Chapter 5 Relational Databases and SQL: Further Issues

- Data Definition Language (DDL): schema generation
- Data Manipulation Language (DML):
 - queries
 - insertions, deletions, modifications
- Database behavior?

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5.1 Database Schema

The database schema is the complete model of the structure of the application domain (here: relational schema):

- relations
 - names of attributes
 - domains of attributes
 - keys
- · additional constraints
 - value constraints
 - referential integrity constraints
- storage issues of the physical schema: indexes, clustering etc. also belong to the schema

5.1.1 Schema Generation in SQL

Definition of Tables

Basic form: attribute names and domains

```
CREATE TABLE 
      (<col> <datatype>,
       <col> <datatype>)
domains: NUMBER, CHAR(n), VARCHAR2(n), DATE ...
   CREATE TABLE City
         ( Name
                        VARCHAR2(35),
           Country
                       VARCHAR2(4),
          Province VARCHAR2(32),
           Population
                        NUMBER,
                        NUMBER,
           Longitude
           Latitude
                        NUMBER );
```

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Integrity constraints

Simple constraints on individual attributes are given with the attribute definitions as "column constraints":

- domain definitions are already integrity constraints
- further constraints on individual attribute values more detailed range restrictions:

```
City: CHECK (population > 0) or CHECK (longitude BETWEEN -180 AND 180)
```

- NULL values allowed? : Country: name NOT NULL
- Definition of key/uniqueness constraints:
 Country: code PRIMARY KEY or name UNIQUE

Integrity constraints (Cont'd)

Constraints on several attributes are given separately as "table constraints":

- table-constraints have a name
- · must state which columns are concerned

```
CREATE TABLE City
  ( Name VARCHAR2(35),
    Country VARCHAR2(4),
    Province VARCHAR2(32),
    Population NUMBER CONSTRAINT CityPop CHECK (Population >= 0),
    Longitude NUMBER CONSTRAINT CityLong CHECK (Longitude BETWEEN -180 AND 180),
    Latitude NUMBER CONSTRAINT CityLat CHECK (Latitude BETWEEN -90 AND 90),
    CONSTRAINT CityKey PRIMARY KEY (Name, Country, Province));
```

... for details see "Practical Training SQL".

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Integrity constraints (Cont'd)

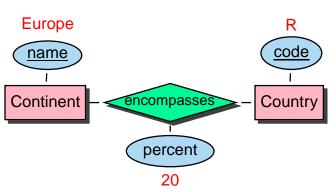
• up to now: only intra-table constraints

General Assertions

- inter-table constraints
 e.g., "sum of inhabitants of provinces equals the population of the country",
 "sum of inhabitants of all cities of a country must be smaller the than population of the country"
- SQL standard: CREATE ASSERTION
- not supported by most systems
- other solution: later

5.1.2 Referential Integrity Constraints

- important part of the schema
- relate foreign keys with their corresponding primary keys:



encompasses				
Country	Continent	Percent		
VARCHAR(4)	VARCHAR(20)	NUMBER		
R	Europe	20		
R	Asia	80		
D	Europe	100		

encompasses.country \rightarrow country.code and encompasses.continent \rightarrow continent.name

other examples:

city.country → country.code and country.(capital,province,code) → city.(name,province,country)

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Referential Integrity Constraints: SQL Syntax

- as column constraints (only single-column keys):
 <column-name> <datatype> REFERENCES (<column>)
- as table constraints (also compound keys):

```
CREATE TABLE is_member
```

CREATE TABLE City

```
( Name VARCHAR2(35),
  Country VARCHAR2(4) REFERENCES Country(Code),
  Province VARCHAR2(32),
  Population NUMBER ..., Longitude NUMBER ..., Latitude NUMBER ...,
  CONSTRAINT CityKey PRIMARY KEY (Name, Country, Province),
  FOREIGN KEY (Country, Province) REFERENCES Province (Country, Name) );
```

5.1.3 Virtual Tables: Views

Views are tables that are not materialized, but defined by a query against the database:

```
CREATE VIEW <name> AS <query>

CREATE OR REPLACE VIEW symm_borders AS

SELECT * FROM borders

UNION

SELECT Country2, Country1, Length FROM borders;

SELECT country2

FROM symm_borders

WHERE country1='D';
```

- classical views: the content of a view is always computed when it is queried.
- Materialized Views: view is materialized and automatically maintained

 → view maintenance problem: when a base table changes, what modifications have to be applied to which views?

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5.2 SQL: Data Manipulation Language

... everything is based on the structure of the SELECT-FROM-WHERE clause:

• Deletions:

```
DELETE FROM ... WHERE ...
```

Updates:

```
UPDATE 
SET <attribute> = <value>, ..., <attribute> = <value>
WHERE ...
```

value can be a subquery (also a correlated one)

Insertions:

```
INSERT INTO  VALUES (...)
INSERT INTO  (SELECT ... FROM ... WHERE ...)
```

SQL: THE DATE DATATYPE

- many applications in business and administration use dates
- computations on dates (e.g., "last of the third month after ...", "number of days between")
- ⇒ SQL provides comprehensive datatypes DATE, TIME, TIMESTAMP
 - semantics: year/month/date/hour/minute/second timestamp: additionally fractions of seconds as decimal (Oracle: only DATE and TIMESTAMP) built-in calendar knows about length of months, leap years etc.
 - operators on date and time:
 - date + days
 - MONTHS_BETWEEN($date_1, date_2$), ADD_MONTHS(date, n), LAST_DAY(date)
 - SYSDATE
 - syntax: different syntactical representations/formats can be chosen for input/output: (Oracle: SET NLS_DATE_FORMAT *string-pattern*, e.g. "DD.MM.YYYY", "DD-MON-YY", "DD hh:mm:ss")

