Chapter 3
Relational Database Languages: Relational Algebra

We first consider only *query* languages.

Relational Algebra: Queries are expressions over operators and relation names.

Relational Calculus: Queries are special formulas of first-order logic with free variables.

SQL: Combination from algebra and calculus and additional constructs. Widely used DML for relational databases.

QBE: Graphical query language.

Deductive Databases: Queries are logical rules.

Relational Database Languages: Comparison and Outlook

Remark:

- Relational Algebra and (safe) Relational Calculus have the same expressive power.
  For every expression of the algebra there is an equivalent expression in the calculus, and vice versa.
- A query language is called **relationally complete**, if it is (at least) as expressive as the relational algebra.
- These languages are compromises between efficiency and expressive power; they are not computationally complete (i.e., they cannot simulate a Turing Machine).
- They can be embedded into host languages (e.g. C++ or Java) or extended (PL/SQL), resulting in full computational completeness.
- Deductive Databases (Datalog) are more expressive than relational algebra and calculus.

3.1 Relational Algebra: Computations over Relations

Operations on Tuples – Overview Slide

Let $\mu \in \text{Tup}(\bar{X})$ where $\bar{X} = \{A_1, \ldots, A_k\}$.

(Formal definition of μ see Slide 60)

- For $\emptyset \subset \bar{Y} \subseteq \bar{X}$, the expression $\mu[\bar{Y}]$ denotes the **projection** of μ to $\bar{Y}$.

  Result: $\mu[\bar{Y}] \in \text{Tup}(\bar{Y})$ where $\mu[\bar{Y}](A) = \mu(A), A \in \bar{Y}$.

- A **selection condition** α (wrt. $\bar{X}$) is an expression of the form $A\,\theta\,B$ or $A\,\theta\,c$, or $c\,\theta\,A$ where $A, B \in \bar{X}$, dom(A) = dom(B), $c \in$ dom(A), and θ is a comparison operator on that domain like e.g. $\{=, \neq, \leq, <, \geq, >\}$.

  A tuple $\mu \in \text{Tup}(\bar{X})$ **satisfies** a selection condition α, if – according to α – $\mu(A)\,\theta\,\mu(B)$ or $\mu(A)\,\theta\,c$, or $c\,\theta\,\mu(A)$ holds.

  These (atomic) selection conditions can be combined to formulas by using $\wedge$, $\vee$, $\neg$, and $(,)$.

- For $\bar{Y} = \{B_1, \ldots, B_k\}$, the expression $\mu[A_1 \to B_1, \ldots, A_k \to B_k]$ denotes the **renaming** of μ.

  Result: $\mu[\ldots, A_i \to B_i, \ldots] \in \text{Tup}(\bar{Y})$ where $\mu[\ldots, A_i \to B_i, \ldots](B_i) = \mu(A_i)$ for $1 \leq i \leq k$.

Let $\mu \in \text{Tup}(\bar{X})$ where $\bar{X} = \{A_1, \ldots, A_k\}$.

Projection

For $\emptyset \subset \bar{Y} \subseteq \bar{X}$, the expression $\mu[\bar{Y}]$ denotes the **projection** of μ to $\bar{Y}$.

Result: $\mu[\bar{Y}] \in \text{Tup}(\bar{Y})$ where $\mu[\bar{Y}](A) = \mu(A),\ A \in \bar{Y}$.

projection to a given set of attributes

Example 3.1

Consider the relation schema $R(\bar{X}) = continent(Name, Area)$: $\bar{X} = [Name,\ Area]$

and the tuple μ = | Name → "Asia", Area → 4.50953e+07 |.

formally: $\mu(Name)$ = "Asia", $\mu(Area) = 4.5E7$

projection attributes: Let $\bar{Y} = [Name]$

Result: $\mu[Name]$ = | Name → "Asia" | □

Again, $\mu \in \text{Tup}(\bar{X})$ where $\bar{X} = \{A_1, \ldots, A_k\}$.

Selection

A **selection condition** α (wrt. $\bar{X}$) is an expression of the form $A\,\theta\,B$ or $A\,\theta\,c$, or $c\,\theta\,A$ where $A, B \in \bar{X}$, $\text{dom}(A) = \text{dom}(B)$, $c \in \text{dom}(A)$, and θ is a comparison operator on that domain like e.g. $\{=, \neq, \leq, <, \geq, >\}$.

A tuple $\mu \in \text{Tup}(\bar{X})$ **satisfies** a selection condition α, if – according to α – $\mu(A)\,\theta\,\mu(B)$ or $\mu(A)\,\theta\,c$, or $c\,\theta\,\mu(A)$ holds.

yes/no-selection of tuples (without changing the tuple)

Example 3.2

Consider again the relation schema $R(\bar{X}) = continent(Name, Area)$*:* $\bar{X} = [Name,\ Area]$.

Selection condition: $Area > 10.000.000$.

Consider again the tuple $\mu =$ | *Name → "Asia", Area → 4.50953e+07* |.

formally: $\mu(Name) =$ *"Asia"*, $\mu(Area) = 4.5E7$

check: $\mu(Area) > 10.000.000$

Result: yes. □

These (atomic) selection conditions can be combined to formulas by using $\wedge$, $\vee$, $\neg$, and $($, $)$.

