

Unsafe Formulas :

$$F_1(x) = x < 3$$

$$F_2(x) = \neg \text{borders}(\cancel{x}, x)$$

$$F_3(x, y, z) = p(x, y) \vee q(x, z)$$

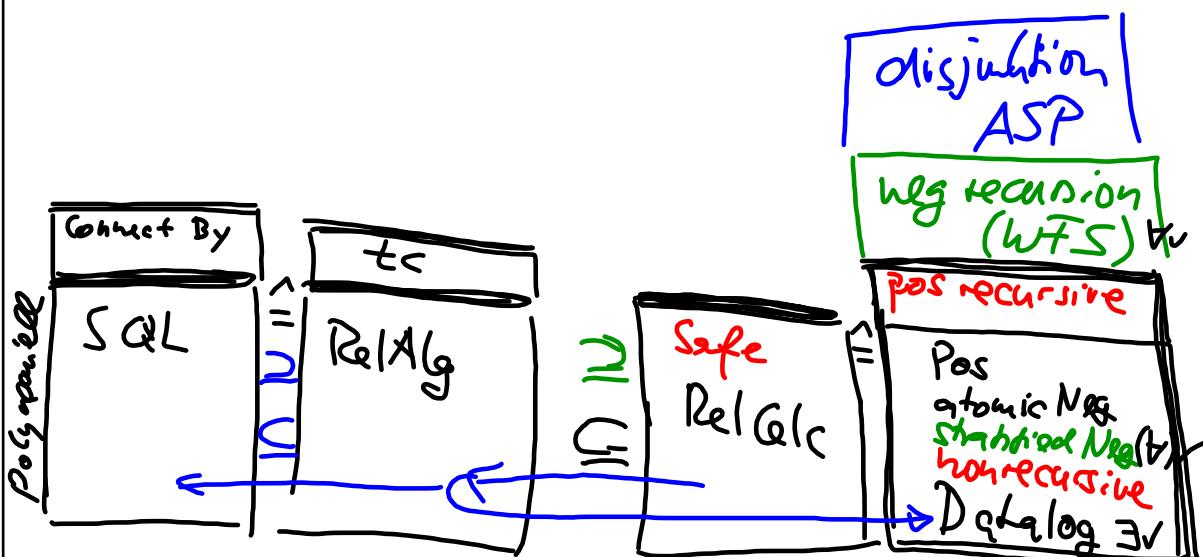
$$F_4(x) = \exists c: \text{cont}(c) \wedge \text{area}(c, x) \wedge x < 3$$

$$\{(x_{1/0,44}), (x_{1/1,43}), \dots\}$$

$$F_2'(x) = \exists c: \text{cont}(c) \wedge \text{code}(c, x) \wedge \neg \text{border}(\cancel{x}, x)$$

Nov 4-14:08

$$\forall x p(x) = \neg \exists x: \neg p(x)$$



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SMART Ink

Example: Win-Move-Game

- 2 players,
- places on a board that are connected by (directed) moves (relation "move(x,y)"),
- first player puts a pebble on a position,
- players alternately move the pebble from x to a connected y ,
- if a player cannot move, he loses.
- Question: which positions are "winning" positions, "losing" position, or "drawn" positions?

The following program "describes" the game:

```
win(X) :- move(X,Y), not win(Y).
```

- the dependency graph contains a negative cycle:
no recursion

589

Nov 4-14:51

$$\begin{array}{c}
 p(x, z) \wedge z = y \wedge q(z, y) \\
 \underbrace{\qquad\qquad}_{m = \{x\}} : \qquad \underbrace{\qquad\qquad}_{m = \{y\}} \\
 \underbrace{\qquad\qquad}_{m = \{x\}} : \\
 \qquad\qquad\qquad \underbrace{\qquad\qquad\qquad}_{m = \{x, y, z\}}
 \end{array}$$

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