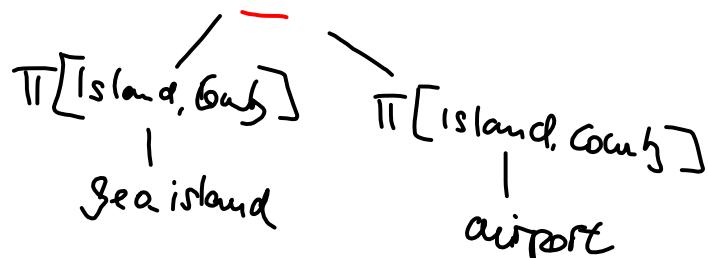


SQL: negation and closed (?) world:  
Calculus query:

$$T(I, C) \equiv \exists \text{Prov}: \text{geo\_island}(I, C, \text{Prov}) \\ \wedge \neg \exists \text{Code}, \text{Name}, \text{City}, \text{Prov}: \\ \text{airport}(\text{Code}, \text{Name}, \\ C, \text{City}, \text{Prov}, I)$$


Apr 29-10:04

SQL:

```
select island, country
from geo_island
where (island, country) not in (select island, country from airport)
```

51 results

```
(select island, country
from geo_island g)
minus
(select island, country
from airport a)
```

distinct

```
select island, country
from geo_island g
where not exists
  (select * from airport a
  where a.island = g.island
  and a.country = g.country)
```

↑  
Same?

132 results

✓?  
.

135 results

Apr 29-10:39

Reason: NULL values

SQL: null values do not violate any constraint?  
 Condition

select name, country, city, island from airport where country='D'

$\Rightarrow$  null values for island - column  
 these airports might be on any island

$\Rightarrow$  actual answer is based on  
 Open world  
 ( Null values with " | N" )

Apr 29-10:57

Slide 477

relational style:

$\forall x, a : \text{person}(x) \wedge \text{age}(x, a) \wedge a \geq 18 \rightarrow \text{adult}(x)$

functional style.

$\forall x : \text{person}(x) \wedge \text{age}(x) \geq 18 \rightarrow \text{adult}(x)$

( Description logics (SemWeb)  
 $\text{Person} \sqcap \text{age} \geq 18 \sqsubseteq \text{Adult}$  )

SQL view:

Create View Adult as

(Select \* from Person

where age  $\geq 18$  )

"Simple" positive RULE

$\rightarrow$  can be expressed as SQL view

Apr 29-11:08

$\forall x : (\exists y : \text{hasChild}(x,y) \rightarrow \text{parent}(x))$  ← existential knowledge  
 not a simple rule

$\forall x, y : \text{hasChild}(x,y) \rightarrow \text{parent}(x)$   
 can be converted into a simple rule  
 no " $\leftarrow$ " direction

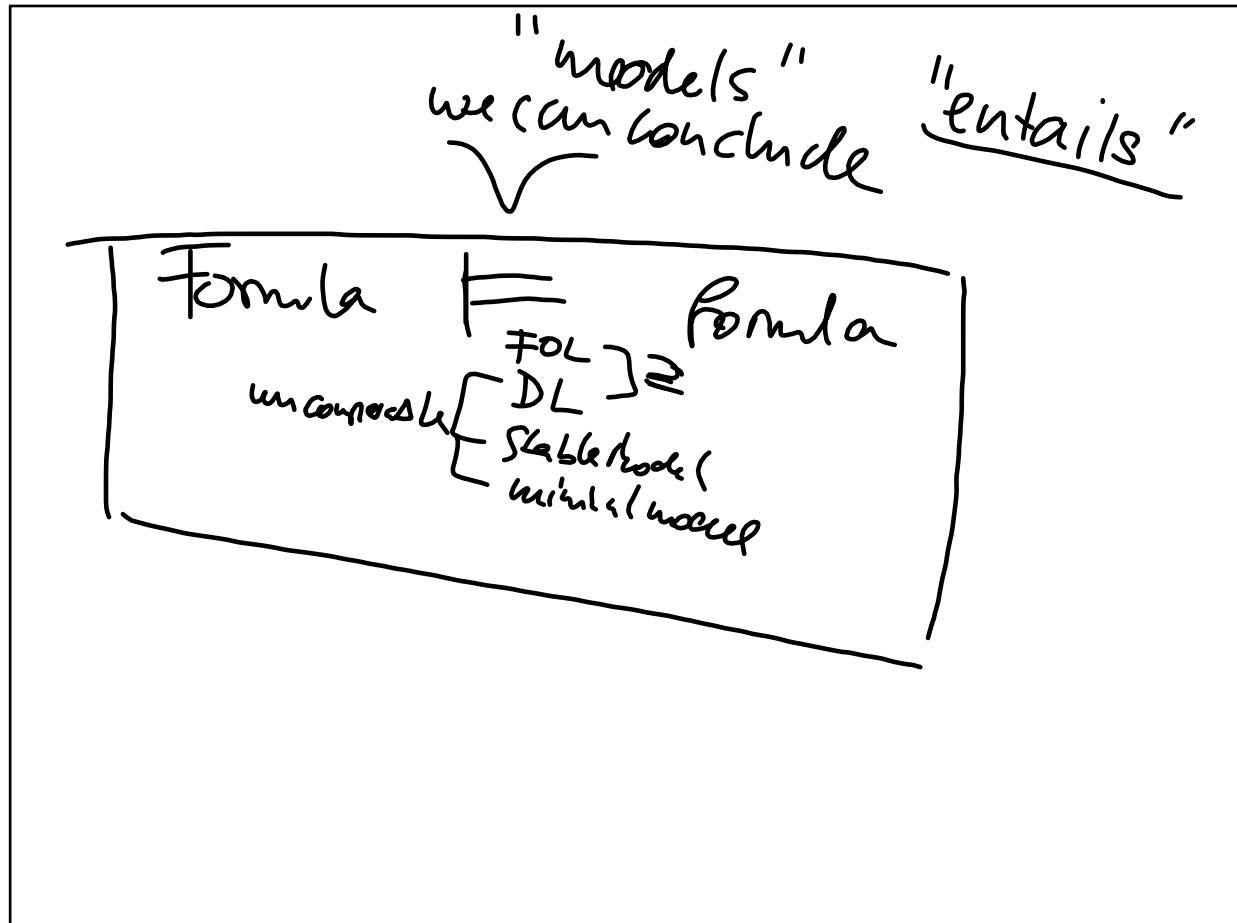
DL:  
 $\text{Parent} \equiv \exists \text{hasChild}. T$   
 → cannot be closed in SQL!

Apr 29-11:10

Consider FOL tableau calculus

$\forall x (\exists y : \text{hasChild}(x,y) \leftarrow \text{parent}(x))$   
 parent (many)  
 $\neg \text{parent}(x)$        $\exists y : \text{hasChild}(x,y)$   
 ;      hasChild( $x, f_{hc}(x)$ )  
 $\square x \rightarrow \text{many}$       hasChild(many),  $f_{hc}(\text{many})$   
 3. rule       $\Rightarrow$  one of the children of many  
 Sholem fact:      children of many  
 $\Rightarrow$  reasoning can deal with existential knowledge

Apr 29-11:16



Apr 29-11:21