Let $\mu \in \text{Tup}(\bar{X})$ where $\bar{X} = \{A_1, \ldots, A_k\}$.

Renaming

For $\bar{Y} = \{B_1, \ldots, B_k\}$, the expression $\mu[A_1 \to B_1, \ldots, A_k \to B_k]$ denotes the **renaming** of μ.

Result: $\mu[\ldots, A_i \to B_i, \ldots] \in \text{Tup}(\bar{Y})$ where $\mu[\ldots, A_i \to B_i, \ldots](B_i) = \mu(A_i)$ for $1 \leq i \leq k$.

renaming of attributes (without changing the tuple)

Example 3.3

Consider (for a tuple of the table $R(\bar{X}) = encompasses(Country, Continent, Percent)$*):*

$\bar{X} = [Country,\ Continent,\ Percent]$.

Consider the tuple $\mu =$ | *Country → "R", Continent → "Asia", Percent → 80* |.

formally: $\mu(Country) =$ *"R"*, $\mu(Continent) =$ *"Asia"*, $\mu(Percent) = 80$

Renaming: $\bar{Y} = [Code, Name, Percent]$

Result: a new tuple $\mu[Country \to Code,\ Continent \to Name,\ Percent \to Percent] =$ | *Code → "R", Name → "Asia", Percent → 80* | *that now fits into the schema* $new\_encompasses(Code, Name, Percent)$. □

The usefulness of renaming will become clear later ...

Expressions in the Relational Algebra

What is an algebra?

- An algebra consists of a "domain" (i.e., a set of "things"), and a set of operators.
- Operators map elements of the domain to other elements of the domain.
- Each of the operators has a "semantics", that is, a definition how the result of applying it to some input should look like.
- **Algebra expressions** are built over basic constants and operators (inductive definition).

Relational Algebra

- The "domain" consists of all relations (over arbitrary sets of attributes).
- The operators are then used for combining relations, and for describing computations - e.g., in SQL.
- **Relational algebra expressions** are defined inductively over relations and operators.
- Relational algebra expressions define queries against a relational database.

Inductive Definition of Expressions

Atomic Expressions

- For an arbitrary attribute A and a constant $a \in$ dom(A), the **constant relation** $A : \{a\}$ is an algebra expression.

  Format: $[A]$
  Result relation: $\{a\}$

A:{a}
A
a

- Given a database schema $\mathbf{R} = \{R_1(\bar{X}_1), \ldots, R_n(\bar{X}_n)\}$, every relation name R_i is an algebra expression.

  Format of R_i: $\bar{X}_i$
  Result relation (wrt. a given database state $\mathcal{S}$): the relation $\mathcal{S}(R_i)$ that is currently stored in the database.

Structural Induction: Applying an Operator

- takes one or more input relations $in_1, in_2, \ldots$
- produces a result relation out:
  - out has a **format**,
    depends on the formats of the input relations.
  - out is a relation, i.e., it contains some tuples,
    depends on the content of the input relations.

BASE OPERATORS

Let $\bar{X}, \bar{Y}$ formats and $r \in \text{Rel}(\bar{X})$ and $s \in \text{Rel}(\bar{Y})$ relations over $\bar{X}$ and $\bar{Y}$.

Union

Assume $r, s \in \text{Rel}(\bar{X})$.
Result format of $r \cup s$: $\bar{X}$
Result relation: $r \cup s = \{\mu \in \text{Tup}(\bar{X}) \mid \mu \in r \text{ or } \mu \in s\}$.

$r =$

A	B	C
a	b	c
d	a	f
c	b	d

$s =$

A	B	C
b	g	a
d	a	f

$r \cup s =$

A	B	C
a	b	c
d	a	f
c	b	d
b	g	a

Set Difference

Assume $r, s \in \text{Rel}(\bar{X})$.
Result format of $r \setminus s$: $\bar{X}$
Result relation: $r \setminus s = \{\mu \in r \mid \mu \notin s\}$.

$$r = \begin{array}{ccc} A & B & C \\ \hline a & b & c \\ d & a & f \\ c & b & d \end{array} \qquad s = \begin{array}{ccc} A & B & C \\ \hline b & g & a \\ d & a & f \end{array} \qquad r \setminus s = \begin{array}{ccc} A & B & C \\ \hline a & b & c \\ c & b & d \end{array}$$

Projection

Assume $r \in \text{Rel}(\bar{X})$ and $\bar{Y} \subseteq \bar{X}$.
Result format of $\pi[\bar{Y}](r)$: $\bar{Y}$
Result relation: $\pi[\bar{Y}](r) = \{\mu[\bar{Y}] \mid \mu \in r\}$.

Example 3.4

Continent	
Name	***Area***
Europe	*9562489.6*
Africa	*3.02547e+07*
Asia	*4.50953e+07*
America	*3.9872e+07*
Australia	*8503474.56*

Let $\bar{Y} = [Name]$

$\mu_1[Name] =$ *Name → "Europe"*
$\mu_2[Name] =$ *Name → "Africa"*
$\mu_3[Name] =$ *Name → "Asia"*
$\mu_4[Name] =$ *Name → "America"*
$\mu_5[Name] =$ *Name → "Australia"*

$\pi[Name]$(***Continent***)
Name
Europe
Africa
Asia
America
Australia

□

Selection

Assume $r \in \text{Rel}(\bar{X})$ and a selection condition α over $\bar{X}$.

