

Ex. 2.4, Slide 605 (borders, reachable)

- stratified negation

consider first naive T_P^ω strategy (cf. T_P with $\overline{T}_P(\emptyset) = \emptyset$ (cf. Definition 3.1.5))

$$\overline{T}_P(\emptyset) = \emptyset \quad (\text{cf. Definition 3.1.5})$$

$$\overline{T}_P(\emptyset) = \overline{T}_P(\emptyset) = EDB$$

$$\overline{T}_P^2(\emptyset) = \overline{T}_P(EDB) = [EDB \cup \{ \text{borders}(y, x) \mid \text{borders}(y, x) \in EDB \}]$$

$$\overline{T}_P^3(\emptyset) = [..] \cup \text{reachable}^1 \text{borders} \cup \{ \text{notReachable}(x, y) \mid x, y \text{ reached} \}$$

$$\vdots \cup \{ \text{notReachable}(x, y) \mid x, y \text{ reached} \}$$

$$\overline{T}_P^{h+2}(\emptyset) = EDB \cup \text{borders}^{\text{from}} \cup \text{reachable}^{\text{all}}$$

$$\overline{T}_P^{h+3}(\emptyset) = EDB \cup \text{borders}^{\text{from}} \cup \text{reachable}^{\text{all}}$$

$$\overline{T}_P^{h+4}(\emptyset) = \overline{T}_P^{h+3}$$

$\Rightarrow \lim_{h \rightarrow \infty} \overline{T}_P^h(\emptyset)$ is often (but always...) also the minicla model.

If there is a rule " $P \leftarrow P$ "

$\text{notReachable}(x, y) :- \text{notReachable}(x, y).$

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Consider stability of the stratified model Ψ_P wrt. transduction 2. P (Slide 605):

$$=: P$$

$$P_\Psi = \{ \begin{array}{l} \text{borders}(b, a, l) :- \text{borders}(a, b, l) \mid a, b \text{ can be orders}, \\ \quad l \in T[\overline{T}_P^h(\text{borders})] \\ \{ \text{reachable}(a, b) :- \text{borders}(a, b) \mid a, b \text{ can be orders} \} \\ \{ \text{reachable}(a, b) :- \text{reachable}(a, c), \text{borders}(c, b) \mid a, b, c \text{ can be orders} \} \\ \{ \text{notReachable}(a, b) :- \text{notReachable}(a, b), \text{borders}(a, b) \mid a, b \text{ can be orders} \} \end{array}$$

\Rightarrow all (relevant) ground instances of rules

P_Ψ , take P_Ψ , delete $\{ \text{notReachable}(a, b) :- \text{notReachable}(a, b), \text{borders}(a, b) \mid a, b \text{ can be orders} \}$

and if the other rules, delete the "notReachable(a, b)" (because $\text{reachable}(a, b) \notin \Psi_P$)

\Rightarrow amongst the blue rules,

only the instances for (a, b) that are not inReachable remain.

$$\Rightarrow \overline{T}_{P_\Psi}^\omega(\emptyset) = P_\Psi$$

is stable. \square

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Ex 2.4. "b" several minimal models

P border, readable, notReachable :

- \mathcal{Y}^P is a minimal model
- $\mathcal{Y}' := \mathcal{Y}^P \cup \{ \text{reachable } (IS, R) \}$
 $\cup \{ \text{reachable } (IS, X) \mid \text{readable } (R, X) \in \mathcal{Y}^P \}$

\mathcal{Y}' is a model of P

is not minimal

(contains atoms that are not (reachable) required by P)

- $\mathcal{Y}'' := \mathcal{Y}' \setminus \left(\{ \text{notReachable } (IS, X) \mid \text{reachable } (R, X) \in \mathcal{Y}^P \} \right)$
 $\cup \{ \text{notReachable } (IS, R) \}$
- is a model of P
- is minimal

is \mathcal{Y}'' stable wrt (original) P ?

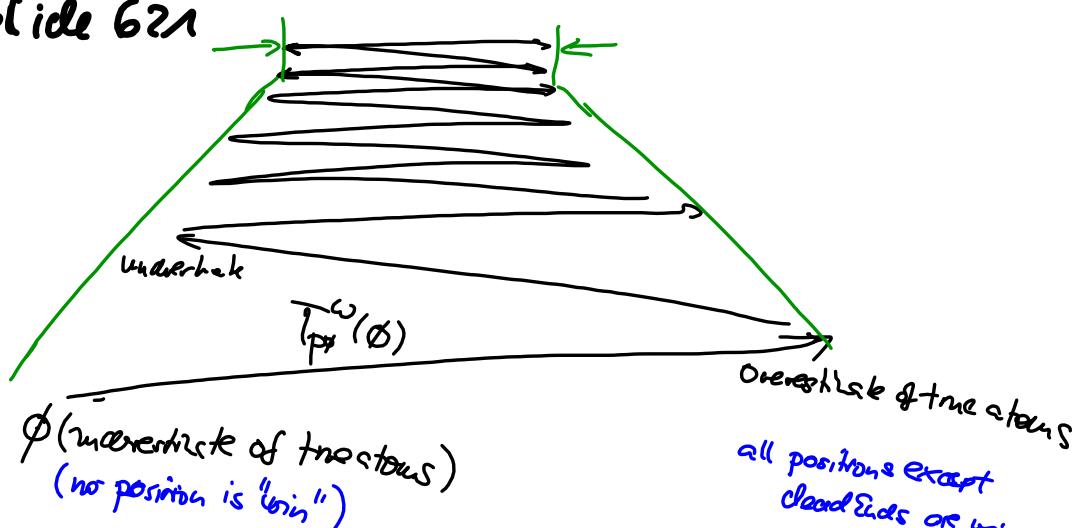
\Rightarrow no!

$\text{reachable } (IS, R) \notin T_{P^{\text{min}}}^{\omega}(\emptyset)$

(has no support by the EDB facts r_1, r_2, r_3)

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With more game : (only win(X) listed)

$$J_0 = \emptyset$$

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$J_0 = \emptyset \rightsquigarrow P^P = \{ \text{win}(a) :- \text{move}(a,b), \text{win}(b). \}$ *satisfied in \emptyset*

$\text{win}(b) :- \text{move}(b,c), \text{win}(c).$ *take any grand
children where
move... is true*

$\text{win}(b) :- \text{move}(b,d), \text{win}(d).$

$\text{win}(a) :- \text{move}(a,e), \text{win}(e).$

\vdots

$= T_{P^P(\emptyset)} = \{ \text{win}(a), \text{win}(b), \dots \}$ *$\text{win}(f), \text{win}(g), \text{win}(h), \text{win}(i)$ are not true*

= an overestimate of 'win'

but some nodes are already excluded :-

\Rightarrow Continue : exercise \rightarrow 20.6.

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