

Ex 1, Sheet 4

Algebra	Algebra	Algebra + tc
Difference Stratified Data Log	Union \cup Projection Π Selection σ (Projection σ join \bowtie)	positive Data Log could is normal

Union $u = p \cup q$
 $\Rightarrow p$ and q must have the same format

$u(x_1 \dots x_n) := p(x_1 \dots x_n)$
 $u(x_1 \dots x_n) := q(x_1 \dots x_n)$

Difference $d = p \setminus q \Rightarrow p, q$ have the same format
 $d(x_1 \dots x_n) := p(x_1 \dots x_n) \wedge \neg q(x_1 \dots x_n)$
 is this rule safe? - yes: all vars in $\neg q(\dots)$ also occur positively in p

Projection $pr(x_1 \dots x_k) = \Pi_{x_1 \dots x_k} p(x_1 \dots x_n)$
 $pr(x_1 \dots x_k) := p(x_1 \dots x_n)$

Selection $s(x_1 \dots x_n) = \sigma[\alpha](p(x_1 \dots x_n))$
 $s(x_1 \dots x_n) := p(x_1 \dots x_n) \wedge \alpha$

$\alpha(x_1 \dots x_n, x_{n+1} \dots x_m) = p(x_1 \dots x_n, x_{n+1} \dots x_m) \wedge q(x_1 \dots x_n, x_{n+1} \dots x_m)$
 \Rightarrow recursive De Morgan - program is stratifiable, because each head atom repeats a level in the algebra tree and uses only lower predicates
 we keep with stratification

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Ex 2

consider first a rule with only positive body literals:

$$B(x_1 \dots x_n, y_1 \dots y_m, z_1 \dots z_k) :- p_1(x_1 \dots x_n), p_2(y_1 \dots y_m), \dots, p_k(z_1 \dots z_k).$$

\Rightarrow Algebra

general case: also negative literals in the body

idea

$B(\dots) :- C_1 \dots C_n, D_1 \dots D_k$
 Each $D_i = \neg E_i$

! for " \neg " both sides must have the same format

all vars/atom names here must also occur

all vars occur in the left side \Rightarrow must also occur on the right side!

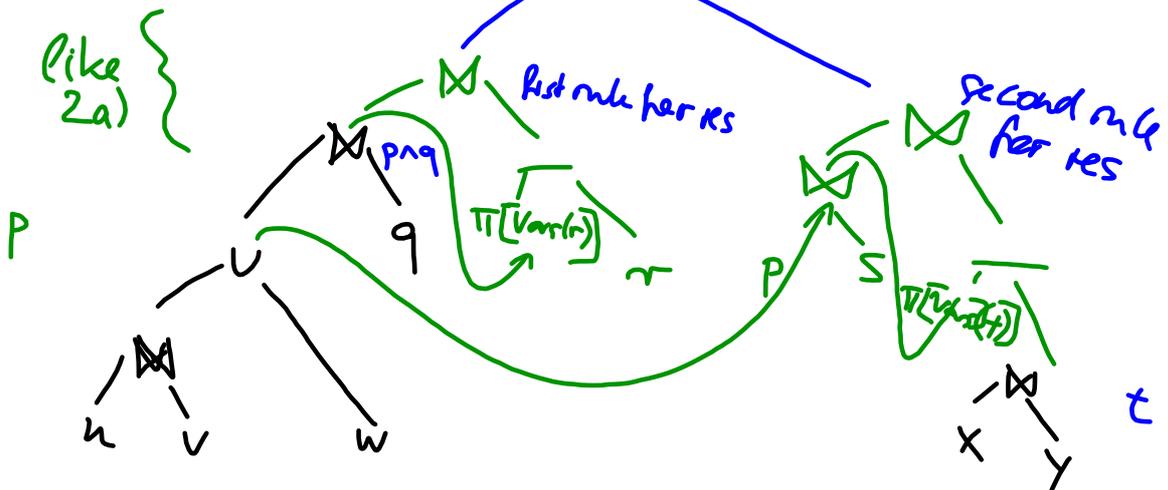
sufficient: take $\Pi[V]C_i \theta$ = was the vars occur in $E_1 \dots E_k = V$ same problem as for RAUF

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2b) done: individual rules ✓
 other necessary construct.
 two rules with same head
 $P :- \text{body}_1.$ \Rightarrow from c): body_1 equiv to ex_1
 $P :- \text{body}_2.$ body_2 equiv to ex_2
 $\rightarrow P$ is equiv to $ex_1 \cup ex_2$

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2c) example



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Ex 3)

→ example: take the reachable/NotReachable prog from slide 509

non-terminating:

$\mathcal{M}(P \cup \text{non-terminating}) \models \text{reachable}(D, P)$
 $\mathcal{M}(P \cup \text{non-terminating}) \models \text{not_reachable}(D, GS)$
 $\mathcal{M}(P \cup \text{non-terminating} \cup \{\text{order}(GS, F, 100)\}) \models \text{reachable}(D, GS)$
 $\not\models \text{not_reachable}(D, GS)$

multiple minimal models:

$\mathcal{M}(P \cup \text{non-terminating})$ is minimal

∴ one cannot remove any atoms to obtain a smaller model

$\mathcal{M}(P \cup \text{non-terminating}) \setminus \{\text{not_reachable}(D, GS)\} \cup \{\text{reachable}(D, GS)\}$
 is also minimal, and it is a model!