Result format of $\sigma[\alpha](r)$: $\bar{X}$
Result relation: $\sigma[\alpha](r) = \{\mu \in r \mid \mu \text{ satisfies } \alpha\}$.

Example 3.5

Continent	
Name	***Area***
Europe	*9562489.6*
Africa	*3.02547e+07*
Asia	*4.50953e+07*
America	*3.9872e+07*
Australia	*8503474.56*

Let $\alpha =$ *"*$Area > 10.000.000$*"*

$\mu_1(Area) > 10.000.000$?– *no*
$\mu_2(Area) > 10.000.000$?– *yes*
$\mu_3(Area) > 10.000.000$?– *yes*
$\mu_4(Area) > 10.000.000$?– *yes*
$\mu_5(Area) > 10.000.000$?– *no*

$\sigma[Area > 10E6]$(***Continent***)	
Name	***Area***
Africa	*3.02547e+07*
Asia	*4.50953e+07*
America	*3.9872e+07*

□

Renaming

Assume $r \in \text{Rel}(\bar{X})$ with $X = [A_1, \ldots, A_k]$ and a renaming $[A_1 \to B_1, \ldots, A_k \to B_k]$.

Result format of $\rho[A_1 \to B_1, \ldots, A_k \to B_k](r)$: $[B_1, \ldots, B_k]$
Result relation: $\rho[A_1 \to B_1, \ldots, A_k \to B_k](r) = \{\mu[A_1 \to B_1, \ldots, A_k \to B_k] \mid \mu \in r\}$.

Example 3.6
Consider the renaming of the table $encompasses(Country, Continent, Percent)$*:*

$\bar{X} = [Country,\ Continent,\ Percent]$
Renaming: $\bar{Y} = [Code, Name, Percent]$

$\rho[Country \to Code,\ Continent \to Name,\ Percent \to Percent]$(***encompasses***)		
Code	***Name***	***Percent***
R	*Europe*	*20*
R	*Asia*	*80*
D	*Europe*	*100*
⋮	⋮	⋮

□

(Natural) Join

Assume $r \in \text{Rel}(\bar{X})$ and $s \in \text{Rel}(\bar{Y})$ for arbitrary $\bar{X}, \bar{Y}$.

Convention: Instead of $\bar{X} \cup \bar{Y}$, we also write $\overline{XY}$.
for two tuples $\mu_1 = \boxed{v_1, \ldots, v_n}$ and $\mu_2 = \boxed{w_1, \ldots, w_m}$, $\mu_1\mu_2 := \boxed{v_1, \ldots, v_n, w_1, \ldots, w_m}$.

Result format of $r \bowtie s$: $\overline{XY}$.
Result relation: $r \bowtie s = \{\mu \in \text{Tup}(\overline{XY}) \mid \mu[\bar{X}] \in r \text{ and } \mu[\bar{Y}] \in s\}$.

Motivation

Simplest Case: $\bar{X} \cap \bar{Y} = \emptyset \Rightarrow$ Cartesian Product $r \bowtie s = r \times s$
$r \times s = \{\mu_1\mu_2 \in \text{Tup}(\overline{XY}) \mid \mu_1 \in r \text{ and } \mu_2 \in s\}$.

$r =$

A	B
1	2
4	5

$s =$

C	D
a	b
c	d
e	f

$r \bowtie s =$

A	B	C	D
1	2	a	b
1	2	c	d
1	2	e	f
4	5	a	b
4	5	c	d
4	5	e	f

Example 3.7 (Cartesian Product of Continent and Encompasses)

Continent $\times$ ***encompasses***				
Name	***Area***	***Continent***	***Country***	***Percent***
Europe	*9562489.6*	*Europe*	*Germany*	*100*
Europe	*9562489.6*	*Europe*	*Russia*	*20*
Europe	*9562489.6*	*Asia*	*Russia*	*80*
Europe	*9562489.6*	:	:	:
Africa	*3.02547e+07*	*Europe*	*Germany*	*100*
Africa	*3.02547e+07*	*Europe*	*Russia*	*20*
Africa	*3.02547e+07*	*Asia*	*Russia*	*80*
Africa	*3.02547e+07*	:	:	:
Asia	*4.50953e+07*	*Europe*	*Germany*	*100*
Asia	*4.50953e+07*	*Europe*	*Russia*	*20*
Asia	*4.50953e+07*	*Asia*	*Russia*	*80*
Asia	*4.50953e+07*	:	:	:
:	:	:	:	:

□

Back to the Natural Join

General case $\bar{X} \cap \bar{Y} \neq \emptyset$: shared attribute names constrain the result relation.

Again the definition: $r \bowtie s = \{\mu \in \text{Tup}(\overline{XY}) \mid \mu[\bar{X}] \in r \text{ and } \mu[\bar{Y}] \in s\}$.

