 $wih(g) :- " \neg win(i).$   
 $wih(h) :- " \neg win(m).$   
 $wih(i) :- " \neg win(j).$   
 $wih(j) :- " \neg win(i).$   
 $wih(k) :- " \neg win(d).$   
 $wih(l) :- " \neg win(h).$   
 $wih(m) :- " \neg win(h).$   
 $wih(n) :- " \neg win(h).$

$\mathcal{I}_0 = \emptyset \rightarrow \neg win(x)$  is true  
 $\mathcal{I}_0 = \emptyset$   
 $\mathcal{I}_1 = P_{\mathcal{I}_0}$   
 $\mathcal{I}_2 = \{all\ moves\}$   
 $\mathcal{I}_3 = \{all\ moves\} \cup \{win(\dots)\}$  as previous slide  
 $T_{P_{\mathcal{I}_0}}^3(p)$  fixes all  $\neg win$  to true( $\mathcal{I}_0$ )  
 $wih(T_2) = win(T_2)$   
 $wih(T_3) = win(T_3)$   
 $wih(T_4) = \dots$   
 $\mathcal{I}_1$   
 $\Rightarrow$  all  $x$  s.t.  $\exists$  move  $ac$  in "win"  
 $\Rightarrow$  overestimate of "win"

Reduct wrt.  $\mathcal{I}_n$ : ( $\mathcal{I}_n$ :  $wih = all\ except\ \{f,j,k,n\}$ )

$wih(a) :- move(a,b), \neg win(b) \text{ false}$   
 $wih(a) :- move(a,f), \neg win(f) \text{ true}$   
 $wih(b) :- move(b,c), \neg win(c) \text{ false}$   
 $wih(b) :- " \neg win(g) \text{ true}$   
 $wih(b) :- " \neg win(e) \text{ true}$   
 $wih(c) :- " \neg win(d) \text{ true}$   
 $wih(c) :- " \neg win(f) \text{ true}$   
 $wih(d) :- " \neg win(e) \text{ true}$   
 $wih(e) :- " \neg win(a) \text{ true}$   
 $wih(g) :- " \neg win(i) \text{ true}$   
 $wih(g) :- " \neg win(m) \text{ true}$   
 $wih(h) :- " \neg win(j) \text{ true}$   
 $wih(l) :- " \neg win(d) \text{ true}$   
 $wih(m) :- " \neg win(h) \text{ true}$

$\mathcal{I}_n$   
 $T_{P_{\mathcal{I}_n}}^1(\emptyset) = \{all\ moves\}$   
 $T_{P_{\mathcal{I}_n}}^2(\emptyset) = \{all\ moves\} \cup \{wih(a), wih(b), wih(i)\}$   
 $= T_{P_{\mathcal{I}_n}}^3(\emptyset) = \mathcal{I}_2$

definitely not in win

**Reduct wrt.  $\mathcal{I}_2$**   $\mathcal{I}_2$ : win (a, b, i)

<del>win(a) :- move(a, b), <math>\neg</math> win(b).</del>	<del><math>\neg</math> win(b) true =:</del>
win(a) :- move(a, f), $\neg$ win(f) true =:	$T_{P_{\mathcal{I}_2}}^1(\emptyset) = \{\text{all moves}\}$
<del>win(b) :- move(b, c), <math>\neg</math> win(c).</del>	$T_{P_{\mathcal{I}_2}}^2(\emptyset) = \{\text{all moves}\}$
<del>win(b) :- " <math>\neg</math> win(g).</del>	$\cup \{\text{win: a, b, c, d, g, h, i, l, m}\}$
<del>win(b) :- " <math>\neg</math> win(e).</del>	
<del>win(c) :- " <math>\neg</math> win(d).</del>	
<del>win(c) :- " <math>\neg</math> win(l).</del>	
<del>win(d) :- " <math>\neg</math> win(e).</del>	
<del>win(e) :- " <math>\neg</math> win(a).</del>	
<del>win(f) :- " no move, no satisfiable ground instance of rule.</del>	
win(g) :- " $\neg$ win(h)	
<del>win(g) :- " <math>\neg</math> win(i)</del>	
<del>win(h) :- " <math>\neg</math> win(m)</del>	
win(i) :- " $\neg$ win(j)	
<del>win(j) :- " "</del>	
<del>win(k) :- " "</del>	
win(l) :- " $\neg$ win(d)	
win(m) :- " $\neg$ win(h)	
<del>win(n) :- " "</del>	

$\Rightarrow$  not in win:  $\mathcal{I}_3$   
 $\{e, f, k, j, n\}$

**reduct wrt.  $\mathcal{I}_3$**  not win  $\mathcal{I}_3$ : e, f, k, j, n

<del>win(a) :- move(a, b), <math>\neg</math> win(b).</del>	$\rightarrow$ (win(..) is true)
win(a) :- move(a, f), $\neg$ win(f) true	$\rightarrow$ all others: $\neg$ win(.) is false
<del>win(b) :- move(b, c), <math>\neg</math> win(c).</del>	
<del>win(b) :- " <math>\neg</math> win(g).</del>	
<del>win(b) :- " <math>\neg</math> win(e).</del>	
<del>win(c) :- " <math>\neg</math> win(d).</del>	
<del>win(c) :- " <math>\neg</math> win(l).</del>	
<del>win(d) :- " <math>\neg</math> win(e) true</del>	
<del>win(e) :- " <math>\neg</math> win(a).</del>	
<del>win(f) :- " no move, no satisfiable ground instance of rule.</del>	
<del>win(g) :- " <math>\neg</math> win(h)</del>	
<del>win(g) :- " <math>\neg</math> win(i)</del>	
<del>win(h) :- " <math>\neg</math> win(m)</del>	
win(i) :- " $\neg$ win(j) true	
<del>win(j) :- " "</del>	
<del>win(k) :- " "</del>	
<del>win(l) :- " <math>\neg</math> win(d)</del>	
<del>win(m) :- " <math>\neg</math> win(h)</del>	
<del>win(n) :- " "</del>	

$\Rightarrow$  can't be exercise

$\mathcal{K}_0 = \emptyset \rightarrow$  underestimate of true atoms  
 overestimate of false atoms  
 $p(\dots) :- q(\dots), \neg r(\dots)$  *time*  
 $\Rightarrow$  reduct  $P_{\mathcal{K}_0}$  is a "big" program  
 $\Rightarrow$  denies relatively much  $\hat{=}$  everything "possible"  
 $\hat{=}$  overestimate of the true atoms  $=: \mathcal{K}_1$   
 (overestimate of "win")  
 $P_{\mathcal{K}_1} \Rightarrow$  many needed literals are "false"  
 $\Rightarrow$  many rules are deleted  
 "small" program  $\rightarrow$  small  $T_{P_{\mathcal{K}_1}}^w(\emptyset)$   
 $\rightarrow$  Underestimate of true atoms  
 "everything necessary"

