

Ex. 1 d)

$F(x) \equiv p(x) \wedge \exists y (q(y) \wedge \neg r(x,y))$

customer(x) \wedge $\exists y$ (product(y) \wedge \neg ordered(x,y))

\Rightarrow ~~$p(x) \wedge \exists y (q(y) \wedge \neg r(x,y))$~~ *redundant*

$\pi[x] (p(x) \wedge q(y) \wedge \neg r(x,y))$

Tree diagram for $\pi[x]$:
 $\pi[x]$ branches to x , which branches to p and q . $\neg r$ is also shown as a branch.

Handwritten annotations:
 $\pi = \emptyset$ (tree = $\{x\}$)
 $\pi = \{y\}$ (tree = $\{x, y\}$)
 $\pi = \emptyset$
 $\pi = \{x, y\}$

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1 ~~Q~~: $\neg \exists$

$p(x) \wedge \neg \exists y (q(y) \wedge \neg r(x,y))$

\Leftarrow relational division pattern

\Rightarrow "minus"

intuitive way: as always done for the relational division

Tree diagram for $\pi[x]$:
 $\pi[x]$ branches to x , which branches to p and q . $\neg r$ is also shown as a branch.

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10)

free = {x, y, v}
 π = {v, x}

free = {y, x, w}
 π = {w, y}

$$\exists v \left(\underbrace{\tau(v, x) \wedge \tau s(x, y, v)}_{\exists v} \right) \wedge \left(\underbrace{\tau(w, y) \wedge \tau s(y, x, w)}_{\exists w} \right)$$

← push-into-exists

π = {x, y}
 free = {x, y} ... all ... safe

push into ... right side, but not RANF

$$F_1 \wedge \exists w (F_1 \wedge Q)$$

$$F_1 \wedge \exists w \left(\exists v \left(\tau(v, x) \wedge \tau s(x, y, v) \right) \wedge \left(\tau(w, y) \wedge \tau s(y, x, w) \right) \right)$$

π = {v, x, y}
 free = {v, x, y}

$$\exists v \exists w \left(\tau(v, x) \wedge \tau s(x, y, v) \wedge \tau(w, y) \wedge \tau s(y, x, w) \right)$$

free = {v, x, y, w}
 π = {v, x, y, w}

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instead: reformulate formula at the beginning:

F(x, y)

$$\exists v \left(\tau(v, x) \wedge \tau s(x, y, v) \right) \wedge \exists w \left(\tau(w, y) \wedge \tau s(y, x, w) \right)$$

$$\equiv \exists v, w : \left(\tau(v, x) \wedge \tau(w, y) \wedge \tau s(x, y, v) \wedge \tau s(y, x, w) \right)$$

π = {v, x, w, y} ✓
 free = {v, x, w, y}

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another possibility with push:

$$\exists v \left(\underbrace{\tau(v,x) \wedge \neg s(x,y,v)}_{\text{problem: only free } x \rightarrow \text{need pos. bindid for free } x} \right) \wedge \exists w \left(\tau(w,y) \wedge \neg s(y,x,w) \right)$$

$$\tau = \{w, y\}$$

$$\text{free} = \{y, x, w\}$$

$$F_1 \wedge \exists w \left(\tau(v,x) \wedge \tau(w,y) \wedge \neg s(y,x,w) \right)$$

$$\tau = \{v, x, w, y\}$$

$$\text{free} = \{v, x, w, y\} \checkmark$$

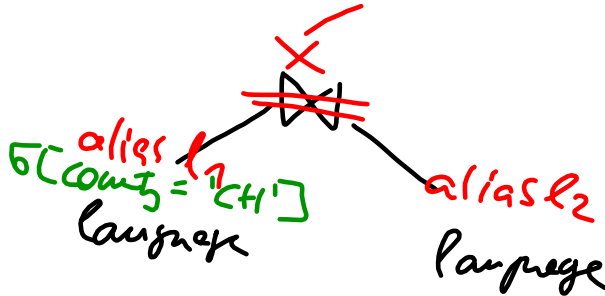
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Ex. 2a)

select I2.country, I1.name
 from language I1, language I2
 where I1.country = 'CH'
 and I1.name = I2.name

$\pi [I2.country]$

$\sigma [I1.name = I2.name]$



$F(\sigma) = \exists$

$L, P_1, P_2 :$

$language('CH', L, P_1) \wedge language(L, P_2)$

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2b)

$$F(C) \equiv \exists N, P, A, G, P$$

country(C, N, P, A, G, P) :

$\wedge \forall$ Languages in that country

$\forall L (\exists P_2$ language(C, L, P₂)

\rightarrow not spoken in CH

$\neg \exists P_1$: language('CH', L, P₁))

$\forall \rightarrow \exists$

\equiv

$$\exists N, P, A, G, P$$

country(C, N, P, A, G, P) :

$\wedge \exists L (\exists P_2$: language(C, L, P₂)

\wedge language('CH', L, P₁))

Ex 2 a (select code from country)

MINUS

(select l2.country

from language l1, language l2

where l1.country = 'CH'

and l1.name = l2.name)

2a

country

[code]

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c)

$$F(C) \equiv \exists N, P, A, G, P, CP$$

country(C, N, P, A, G, P, CP)

$\wedge \forall L (\exists P_1$ language(C, L, P₁)

$\rightarrow \exists P_2$ language('CH', L, P₂))

$\equiv \exists N, P, A, G, P, CP$:

country(C, N, P, A, G, P, CP)

$\wedge \neg \exists L (\exists P_1$ language(C, L, P₁)

$\wedge \neg \exists P_2$: language('CH', L, P₂))

$\hat{=}$ all countries, except those

where a non-swiss lg. is spoken

NonSwiss lg country

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2c SQL:

(select code from country)

MINUS

(select l1.country

from language l1

where l1.name not in (select name from language

where country='CH'))

$\hat{=}$ $\neg \exists$

Swiss languages

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d)

$F(C) = \exists N, P, A, C_0, C_P$

$country(C, N, P, A, C_0, C_P)$

$\wedge \forall l \text{ in 'CH'} \rightarrow \text{is spoken in C}$

$\forall L: (\exists P_1: language('CH', L, P_1)$

$\rightarrow \exists P_2: language(C, L, P_2))$

$\wedge \exists \dots country(C, \dots)$

$\wedge \neg \exists L (\exists P_1: language('CH', L, P_1) \wedge \neg \exists P_2: language(C, L, P_2))$

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