(Note: this implies that the tuples $\mu_1 := \mu[\bar{X}] \in r$ and $\mu_2 := \mu[\bar{Y}] \in s$ coincide in the shared attributes $\bar{X} \cap \bar{Y}$)

Example 3.8

Consider encompasses(country,continent,percent) and isMember(organization,country,type):

encompasses		
Country	***Continent***	***Percent***
R	*Europe*	*20*
R	*Asia*	*80*
D	*Europe*	*100*
:	*:*	*:*

isMember		
Organization	***Country***	***Type***
EU	*D*	*member*
UN	*D*	*member*
UN	*R*	*member*
:	*:*	*:*

$$encompasses \bowtie isMember = \{\mu \in \text{Tup}(country, cont, perc, org, type) \mid \mu[country, cont, perc] \in encompasses \text{ and } \mu[org, country, type] \in isMember\}$$

□

Example 3.8 (Continued)

$$encompasses \bowtie isMember = \{\mu \in \text{Tup}(country, cont, perc, org, type) \mid \mu[country, cont, perc] \in encompasses \text{ and } \mu[org, country, type] \in isMember\}$$

start with $(R, Europe, 20) \in encompasses$.
check which tuples in $isMember$ *match:*
- $(UN, R, member) \in isMember$ *matches:*
  result: $(R, Europe, 20, UN, member)$ *belongs to the result.*
  (some more matches ...)

continue with $(R, Asia, 80) \in encompasses$.
- $(UN, R, member) \in isMember$ *matches:*
  result: $(R, Asia, 80, UN, member)$ *belongs to the result.*
  (some more matches ...)

continue with $(D, Europe, 100) \in encompasses$.
- $(EU, D, member) \in isMember$ *matches:*
  result: $(D, Europe, 100, EU, member)$ *belongs to the result.*
- $(UN, D, member) \in isMember$ *matches:*
  result: $(D, Europe, 100, UN, member)$ *belongs to the result.*
  (some more matches ...)

□

Example 3.8 (Continued)

Result:

encompasses $\times$ ***isMember***				
Country	***Continent***	***Percent***	***Organization***	***Type***
R	*Europe*	*20*	*UN*	*member*
R	*Europe*	*20*	*:*	*:*
R	*Asia*	*80*	*UN*	*member*
R	*Asia*	*80*	*:*	*:*
D	*Europe*	*100*	*UN*	*member*
D	*Europe*	*100*	*EU*	*member*
D	*Europe*	*100*	*:*	*:*
:	*:*	*:*	*:*	*:*

□

Example 3.9 (and Exercise)

Consider the expression

$continent \bowtie \rho[Country \to Code,\ Continent \to Name,\ Percent \to Percent]($***encompasses***$)$

□

Functionalities of the Join

- Combining relations
- Selective functionality: only matching tuples survive
  (consider joining cities and organizations on headquarters)

DERIVED OPERATORS

Intersection

Assume $r, s \in$ Rel$(\bar{X})$.

Then, $r \cap s = \{\mu \in$ Tup$(\bar{X}) \mid \mu \in r$ and $\mu \in s\}$.

Theorem 3.1

Intersection can be expressed by Difference: $r \cap s = r \setminus (r \setminus s)$. □

θ-Join

Combination of Cartesian Product and Selection:

Assume $r \in \text{Rel}(\bar{X})$, and $s \in \text{Rel}(\bar{Y})$, such that $\bar{X} \cap \bar{Y} = \emptyset$, and $A\,\theta\,B$ a selection condition.

$$r \bowtie_{A\theta B} s = \{\mu \in \text{Tup}(\overline{XY}) \mid \mu[\bar{X}] \in r,\ \mu[\bar{Y}] \in s \text{ and } \mu \text{ satisfies } A\theta B\} = \sigma[A\theta B](r \times s).$$

Equi-Join

θ-join that uses the "="-predicate.

Example 3.10 (and Exercise)

Consider again Example 3.7:

$Continent \times encompasses$ *contained tuples that did not really make sense.*

$(Continent \times encompasses)_{continent=name}$ *would be more useful.*

Furthermore, consider
$\pi[continent, area, code, percent]((Continent \times encompasses)_{continent=name})$*:*

- *removes the - now redundant - "name" column,*
- *is equivalent to the natural join* $(\rho[name \to continent]continent) \bowtie encompasses$. □

Semi-Join

- recall: joins combine, but are also selective
- semi-join acts like a selection on a relation r:
  selection condition not given as a boolean formula on the attributes of r, but by "looking into" another relation (a subquery)

Assume $r \in \text{Rel}(\bar{X})$ and $s \in \text{Rel}(\bar{Y})$ such that $\bar{X} \cap \bar{Y} \neq \emptyset$.