⇒ it is not a "nice" model because reachable is a EDB predicate, and the literal "reachable(D, GS)" is not supported by facts.

⇒ it is not well-founded!

⇒ the well-founded semantics coincides with the stratified semantics for stratifiable Datalog programs

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Ex 4

IDB predicate lang('CH', gene, GS)

:- include (mandiant). country(C, N, Pop, A, Gp, Gprou)
 :- autotable.

exA(C) :- language('CH', X, -),
 language(C, X, -).

exB(C) :- country(C, -, -|-|-|-|-),
 not exA(C).

c) nonCHlgGou(C) :- all countries, where some Co. is broken which is not ghou in CH
 lang(C, -, L, -)
 not lang('CH', -, L, -).
 ...and we again need the remaining countries

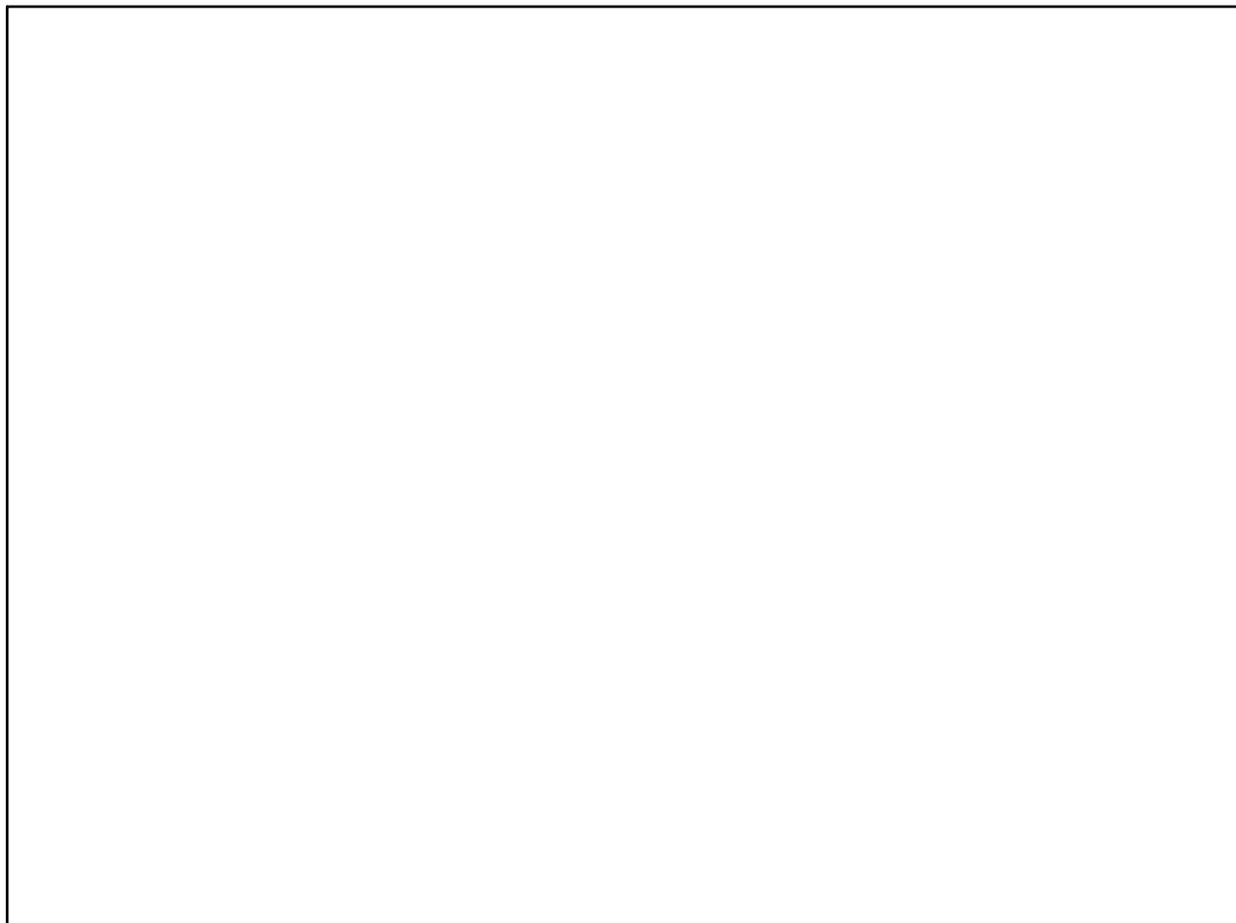
onlyCHlgGou(C) :- country(C, -, -, -, -|-|-|-),
 not nonCHlgGou(C).

d) ~~chlg(L) :- language('CH', L).~~

chlgLissig(C) :- if some str or lg. is missing in C
 language('CH', L, -), not lang(C, L, -)

exD(C) :- ~~country(C, -----),
 not chlgLissig(C).~~ country(C, -----)

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