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The DATE Datatype: Example

```
CREATE TABLE Politics
  ( Country VARCHAR2(4),
        Independence DATE,
        Government VARCHAR2(120));
ALTER SESSION SET NLS_DATE_FORMAT = 'DD MM YYYY';
INSERT INTO politics VALUES
      ('B','04 10 1830','constitutional monarchy');
```

All countries that have been founded between 1200 und 1600:

SELECT Country, Independence FROM Politics WHERE Independence BETWEEN '01 01 1200' AND '31 12 1599';

Land	Datum	
МС	01 01 1419	
NL	01 01 1579	
Е	01 01 1492	
THA	01 01 1238	

5.3 Beyond Relational Completeness

- The Relational Algebra and SQL are only relationally complete.
- can e.g. not compute the transitive closure of a relation
- applications require a more complex behavior:
 - SQL als the "core query language"
 - with something around it ...

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MAKING SQL TURING-COMPLETE

embedded SQL in C/Pascal:

```
EXEC SQL SELECT ... FROM ... WHERE ... embedded into Java: JDBC (Java Database Connectivity)
```

- SQL-92: Procedural Extensions to SQL:
 - CREATE procedures and functions as compiled things *inside* the database
 - standardized concepts, but product-specific syntax
 - basic programming constructs of a "typical" Turing-complete language:
 Variables, BEGIN ... END, IF ... THEN ... ELSIF ..., WHILE ... LOOP ..., FOR ... LOOP
 - SQL can be used inside PL/SQL statements

"IMPEDANCE MISMATCH" BETWEEN DB AND PROGRAMMING LANGUAGES

(cf. Slide 3)

Set-oriented (relations) vs. value-oriented (variables)

• how to handle the result of a query in C/Pascal/Java?

Iterators (common programming pattern for all kinds of collections)

- explicit:
 - new/init(<query>)/open()
 - first(), next(), isempty()
 - fetch() (into a record/tuple variable)
- implicit (PL/SQL's "Cursor FOR LOOP"):

```
FOR <record-variable> IN <query>
LOOP
   do something with <record-variable>
END LOOP;
```

... for details see "Practical Training SQL".

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5.4 Integrity Maintenance

- if a tuple is changed/inserted/deleted it is immediately checked whether all constraints in the current database state are satisfied afterwards.
 - Otherwise the operation is rejected.
- if a constraint is defined/enabled, it is immediately checked whether it is satisfied by the current database state.

Otherwise the operation is rejected.

Any further possibilities?

Integrity Maintenance (Cont'd): referential integrity

Consider again country - organization - is member:

is_member.organization → organization.abbrev $is_member.country \rightarrow country.code$

- deletion of a membership entry: no problem
- deletion of a country: any membership entries for it are now "dangling"
- ⇒ remove them!

Referential Actions

FOREIGN KEY is_member(country) REFERENCES country(code) ON DELETE CASCADE

- ON DELETE CASCADE: delete referencing tuple
- ON DELETE RESTRICT: referenced tuple cannot be deleted
- ON DELETE NO ACTION: referenced tuple can be deleted if the same transaction also deletes the referencing tuple
- ON DELETE SET NULL: foreign key of referencing tuple is set to NULL
- ON DELETE SET DEFAULT: foreign key of referencing tuple is set to a default value
- same for ON UPDATE

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Referential Actions

Country					
Name	Code	Capital	Province		
Germany	Ď	Berlin/	Berlin		
United States	u\$A	Washington	Distr. Columbia		
		·/··			
CASCADE					
NO ACTION					
			1.		
City					
Name	Country	Province	2.		
Berlin	D	Berlin			
Washington	USA	Distr. Colum	bia 3.		

- 1. DELETE FROM City WHERE Name='Berlin';
- 2. DELETE FROM Country WHERE Name='Germany';
- 3. UPDATE Country SET code='DE' WHERE code='D';

Referential Actions: Problems Country Capital Name Code Province D Germany Berlin Berlin CASCADE **United States** US Washington Distr.Col. **Province** Name Country Capital Berlin Berlin **CASCADE** D SET NULL Washington Distr.Col. US City Name Country **Province** Berlin D В Washington USA Distr.Col. DELETE FROM Country WHERE Code='D'

... ambiguous semantics!

See http://dbis.informatik.uni-goettingen.de/RefInt.

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... active behavior/reaction on events!

5.5 Active Databases/Triggers

- reacting on an event
 - external event/signal
 - internal event: modification/insertion/deletion
 - internal event: time
- if a condition is satisfied
- then do something/execute an action

ECA: Event-Condition-Action rules

ECA-Rules

Consider database updates only: one or more tuples of a table are changed.

- Granularity:
 - execute action once for "all updates together" (e.g., afterwards, update a sum)
 - execute action for each changed tuple (e.g. cascading update)
- Timepoint:
 - after execution of original update
 - before execution of original update
 - instead of original update
- Actions:
 - can read the before- and after value of the updated tuple
 - read and write other tables

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Triggers

The SQL standard provides "Triggers" for implementation of ECA rules:

CREATE TRIGGER

specify event:

```
ON {DELETE | UPDATE | INSERT} ON
```

- specify condition WHEN <condition>
- specify granularity FOR EACH STATEMENT | ROW
- specify action

Actions are programmed using the above-mentioned procedural extensions to SQL.

Applications

- implement application-specific business rules
- integrity maintenance
- monitoring of assertions

... for details see "Practical Training SQL".