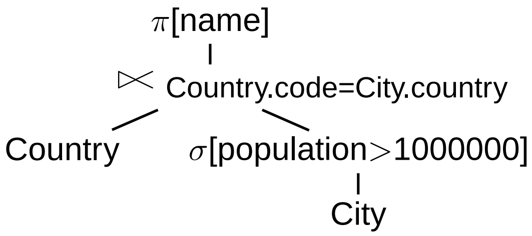
Result format of $r \ltimes s$: $\bar{X}$
Result relation: $r \ltimes s = \pi[\bar{X}](r \bowtie s)$

The semi-join $r \ltimes s$ does *not* return the join, but checks which tuples of r "survive" the join with s (i.e., "which find a counterpart in s wrt. the shared attributes"):

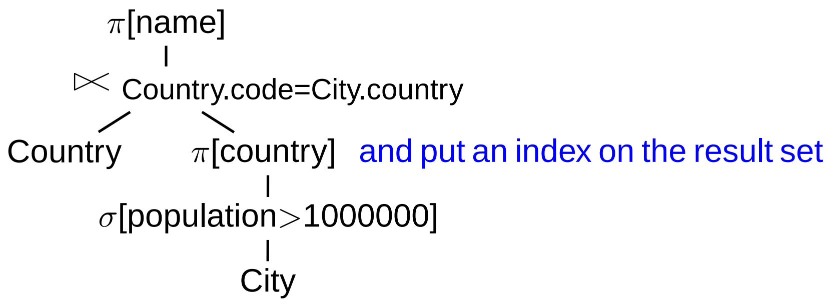
- Used with subqueries: (main query) $\ltimes$ (subquery)
- $r \ltimes s \subseteq r$
- Used for optimizing the evaluation of joins (often in combination with indexes).

Semi-Join: Example

Give the names of all countries where a city with at least 1.000.000 inhabitants is located:



- Have a short look “inside” the subquery, but dont’ actually use it:
- look only if there is a big city in this country.
- “if the country code is in the set of country codes ...”:



Outer Join

- Join is the operator for combining relations

Example 3.11

- *Persons work in divisions of a company, tools are assigned to the divisions:*

Works	
Person	***Division***
John	*Production*
Bill	*Production*
John	*Research*
Mary	*Research*
Sue	*Sales*

Tools	
Division	***Tool***
Production	*hammer*
Research	*pen*
Research	*computer*
Admin.	*typewriter*

Works $\bowtie$ ***Tools***		
Person	***Division***	***Tool***
John	*Production*	*hammer*
Bill	*Production*	*hammer*
John	*Research*	*pen*
John	*Research*	*computer*
Mary	*Research*	*pen*
Mary	*Research*	*computer*

- *join contains no tuple that describes Sue*
- *join contains no tuple that describes the administration or sales division*
- *join contains no tuple that shows that there is a typewriter* □

Outer Join

Assume $r \in \text{Rel}(\bar{X})$ and $s \in \text{Rel}(\bar{Y})$.

Result format of $r ⟗ s$: $\overline{XY}$

The outer join extends the “inner” join with all tuples that have no counterpart in the other relation (filled with null values):

Example 3.12 (Outer Join)

Consider again Example 3.11

Works ⟗ Tools		
Person	***Division***	***Tool***
John	*Production*	*hammer*
Bill	*Production*	*hammer*
John	*Research*	*pen*
John	*Research*	*computer*
Mary	*Research*	*pen*
Mary	*Research*	*computer*
Sue	*Sales*	*NULL*
NULL	*Admin*	*typewriter*

Works ⋉ Tools	
Person	***Division***
John	*Production*
Bill	*Production*
John	*Research*
Mary	*Research*

Works ⋊ Tools	
Division	***Tool***
Production	*hammer*
Research	*pen*
Research	*computer*

Formally, the result relation is defined as follows:

$J = r \bowtie s$ — take the (“inner”) join as base

$r_0 = r \setminus \pi[\bar{X}](J) = r \setminus (r \ltimes s)$ — r-tuples that “are missing”

$s_0 = s \setminus \pi[\bar{Y}](J) = s \setminus (r \rtimes s)$ — s-tuples that “are missing”

$Y_0 = \bar{Y} \setminus \bar{X}$, $X_0 = \bar{X} \setminus \bar{Y}$

Let $\mu_1 \in \text{Tup}(Y_0)$, $\mu_2 \in \text{Tup}(X_0)$ such that μ_1, μ_2 consist only of *null* values

$$r ⟗ s = J \cup (r_0 \times \{\mu_1\}) \cup (s_0 \times \{\mu_2\}).$$

Example 3.12 (Continued)

For the above example,

$J = Works \bowtie Tools$

$r_0 =$ [*“Sue”,“Sales”*], $s_0 =$ [*“Admin”,“Typewriter”*]

$Y_0 = Tool$, $X_0 = Person$

$\mu_1 =$

Tool
null

$\mu_2 =$

Person
null

$r_0 \times \{\mu_1\} =$

Person	***Division***	***Tool***
Sue	*Sales*	*null*

$s_0 \times \{\mu_2\} =$

Person	***Division***	***Tool***
null	*Admin*	*Typewriter*

□

Generalized Natural Join

Assume $r_i \subseteq \text{Tup}(\bar{X}_i)$.

Result format: $\cup_{i=1}^{n} \bar{X}_i$
Result relation: $\bowtie_{i=1}^{n} r_i = \{\mu \in \text{Tup}(\cup_{i=1}^{n} \bar{X}_i) \mid \mu[\bar{X}_i] \in r_i\}$

Exercise 3.1
Prove that the natural join is associative (which makes the generalized natural join well-defined):

$$\begin{aligned} \bowtie_{i=1}^{n} r_i &= ((\ldots((r_1 \bowtie r_2) \bowtie r_3) \bowtie \ldots) \bowtie r_n) \\ &= (r_1 \bowtie (r_2 \ldots (r_{n-1} \bowtie r_n) \ldots)) \end{aligned}$$

□

Relational Division

Assume $r \in \text{Rel}(\bar{X})$ and $s \in \text{Rel}(\bar{Y})$ such that $\bar{Y} \subsetneq \bar{X}$.
Result format of $r \div s$: $\bar{Z} = \bar{X} \setminus \bar{Y}$.

The result relation $r \div s$ is specified as "all $\bar{Z}$-values that occur in $\pi[\bar{Z}](r)$, with the additional condition that they occur in r together with each of the $\bar{Y}$ values that occur in s".

Formally,

$$r \div s = \{\mu \in \text{Tup}(\bar{Z}) \mid \{\mu\} \times s \subseteq r\} = \pi[\bar{Z}](r) \setminus \pi[\bar{Z}]((\pi[\bar{Z}](r) \times s) \setminus r).$$

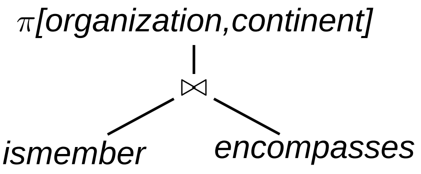
this implies that $\mu \in \pi[\bar{Z}](r)$

- Simple observation: $\pi[\bar{Z}](r) \supseteq r \div s$.
  This constrains the set of possible results.
- Often, $\bar{Z}$ and $\bar{Y}$ correspond to the keys of relations that represent the instances of entity types.

Example 3.13 (Relational Division)

Compute those organizations that have at least one member on each continent:

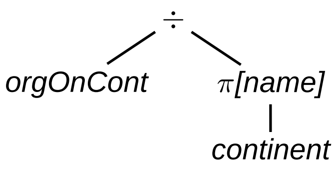
First step: which organizations have (some) member on which continents:




```
SELECT DISTINCT i.organization, e.continent
FROM ismember i, encompasses e
WHERE i.country=e.country
ORDER by 1
```

orgOnCont	
organization	***continent***
UN	*Europe*
UN	*Asia*
UN	*America*
UN	*Africa*
UN	*Australia*
NATO	*Europe*
NATO	*America*
NATO	*Asia*
:	:

Example 3.13 (Cont'd)



$r(\bar{X}),\quad s(\bar{Y}),\quad \bar{Z} := \bar{X} \setminus \bar{Y}$

$r \div s = \{\mu \in \mathit{Tup}(\bar{Z}) \mid \{\mu\} \times s \subseteq r\}$

$\bar{X} = [organization, continent]$
$\bar{Y} = [continent]$
Thus, $\bar{Z} = [organization]$.

orgOnCont	
organization	***continent***
UN	*Europe*
UN	*Asia*
UN	*America*
UN	*Africa*
UN	*Australia*
NATO	*Europe*
NATO	*America*
NATO	*Asia*
:	:

π[name](continent)
continent
Asia
Europe
Australia
America
Africa

- *UN: occurs with each continent in orgOnCont*
  $\Rightarrow$ *belongs to the result.*
- *NATO: does not occur with each continent in orgOnCont*
  $\Rightarrow$ *does not belong to the result.*

Example 3.13 (Cont'd)
Consider again the formal algebraic characterization of Division:

$$r \div s = \{\mu \in \textit{Tup}(\bar{Z}) \mid \{\mu\} \times s \subseteq r\} = \pi[\bar{Z}](r) \setminus \pi[\bar{Z}]((\pi[\bar{Z}](r) \times s) \setminus r).$$

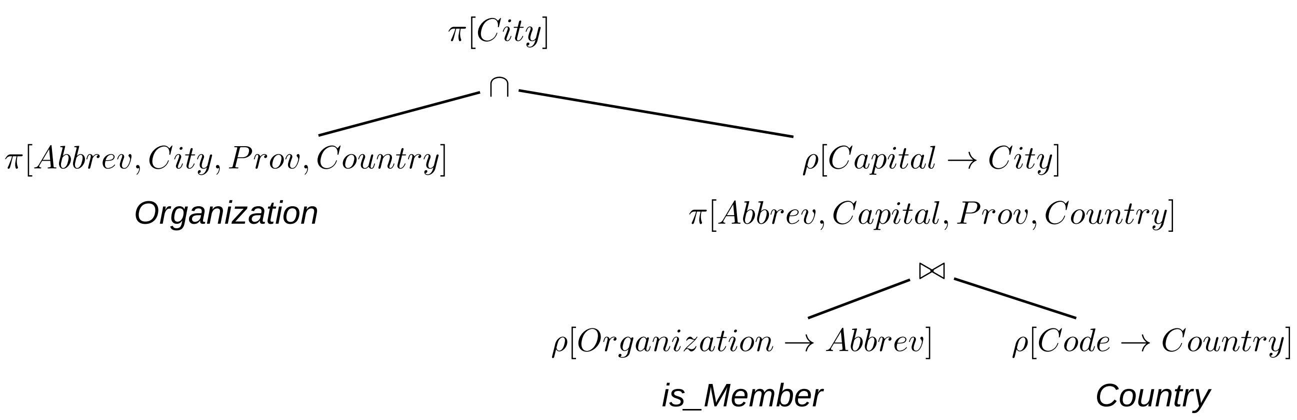
1. $r = orgOnCont$, $s = \pi$*[name](continent)*, $Z = Country$.
2. $(\pi[\bar{Z}](r) \times s)$ *contains all tuples of organizations with each of the continents, e.g., (NATO,Europe), (NATO,Asia), (NATO,America), (NATO,Africa), (NATO,Australia).*
3. $((\pi[\bar{Z}](r) \times s) \setminus r)$ *contains all such tuples which are not "valid", e.g., (NATO,Africa).*
4. *projecting this to the organizations yields all those organizations where a non-valid tuple has been generated in (2), i.e., that have no member on some continent (e.g., NATO).*
5. $\pi[\bar{Z}](r)$ *is the list of* all *organizations ...*
6. *... subtracting those computed in (4) yields those that have a member on* each *continent.* □

EXPRESSIONS

- inductively defined: combining expressions by operators

Example 3.14
The names of all cities where (i) headquarters of an organization are located, and (ii) that are capitals of a member country of this organization.

As a tree:



□

Note that there are many equivalent expressions.

Expressions in the Relational Algebra as Queries

Let $\mathbf{R} = \{R_1, \ldots, R_k\}$ a set of relation schemata of the form $R_i(\bar{X}_i)$. As already described, an **database state** to $\mathbf{R}$ is a **structure** $\mathcal{S}$ that maps every relation name R_i in $\mathbf{R}$ to a relation $\mathcal{S}(R_i) \subseteq \text{Tup}(\bar{X}_i)$

Every algebra expression Q defines a **query** against the state $\mathcal{S}$ of the database:

- For given $\mathbf{R}$, Q is assigned a **format** Σ_Q (the format of the answer).
- For every database state $\mathcal{S}$, $\mathcal{S}(Q) \subseteq \text{Tup}(\Sigma_Q)$ is a relation over Σ_Q, called the **answer set** for Q wrt. $\mathcal{S}$.
- $\mathcal{S}(Q)$ can be computed according to the inductive definition, starting with the innermost (atomic) subexpressions.
- Thus, the relational algebra has a **functional semantics**.

Summary: Inductive Definition of Expressions

Atomic Expressions

- For an arbitrary attribute A and a constant $a \in \text{dom}(A)$, the **constant relation** $A : \{a\}$ is an algebra expression.

  $\Sigma_{A:\{a\}} = [A]$ and $\mathcal{S}(A : \{a\}) = A : \{a\}$
- Every relation name R is an algebra expression.

  $\Sigma_R = \bar{X}$ and $\mathcal{S}(R) = \mathcal{S}(R)$.

Summary (Cont'd)

Compound Expressions

Assume algebra expressions Q_1, Q_2 that define Σ_{Q_1}, Σ_{Q_2}, $\mathcal{S}(Q_1)$, and $\mathcal{S}(Q_2)$.

Compound algebraic expressions are now formed by the following rules (corresponding to the algebra operators):

Union

If $\Sigma_{Q_1} = \Sigma_{Q_2}$, then $Q = (Q_1 \cup Q_2)$ is the **union** of Q_1 and Q_2.

$\Sigma_Q = \Sigma_{Q_1}$ and $\mathcal{S}(Q) = \mathcal{S}(Q_1) \cup \mathcal{S}(Q_2)$.

Difference

If $\Sigma_{Q_1} = \Sigma_{Q_2}$, then $Q = (Q_1 \setminus Q_2)$ is the **difference** of Q_1 and Q_2.

$\Sigma_Q = \Sigma_{Q_1}$ and $\mathcal{S}(Q) = \mathcal{S}(Q_1) \setminus \mathcal{S}(Q_2)$.

Projection

For $\emptyset \neq \bar{Y} \subseteq \Sigma_{Q_1}$, $Q = \pi[\bar{Y}](Q_1)$ is the **projection** of Q_1 to the attributes in $\bar{Y}$.

$\Sigma_Q = \bar{Y}$ and $\mathcal{S}(Q) = \pi[\bar{Y}](\mathcal{S}(Q_1))$.

Inductive Definition of Expressions (Cont'd)

Selection

For a selection condition α over Σ_{Q_1}, $Q = \sigma[\alpha]Q_1$ is the **selection** from Q_1 wrt. α.

$\Sigma_Q = \Sigma_{Q_1}$ and $\mathcal{S}(Q) = \sigma[\alpha](\mathcal{S}(Q_1))$.

Natural Join

$Q = (Q_1 \bowtie Q_2)$ is the **(natural) join** of Q_1 and Q_2.

$\Sigma_Q = \Sigma_{Q_1} \cup \Sigma_{Q_2}$ and $\mathcal{S}(Q) = \mathcal{S}(Q_1) \bowtie \mathcal{S}(Q_2)$.

Renaming

For $\Sigma_{Q_1} = \{A_1, \ldots, A_k\}$ and $\{B_1, \ldots, B_k\}$ a set of attributes, $\rho[A_1 \to B_1, \ldots, A_k \to B_k]Q_1$ is the **renaming** of Q_1

$\Sigma_Q = \{B_1, \ldots, B_k\}$ and $\mathcal{S}(Q) = \{\mu[A_1 \to B_1, \ldots, A_k \to B_k] \mid \mu \in \mathcal{S}(Q_1)\}$.

Example

Example 3.15

Professor(PNr, Name, Office), Course(CNr, Credits, CName)
teach(PNr, CNr), examine(PNr, CNr)

- *For each professor (name) determine the courses he gives (CName).*

$$\pi \text{ [Name, CName] ((Professor} \bowtie \text{teach)} \bowtie \text{Course)}$$

- *For each professor (name) determine the courses (CName) that he teaches, but that he does not examine.*

$$\begin{array}{l} \pi[\text{Name}, \text{CName}]((\\ (\pi[\text{Name}, \text{CNr}](\text{Professor} \bowtie \text{teach}))\\ \setminus\\ (\pi[\text{Name}, \text{CNr}](\text{Professor} \bowtie \text{examine}))\\ ) \bowtie \text{Course}) \end{array}$$

Simpler expression:

$$\pi \text{ [Name, CName] ((Professor} \bowtie \text{(teach} \setminus \text{examine))} \bowtie \text{Course)} \qquad \square$$

Equivalence of Expressions

Algebra expressions Q, Q' are called **equivalent**, $Q \equiv Q'$, if and only if for all structures $\mathcal{S}$, $\mathcal{S}(Q) = \mathcal{S}(Q')$.

Equivalence of expressions is the basis for **algebraic optimization**.

Let attr(α) the set of attributes that occur in a selection condition α, and $Q, Q_1, Q_2, \ldots$ expressions with formats X, X_1, $\ldots$.

Projections

- $\bar{Z}, \bar{Y} \subseteq \bar{X} \Rightarrow \pi[\bar{Z}](\pi[\bar{Y}](Q)) \equiv \pi[\bar{Z} \cap \bar{Y}](Q)$.
- $\bar{Z} \subseteq \bar{Y} \subseteq \bar{X} \Rightarrow \pi[\bar{Z}](\pi[\bar{Y}](Q)) \equiv \pi[\bar{Z}](Q)$.

Selections

- $\sigma[\alpha_1](\sigma[\alpha_2](Q)) \equiv \sigma[\alpha_2](\sigma[\alpha_1](Q)) \equiv \sigma[\alpha_1 \wedge \alpha_2](Q))$.
- attr$(\alpha) \subseteq \bar{Y} \subseteq \bar{X} \Rightarrow \pi[\bar{Y}](\sigma[\alpha](Q)) \equiv \sigma[\alpha](\pi[\bar{Y}](Q))$.

Joins

- $Q_1 \bowtie Q_2 \equiv Q_2 \bowtie Q_1$.
- $(Q_1 \bowtie Q_2) \bowtie Q_3 \equiv Q_1 \bowtie (Q_2 \bowtie Q_3)$.

Equivalence of Expressions (Cont'd)

Joins and other Operations

- attr$(\alpha) \subseteq \bar{X}_1 \cap \bar{X}_2 \Rightarrow \sigma[\alpha](Q_1 \bowtie Q_2) \equiv \sigma[\alpha](Q_1) \bowtie \sigma[\alpha](Q_2)$.
- attr$(\alpha) \subseteq \bar{X}_1$, attr$(\alpha) \cap \bar{X}_2 = \emptyset \Rightarrow \sigma[\alpha](Q_1 \bowtie Q_2) \equiv \sigma[\alpha](Q_1) \bowtie Q_2$.
- Assume $V \subseteq \overline{X_1 X_2}$ and let $W = \bar{X}_1 \cap \overline{V X_2}$, $U = \bar{X}_2 \cap \overline{V X_1}$.
  Then, $\pi[V](Q_1 \bowtie Q_2) = \pi[V](\pi[W](Q_1) \bowtie \pi[U](Q_2))$;
- $\bar{X}_2 = \bar{X}_3 \Rightarrow Q_1 \bowtie (Q_2 \text{ op } Q_3) = (Q_1 \bowtie Q_2) \text{ op } (Q_1 \bowtie Q_3)$ where op $\in \{\cup, -\}$.

Exercise 3.2
Prove some of the equalities (use the definitions given on the "Base Operators" slide). □

Expressive Power of the Algebra

Transitive Closure

The transitive closure of a binary relation R, denoted by R^* is defined as follows:

$$
\begin{aligned}
R^1 &= R \\
R^{n+1} &= \{(a,b) \mid \text{there is an } s \text{ s.t. } (a,x) \in R^n \text{ and } (x,b) \in R\} \\
R^* &= \bigcup_{1..\infty} R^n
\end{aligned}
$$

Examples:

- child(x,y): child* = descendant
- flight connections
- flows_into of rivers in Mondial

Theorem 3.2
There is no expression of the relational algebra that computes the transitive closure of arbitrary binary relations r. □

EXAMPLES

Time to play. Perhaps postpone examples after comparison with SQL (next subsections)

Aspects

- join as “extending” operation (cartesian product – “all pairs of X and Y such that ...”)
- equijoin as “restricting” operation
- natural join/equijoin in many cases along key/foreign key relationships
- relational division (in case of queries of the style “return all X that are in a given relation with all Y such that ...”)